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Threshold Graph Limits and Random Threshold Graphs

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Abstract. We study the limit theory of large threshold graphs and apply this to a variety of models for random threshold graphs. The results give a nice set of examples for the emerging theory of graph limits.

I. Introduction

I.I. Threshold Graphs

Graphs have important applications in modern systems biology and social sciences. Edges are created between interacting genes or people who know each other. However, graphs are not objects that are naturally amenable to simple statistical analysis. There is no natural average graph, for instance. Being able to predict or replace a graph by hidden (statisticians call them latent) real variables has many advantages. This paper studies such a class of graphs that sits within the larger class of interval graphs [McKee and McMorris 99], itself a subset of intersection graphs [Erdős et al. 89]; see also [Brandstädt et al. 99].

Consider the following properties of a simple graph G on $[n] := \{1, 2, \dots, n\}$:

(1.1) There are real weights w_i and a threshold value t such that there is an edge from i to j if and only if $w_i + w_j > t$. Thus "the rich people always know each other."

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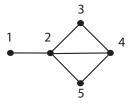


Figure 1. A threshold graph.

- (1.2) The graph G can be built sequentially from the empty graph by adding vertices one at a time, where each new vertex is either isolated (nonadjacent to all the previous) or dominant (connected to all the previous).
- (1.3) The graph is uniquely determined (as a labeled graph) by its degree sequence.
- (1.4) Any induced subgraph has either an isolated or a dominant vertex.
- (1.5) There is no induced subgraph $2K_2$, P_4 , or C_4 . (Equivalently, there is no alternating 4-cycle, i.e., four distinct vertices x, y, z, w with edges xy and zw but no edges yz and xw; the diagonals xz and yw may or may not exist.)

These properties are equivalent and define the class of threshold graphs. The book [Mahadev and Peled 95] contains proofs and several other seemingly different characterizations. Note that the complement of a threshold graph is a threshold graph (by any of (1.1)–(1.5)). By (1.2), a threshold graph is either connected (if the last vertex is dominant) or has an isolated vertex (if the last vertex is isolated); clearly these two possibilities exclude each other when n > 1.

Example 1.1. The graph in Figure 1 is a threshold graph, from (1.1) by taking weights 1, 5, 2, 3, 2 on vertices 1 through 5 with t = 4.5, or from (1.2) by adding vertices 3, 5 (isolated), 4 (dominant), 1 (isolated), and 2 (dominant).

While many familiar graphs are threshold graphs (stars and complete graphs, for example), many are not (e.g., paths and cycles of length 4 or more). For example, of the 64 labeled graphs on four vertices, 46 are threshold graphs; the other 18 are paths P_4 , cycles C_4 , and pairs of edges $2K_2$ (which is the complement of C_4). Considering unlabeled graphs, there are 11 graphs on four vertices, and eight of them are threshold graphs.









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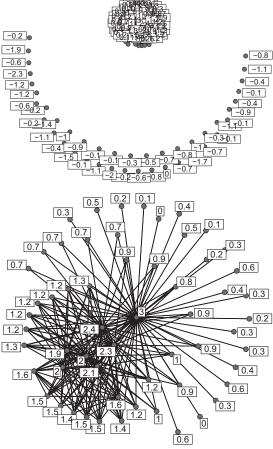


Figure 2. A whole threshold graph with isolates (above) and with only the connected part expanded (below); the labels are the rounded weights w_i .

1.2. Random Threshold Graphs

It is natural to study random threshold graphs. There are several different natural random constructions; we will in particular consider the following three:

- (1.6) From (1.1) by choosing $\{w_i\}_{1 \leq i \leq n}$ as independent and identically distributed (i.i.d.) random variables from some probability distribution. (We also choose some fixed t; we may assume t = 0 by replacing w_i by $w_i t/2$.)
- (1.7) From (1.2) by ordering the vertices randomly and adding the vertices one by one, each time choosing at random between the qualifiers "dominant" and "isolated" with probabilities p_i and $1 p_i$, respectively, $1 \le i \le n$.









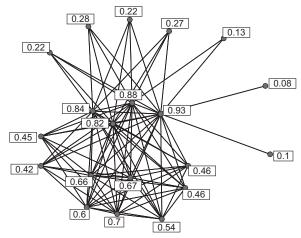


Figure 3. A threshold graph with n = 20 and uniform w_i . It turns out that this instance has no isolates. The labels are the rounded weights w_i .

This is a simple random attachment model in a similar vein to those in [Mitzenmacher 04]. We mainly consider the case that all p_i are equal to a single parameter $p \in [0, 1]$.

(1.8) The uniform distribution on the set of threshold graphs.

Example 1.2. Figure 2 shows a random threshold graph constructed by (1.6) with w_i chosen independently from the standardized normal distribution and t=3. About half of the vertices are isolated, most of those with negative weights.

Example 1.3. Figure 3 shows a random threshold graph constructed by (1.6) with w_i chosen as i.i.d. uniform random variables on [0, 1] and t = 1. This instance is connected; this happens if and only if the maximum and minimum of the w_i add to more than 1 (then there is a dominant vertex); in this example this has probability 1/2.

We show below (Corollaries 6.6 and 6.7) that this uniform-weight model is equivalent to adding isolated or dominant nodes as in (1.7) with probability p=1/2, independently and in random order. It follows that this same distribution appears as the stationary distribution of a Markov chain on threshold graphs that picks a vertex at random and changes it to dominant or isolated with probability 1/2 (this walk is analyzed in [Brown and Diaconis 98]). Furthermore, it follows from Section 2.1 that these models yield a uniform distribution on the set of unlabeled threshold graphs of order n.









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1.3. Bipartite Threshold Graphs

We also study the parallel case of bipartite threshold graphs (difference graphs), both for its own sake and because one of the main theorems is proved by first considering the bipartite case.

By a bipartite graph we mean a graph with an explicit bipartition of the vertex set; it can thus be written as (V_1, V_2, E) , where the edge set E satisfies $E \subseteq V_1 \times V_2$. The following properties of a bipartite graph are equivalent and define the class of bipartite threshold graphs. (See [Mahadev and Peled 95] for further characterizations.)

- (1.9) There are real weights w'_i , $i \in V_1$, and w''_j , $j \in V_2$, and a threshold value t such that there is an edge from i to j if and only if $w'_i + w''_j > t$.
- (1.10) The graph G can be built sequentially starting from n_1 white vertices and n_2 black vertices in some fixed total order. Proceeding in this order, make each white vertex dominant or isolated from all the black vertices that precede it and each black vertex dominant or isolated from all earlier white vertices.
- (1.11) Any induced subgraph has either an isolated vertex or a vertex dominating every vertex in the other part.
- (1.12) There is no induced subgraph $2K_2$.

Remark I.4. Threshold graphs were defined in [Chvátal and Hammer 77]. Bipartite threshold graphs were studied in [Hammer et al. 90] under the name difference graphs because they can equivalently be characterized as the graphs (V, E) for which there exist weights w_v , $v \in V$, and a real number t such that $|w_v| < t$ for every v and $uv \in E \iff |w_u - w_v| > t$; it is easily seen that every such graph is bipartite with $V_1 = \{v : w_v \geq 0\}$ and $V_2 = \{v : w_v < 0\}$ and that this satisfies the definition above (e.g., with $w'_v = w_v$ and $w''_v = -w_v$), and conversely. We will use the name bipartite threshold graph to emphasize that we consider these graphs equipped with a given bipartition. The same graphs were called chain graphs in [Yannakakis 82] because each partition can be linearly ordered for the inclusion of the neighborhoods of its elements.

Remark 1.5. A suite of programs for working with threshold graphs appears in [Hagberg et al. 06] with further developments in [Konno et al. 05] and [Masuda et al. 05].









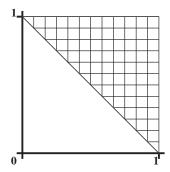


Figure 4. The function W(x,y) for Example 1.3. Hashed values have W(x,y)=1; unhashed, W(x,y)=0.

Remark I.6. The most natural class of graphs built from a coordinate system are commonly called geometric graphs [Penrose 03] or geographical graphs [Konno et al. 05, Masuda et al. 05]. Threshold graphs are a special case of these. Their recognition and manipulation in a statistical context relies on useful measures on such graphs. We will start by defining such measures and developing a limit theory.

1.4. Overview of the Paper

The purpose of this paper is to study the limiting properties of large threshold graphs in the spirit of the theory of graph limits developed in [Lovász and Szegedy 06] and [Borgs et al. 07b] (and in further papers by those authors and others). As explained below, the limiting objects are not graphs, but can rather be represented by symmetric functions W(x, y) from $[0, 1]^2$ to [0, 1]; any sequence of graphs that converges in the appropriate way has such a limit.

Conversely, such a function W may be used to form a random graph G_n by choosing independent random points U_i in [0,1], and then for each pair (i,j) with $1 \le i < j \le n$ flipping a biased coin with heads probability $W(U_i, U_j)$, putting an edge from i to j if the coin comes up heads. The resulting sequence of random graphs is (almost surely) an example of a sequence of graphs converging to W. For Example 1.3, letting $n \to \infty$, there is (as we show in greater generality in Section 6) a limit W that may be pictured as in Figure 4.

One of our main results (Theorem 5.5) shows that graph limits of threshold graphs have unique representations by increasing symmetric zero—one-valued functions W. Furthermore, there is a one-to-one correspondence between these limiting objects and a certain type of "symmetric" probability distribution P_W on [0,1]. A threshold graph is characterized by its degree sequence; normalizing





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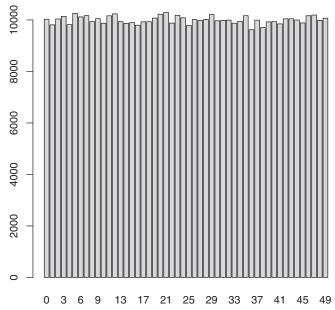


Figure 5. Threshold graphs generated with n = 50 as in Example 1.3 with uniform w_i and t = 1; this is the degree histogram for a sample of 10,000 random graphs.

this to be a probability distribution, say $\nu(G_n)$, we show (Theorem 5.8) that a sequence of threshold graphs converges to W when $n \to \infty$ if and only if $\nu(G_n)$ converges to P_W . (Hence P_W can be regarded as the degree distribution of the limit. The result that a limit of threshold graphs is determined by its degree distribution is a natural analogue for the limit objects of the fact that an unlabeled threshold graph is uniquely determined by its degree distribution.)

Figures 5 and 7 show simulations of these results. In Figure 5, ten thousand graphs with n=50 were generated from (1.6) with uniform weights as in Example 1.3.

In the bipartite case, there is a similar one-to-one correspondence between the limit objects and probability distributions on [0,1]; now all probability distributions on [0,1] appear in the representation of the limits (Theorem 5.1).

Section 2 discusses uniform random threshold graphs (both labeled and unlabeled) and methods to generate them. Section 3 gives a succinct review of notation and graph limits. Section 4 develops the limit theory of degree sequences; this is not restricted to threshold graphs. Section 5 develops the limit theory for threshold graphs both deterministic and random. Section 6 treats examples of random threshold graphs and their limits, and Section 8 gives corresponding examples and results for random bipartite threshold graphs. Section 7









treats the degree distribution of uniform random threshold graphs in greater detail. Section 9 treats the spectrum of the Laplacian of threshold graphs.

We denote the vertex and edge sets of a graph G by V(G) and E(G), and the numbers of vertices and edges by v(G) := |V(G)| and e(G) := |E(G)|. For a bipartite graph we similarly use $V_i(G)$ and $v_i(G)$, i = 1, 2.

Throughout the paper, "increasing" and "decreasing" should be interpreted in the weak sense (nondecreasing and nonincreasing). Unspecified limits are to be taken as $n \to \infty$.

2. Generating Threshold Graphs Uniformly

This section gives algorithms for generating uniformly distributed threshold graphs, in both the labeled case and unlabeled cases. The algorithms are used here for simulation and in Sections 6 and 7 to prove limit theorems.

Let \mathcal{T}_n and \mathcal{LT}_n be the sets of unlabeled and labeled threshold graphs on n vertices. These are different objects; \mathcal{T}_n is a quotient of \mathcal{LT}_n , and we treat counting and uniform generation separately for the two cases. We assume in this section that $n \geq 2$.

2.1. Unlabeled Threshold Graphs

We can code an unlabeled threshold graph on n vertices by a binary code $\alpha_2 \cdots \alpha_n$ of length n-1: Given a code $\alpha_2 \cdots \alpha_n$, we construct G by (1.2), adding vertex i as a dominant vertex if and only if $\alpha_i = 1$ ($i \geq 2$). Conversely, given G of order $n \geq 2$, let $\alpha_n = 1$ if there is a dominant vertex (G is connected) and $\alpha_n = 0$ if there is an isolated vertex (G is disconnected). We then remove one such dominant or isolated vertex and continue recursively to define $\alpha_{n-1}, \ldots, \alpha_2$.

Since all dominant (isolated) vertices are equivalent to each other, this coding gives a bijection between \mathcal{T}_n and $\{0,1\}^{n-1}$. In particular,

$$|\mathcal{T}_n| = 2^{n-1}, \qquad n \ge 1.$$

See Figure 6 for an example.



Figure 6. The four graphs in \mathcal{T}_3 and their codes.









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Algorithm I. (Generating uniform random unlabeled threshold graphs of a given order n.)

1. Add n vertices by (1.2), each time randomly choosing "isolated" or "dominant" with probability 1/2.

This leads to a simple algorithm to generate a uniformly distributed random unlabeled threshold graph: we construct a random code by making n-1 coin flips, according to Algorithm 1.

This is thus the same as the second method in Example 1.3, so Corollary 6.6 shows that the first method in Example 1.3 also yields uniform random unlabeled threshold graphs (if we forget the labels).

The following notation is used to define two further algorithms (Section 2.3) and for proof of the limiting results in Section 7.

Define the extended binary code of a threshold graph to be the binary code with the first binary digit repeated; it is thus $\alpha_1\alpha_2\alpha_3\cdots\alpha_n$ with $\alpha_1:=\alpha_2$. The runs of 0's and 1's in the extended binary code then correspond to blocks of vertices that can be added together in (1.2) as either isolated or dominant vertices, with the blocks alternating between isolated and dominant.

The vertices in each block are equivalent and have in particular the same vertex degrees, while vertices in different blocks can be seen to have different degrees. (The degree increases strictly from one dominant block to the next and decreases strictly from one isolated block to the next, with every dominant block having higher degree than every isolated block; cf. Example 2.1 below.) The number of different vertex degrees thus equals the number of blocks.

If the lengths of the blocks are $b_1, b_2, \ldots, b_{\tau}$, then the number of automorphisms of G is $\prod_{j=1}^{\tau} b_j!$, since the vertices in each block may be permuted arbitrarily.

Note that if b_1, \ldots, b_{τ} are the lengths of the blocks, then

$$b_1 \ge 2, \quad b_k \ge 1 \ (k \ge 2), \quad \sum_{k=1}^{\tau} b_k = n.$$
 (2.1)

Since the blocks are alternately dominant or isolated, and the first block may be either, each sequence b_1, \ldots, b_{τ} satisfying (2.1) corresponds to exactly two unlabeled threshold graphs of order n. (These graphs are the complements of each other. One has isolated blocks where the other has dominant blocks.)



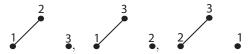






2.2. Labeled Threshold Graphs

The situation is different for labeled threshold graphs. For example, all of the $2^{\binom{3}{2}} = 8$ labeled graphs with n = 3 turn out to be threshold graphs, and for instance, the graphs



are distinguished. Hence the distribution of a uniform random labeled threshold graph differs from the distribution of a uniform unlabeled threshold graph (even if we forget the labels). In particular, Example 1.3 does not produce uniform random labeled threshold graphs.

Let G be an unlabeled threshold graph with an extended code having block lengths (runs) b_1, \ldots, b_{τ} . Then the number of labeled threshold graphs corresponding to G is $n!/\prod_1^{\tau} b_j!$, since every such graph corresponds to a unique assignment of the labels $1, \ldots, n$ to the τ blocks, with b_i labels to block i. (Alternatively and equivalently, this follows from the number $\prod_1^{\tau} b_j!$ of automorphisms given above.)

The number $t(n) := |\mathcal{LT}_n|$ of labeled threshold graphs [Sloane 09, A005840] has been studied in [Beissinger and Peled 87]. Among other things, the authors show that

$$\sum_{n=0}^{\infty} t(n) \frac{x^n}{n!} = \frac{(1-x)e^x}{2-e^x},$$

so by Taylor expansion,

and by expanding the singularities (cf. [Flajolet and Sedgewick 09, Chapter IV]), they obtain the exact formula

$$\frac{t(n)}{n!} = \sum_{k=-\infty}^{\infty} \left(\frac{1}{\log 2 + 2\pi i k} - 1\right) \left(\frac{1}{\log 2 + 2\pi i k}\right)^n, \quad n \ge 2,$$
 (2.2)

where the leading term is the one with k = 0, and thus one has the asymptotics

$$\frac{t(n)}{n!} = \left(\frac{1}{\log 2} - 1\right) \left(\frac{1}{\log 2}\right)^n + \epsilon(n), \quad |\epsilon(n)| \le \frac{2\zeta(n)}{(2\pi)^n},\tag{2.3}$$

where $\zeta(n)$ is the Riemann zeta function and thus $\zeta(n) \to 1$. Furthermore,

$$t(n) = 2R_n - 2nR_{n-1}, \ n \ge 2, \text{ with } R_n = \sum_{k=1}^n k! S(n,k) = \sum_{\ell=0}^\infty \frac{\ell^n}{2^{\ell+1}},$$
 (2.4)







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where S(n,k) are Stirling numbers; R_n is the number of preferential arrangements of n labeled elements, or the number of weak orders on n labeled elements [Sloane 09, A000670], also called surjection numbers [Flajolet and Sedgewick 09, II.3]. (This is easily seen using the blocks above; the number of labeled threshold graphs with a given sequence of blocks is twice—since the first block may be either isolated or dominant—the number of preferential arrangements with the same block sizes; if we did not require $b_1 \geq 2$, this would yield $2R_n$, but we have to subtract twice the number of preferential arrangements with $b_1 = 1$, which is $2nR_{n-1}$.) We note for future use the generating function [Flajolet and Sedgewick 09, (II.15)]

$$\sum_{n=0}^{\infty} R_n \frac{x^n}{n!} = \frac{1}{2 - e^x}.$$
 (2.5)

Let t(n, j) be the number of labeled threshold graphs with j isolated points. Then, as also shown in [Beissinger and Peled 87] (and easily seen), for $n \geq 2$,

$$t(n,0) = t(n)/2,$$

$$t(n,j) = \begin{cases} \binom{n}{j}t(n-j,0) = \frac{1}{2}\binom{n}{j}t(n-j), & 0 \le j \le n-2, \\ 0, & j = n-1, \\ 1, & j = n. \end{cases}$$
 (2.6)

Thus knowledge of t(n) provides t(n, j).

These ingredients allow us to give an algorithm, Algorithm 2, for choosing uniformly in \mathcal{LT}_n .

Alternatively, instead of selecting the subsets in steps 1 and 2 at random, we may choose them in any way, provided the algorithm begins or ends with a random permutation of the points.

The algorithm works because of a characterization of threshold graphs by Chvátal and Hammer [Chvátal and Hammer 77], cf. (1.4): a graph is a threshold graph iff any subset S of vertices contains at least one isolated or one dominant vertex (within the graph induced by S). Thus in step 2, since there are no isolates among the n' vertices left, there must be at least one dominant vertex. (Note that j_0 may be zero, but not j_1, j_2, \ldots) The probability distribution for the number of dominant vertices follows the same law as that of the isolates because the complement of a threshold graph is a threshold graph (or because of the interchangeability of 0's and 1's in the binary coding given earlier in this section).

Note that this algorithm treats vertices in the reverse of the order in (1.2), where we add vertices instead of peeling them off, as here. It follows that we obtain the extended binary code of the graph by taking runs of j_0 0's, j_1 1's,









Algorithm 2. (Generating uniform random labeled threshold graphs of a given order n.)

- 0. Make a list of t(k) for k between 1 and n. Make lists of t(k, j) for k = 1, ..., n and j = 0, ..., k.
- 1. Choose an integer j_0 in $\{0, \ldots, n\}$ with probability that $j_0 = j$ given by t(n, j)/t(n). Choose (at random) a subset of j_0 points in $\{1, \ldots, n\}$. These are the isolated vertices in the graph. Let $n' := n j_0$ be the number of remaining points. If n' = 0 then stop.
- 2. Choose an integer j_1 in $\{1, \ldots, n'\}$ with probability that $j_1 = j$ given by t(n', j)/(t(n') t(n', 0)) = 2t(n', j)/t(n') and choose (at random) j_1 points of those remaining; these will dominate all further points, so add edges between these vertices and from them to all remaining points. Update n' to $n' j_1$, the number of remaining points. If n' = 0 then stop.
- 3. Choose an integer j_2 in $\{1, \ldots, n'\}$ with probability that $j_2 = j$ given by 2t(n', j)/t(n') and choose (at random) j_2 points of those remaining; these will be isolated among the remaining points, so no further edges are added. Update n' to $n' j_1$, the number of remaining points. If n' = 0 then stop.
- 4. Repeat from step 2 with the remaining n' points.

 j_2 0's, and so on, and then reversing the order. Hence, in the notation used above, the sequence (b_k) equals (j_k) in reverse order, ignoring j_0 if $j_0 = 0$. (In particular, note that the last j_k is greater than or equal to 2, since t(n', n'-1) = 0 for $n' \geq 2$, which corresponds to the first block $b_1 \geq 2$.)

Example 2.1. A sequence of j's generated for a threshold graph of size 20 is 0 2 3 1 1 1 3 1 1 3 1 1 2, which yields the sequence d d i i i d i d i i i d i d i d i i i d d d d i d i i of dominant and isolated vertices. A random permutation of $\{1, \ldots, 20\}$ was generated and we obtain

where d signifies that the vertex is connected to all later vertices in this list. The degree sequence is thus, taking the vertices in this order: 19, 19, 2, 2, 2, 16, 3, 15, 4, 4, 4, 12, 5, 11, 11, 11, 8, 10, 9, 9. The extended binary code 00101110100010100011 is obtained by translating i to 0 and d to 1, and reversing the order.











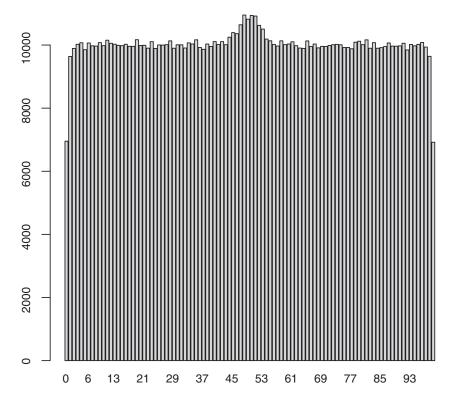


Figure 7. Threshold graphs were generated according to the algorithm of this section; this is the degree histogram.

Remark 2.2. Note that the last j is at least 2, since t(n', n'-1) = 0 for every $n' \geq 2$. Hence, the sequence of d's and i's always ends with at least two identical symbols.

Note that the vertex degrees are constant in each block of vertices assigned in one of the steps, i.e., for each run of i's or d's, and that they decrease from each run of d's to the next and increase from each run of i's to the next, with each vertex labeled d having a higher degree than every vertex labeled with i; the number of different vertex degrees is thus equal to the number of j's chosen by the algorithm, which equals the number of runs in the sequence of d's and i's (or in the extended binary code).

In Figure 7, ten thousand graphs were generated with n=100 according to the uniform distribution over all labeled threshold graphs. We discuss the central "bump" and other features of Figure 7 in Theorem 7.4.









2.3. The Distribution of Block Lengths

We have seen in Section 2.1 that if b_1, \ldots, b_{τ} are the lengths of the blocks of isolated or dominant vertices added to the graph in building it as in (1.2), then (2.1) holds.

Consider now a sequence of independent integer random variables B_1, B_2, \ldots with $B_1 \geq 2$ and $B_j \geq 1$ for $j \geq 2$, and let $S_k := \sum_{j=1}^k B_j$ be the partial sums. If some $S_{\tau} = n$, then stop and output the sequence (B_1, \ldots, B_{τ}) . Conditioning on the event that $S_{\tau} = n$ for some τ , this yields a random sequence b_1, \ldots, b_{τ} satisfying (2.1), and the probability that we obtain a given sequence $(b_j)_1^{\tau}$ equals $c \prod_{j=1}^{\tau} \mathbb{P}(B_j = b_j)$ for some normalizing constant c.

We now specialize to the case that $B_1 \stackrel{\mathrm{d}}{=} (B^* \mid B^* \geq 2)$ and $B_j \stackrel{\mathrm{d}}{=} (B^* \mid B^* \geq 1)$ for $j \geq 2$, for some given random variable B^* . Then the (conditional) probability of obtaining a given b_1, \ldots, b_{τ} satisfying (2.1) can be written

$$c' \prod_{i=1}^{\tau} \frac{\mathbb{P}(B^* = b_j)}{\mathbb{P}(B^* \ge 1)}$$

$$(2.7)$$

(with $c' = c \mathbb{P}(B^* \ge 1) / \mathbb{P}(B^* \ge 2)$).

There are two important cases. First, if we take $B^* \sim \text{Ge}(1/2)$, then $\mathbb{P}(B^* = b_j)/\mathbb{P}(B^* \geq 1) = 2^{-b_j}$, and thus (2.7) yields $c'2^{-\sum_j b_j} = c'2^{-n}$, so the probability is the same for all allowed sequences. Hence, in this case the distribution of the constructed sequence is uniform on the set of sequences satisfying (2.1), so it equals the distribution of block lengths for a random unlabeled threshold graph of size n.

The other case is $B^* \sim \text{Po}(\log 2)$. Then $\mathbb{P}(B^* \geq 1) = 1 - e^{-\log 2} = 1/2$, and $\mathbb{P}(B^* = b_j)/\mathbb{P}(B^* \geq 1) = (\log 2)^{b_j}/b_j!$. Thus, (2.7) yields the probability $c'(\log 2)^n/\prod_j b_j!$, which is proportional to the number $2 \cdot n!/\prod_j b_j!$ of labeled threshold graphs with block lengths b_1, \ldots, b_τ . Hence, in this case the distribution of the constructed sequence equals the distribution of block lengths for a random labeled threshold graph of size n.

We have proved the following result.

Theorem 2.3. Construct a random sequence B_1, \ldots, B_{τ} as above, based on a random variable B^* , stopping when $\sum_{1}^{\tau} B_j \geq n$ and conditioning on $\sum_{1}^{\tau} B_j = n$.

- (i) If $B^* \sim \text{Ge}(1/2)$, then (B_1, \ldots, B_{τ}) has the same distribution as the block lengths in a random unlabeled threshold graph of order n.
- (ii) If $B^* \sim \text{Po}(\log 2)$, then (B_1, \ldots, B_{τ}) has the same distribution as the block lengths in a random labeled threshold graph of order n.









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Algorithm 3. (Generating uniform unlabeled or labeled threshold graphs of a given order n.)

- 1. In the unlabeled case, let $B^* \sim \text{Ge}(1/2)$. In the labeled case, let $B^* \sim \text{Po}(\log 2)$.
- 2. Choose independent random numbers $B_1, B_2, \ldots, B_{\tau}$, with $B_1 \stackrel{\mathrm{d}}{=} (B^* \mid B^* \geq 2)$ and $B_j \stackrel{\mathrm{d}}{=} (B^* \mid B^* \geq 1), j \geq 2$, until the sum[2] $\sum_{1}^{\tau} B_j$ is greater than or equal to n.
- 3. If $\sum_{i=1}^{\tau} B_i > n$, start again with step 2.
- 4. We have found B_1, \ldots, B_{τ} with $\sum_{1}^{\tau} B_j = n$. Toss a coin to decide whether the first block is isolated or dominant; the following blocks alternate. Construct a threshold graph by adding vertices as in (1.2), block by block.
- 5. In the labeled case, make a random labeling of the graph.

It follows that the length of a typical (for example, a random) block converges in distribution to $(B^* \mid B^* \geq 1)$. Theorem 2.3 also leads to another algorithm to construct uniform random threshold graphs (Algorithm 3).

By standard renewal theory, the probability that $\sum_{1}^{\tau} B_{j}$ is exactly n is asymptotically $1/\mathbb{E}(B^{*} \mid B^{*} \geq 1) = \mathbb{P}(B^{*} \geq 1)/\mathbb{E}(B^{*})$, which is 1/2 in the unlabeled case and $1/(2\log 2) \approx 0.72$ in the labeled case, so we do not have to do very many restarts in step 3.

3. Graph Limits

This section reviews needed tools from the emerging field of graph limits.

3.1. Graph Limits

Here we review briefly the theory of graph limits as described in [Lovász and Szegedy 06], [Borgs et al. 07b], and [Diaconis and Janson 08].

If F and G are two graphs, let t(F,G) be the probability that a random mapping $\phi: V(F) \to V(G)$ defines a graph homomorphism, i.e., that $\phi(v)\phi(w) \in E(G)$ when $vw \in E(F)$. (By a random mapping we mean a mapping uniformly chosen among all $v(G)^{v(F)}$ possible ones; the images of the vertices in F are thus independent and uniformly distributed over V(G), i.e., they are obtained by random sampling with replacement.)









The basic definition is that a sequence G_n of (generally unlabeled) graphs converges if $t(F, G_n)$ converges for every graph F; as in [Diaconis and Janson 08], we will further assume $v(G_n) \to \infty$.

More precisely, the (countable and discrete) set \mathcal{U} of all unlabeled graphs can be embedded in a compact metric space $\overline{\mathcal{U}}$ such that a sequence $G_n \in \mathcal{U}$ of graphs with $v(G_n) \to \infty$ converges in $\overline{\mathcal{U}}$ to some limit $\Gamma \in \overline{\mathcal{U}}$ if and only if $t(F, G_n)$ converges for every graph F (see [Lovász and Szegedy 06, Borgs et al. 07b, Diaconis and Janson 08]).

Let $\mathcal{U}_{\infty} := \overline{\mathcal{U}} \setminus \mathcal{U}$ be the set of proper limit elements; we call the elements of \mathcal{U}_{∞} graph limits. The functionals $t(F,\cdot)$ extend to continuous functions on $\overline{\mathcal{U}}$, so $G_n \to \Gamma \in \mathcal{U}_{\infty}$ if and only if $v(G_n) \to \infty$ and $t(F,G_n) \to t(F,\Gamma)$ for every graph F.

Let W be the set of all measurable functions $W:[0,1]^2 \to [0,1]$ and let W_s be the subset of symmetric functions. The main result of [Lovász and Szegedy 06] is that every element of \mathcal{U}_{∞} can be represented by a (nonunique) function $W \in \mathcal{W}_s$. We let $\Gamma_W \in \mathcal{U}_{\infty}$ denote the graph limit defined by W. (We sometimes use the notation $\Gamma(W)$ for readability.) Then, for every graph F,

$$t(F, \Gamma_W) = \int_{[0,1]^{v(F)}} \prod_{ij \in E(F)} W(x_i, x_j) \, \mathrm{d}x_1 \cdots \, \mathrm{d}x_{v(F)}. \tag{3.1}$$

Moreover, define, for every $n \geq 1$, a random graph G(n, W) as follows: first choose a sequence X_1, X_2, \ldots, X_n of i.i.d. random variables uniformly distributed on [0,1], and then, given this sequence, for each pair (i,j) with i < j draw an edge ij with probability $W(X_i, X_j)$, independently for all pairs (i,j) with i < j. Then the random graph G(n, W) converges to Γ_W a.s. as $n \to \infty$.

If G is a graph, with $V(G) = \{1, \ldots, v(G)\}$ for simplicity, we define a function $W_G \in \mathcal{W}_s$ by partitioning [0,1] into v(G) intervals I_i , $i=1,\ldots,v(G)$, and letting W_G be the indicator $\mathbf{1}[ij \in E(G)]$ on $I_i \times I_j$. (In other words, W_G is a step function corresponding to the adjacency matrix of G.) We let $\pi(G) := \Gamma(W_G)$ denote the corresponding object in \mathcal{U}_∞ . It follows easily from (3.1) that $t(F,\pi(G)) = t(F,G)$ for every graph F. In particular, if G_n is a sequence of graphs with $v(G_n) \to \infty$, then G_n converges to some graph limit Γ if and only if $\pi(G_n) \to \Gamma$ in \mathcal{U}_∞ . (Unlike [Lovász and Szegedy 06] and [Borgs et al. 07b], we distinguish between graphs and limit objects and we do not identify G and $\pi(G)$; see [Diaconis and Janson 08].)

3.2. Bipartite Graphs and Their Limits

In the bipartite case, there are analogous definitions and results (see [Diaconis and Janson 08] for further details). We define a bipartite graph to be a graph







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and only if $t(F, G_n)$ converges for every bipartite graph F.

G with an explicit bipartition $V(G) = V_1(G) \cup V_2(G)$ of the vertex set such that the edge set E(G) is a subset of $V_1(G) \times V_2(G)$. Then we define t(F,G) in the same way as above but now for bipartite graphs F, by letting $\phi = (\phi_1, \phi_2)$ be a pair of random mappings $\phi_j : V_j(F) \to V_j(G)$. We let \mathcal{B} be the set of all unlabeled bipartite graphs and embed \mathcal{B} in a compact metric space $\overline{\mathcal{B}}$. A sequence (G_n) of bipartite graphs with $v_1(G_n), v_2(G_n) \to \infty$ converges in $\overline{\mathcal{B}}$ if

Let $\mathcal{B}_{\infty\infty}$ be the (compact) set of all such limits; we call the elements of $\mathcal{B}_{\infty\infty}$ bipartite graph limits. Every element of $\mathcal{B}_{\infty\infty}$ can be represented by a (nonunique) function $W \in \mathcal{W}$. We let $\Gamma_W'' \in \mathcal{B}_{\infty\infty}$ denote the element represented by W and have for every bipartite F,

$$t(F, \Gamma_W'') = \int_{[0,1]^{v_1(F) + v_2(F)}} \prod_{ij \in E(F)} W(x_i, y_j) \, \mathrm{d}x_1 \cdots \, \mathrm{d}x_{v_1(F)} \, \mathrm{d}y_1 \cdots \, \mathrm{d}y_{v_2(F)}.$$
(3.2)

Given $W \in \mathcal{W}$ and $n_1, n_2 \geq 1$, we define a random bipartite graph $G(n_1, n_2, W)$ by an analogue of the construction in Section 3.1: first choose two sequences $X_1, X_2, \ldots, X_{n_1}$ and $Y_1, Y_2, \ldots, Y_{n_2}$ of i.i.d. random variables uniformly distributed on [0, 1], and then, given these sequences, for each pair (i, j) draw an edge ij with probability $W(X_i, Y_j)$, independently for all pairs $(i, j) \in [n_1 \times [n_2]]$.

If G is a bipartite graph we define $W_G \in \mathcal{W}$ similarly as above (in general, with different numbers of steps in the two variables; note that W_G now in general is not symmetric) and let $\pi(G) := \Gamma''(W_G)$. Then by (3.2), $t(F, \pi(G)) = t(F, G)$ for every bipartite graph F. Hence if G_n is a sequence of bipartite graphs with $v_1(G_n), v_2(G_n) \to \infty$, then G_n converges to some bipartite graph limit Γ if and only if $\pi(G_n) \to \Gamma$ in $\mathcal{B}_{\infty\infty}$.

3.3. Cut Distance

The authors of [Borgs et al. 07b, Section 3.4] define a (pseudo)metric δ_{\square} on \mathcal{W}_{s} called the *cut distance*. This is only a pseudometric, since two different functions in \mathcal{W}_{s} may have cut distance 0 (for example, if one is obtained by a measure-preserving transformation of the other; see further [Borgs et al. 07a] and [Diaconis and Janson 08]), and it is shown in [Borgs et al. 07b] that in fact, $\delta_{\square}(W_1, W_2) = 0$ if and only if $t(F, W_1) = t(F, W_2)$ for every graph F, i.e., if and only if $\Gamma_{W_1} = \Gamma_{W_2}$ in \mathcal{U}_{∞} . Moreover, the quotient space $\mathcal{W}_{s}/\delta_{\square}$, where we identify elements of \mathcal{W}_{s} with cut distance 0, is a compact metric space and the mapping $W \mapsto \Gamma_{W}$ is a homeomorphism of $\mathcal{W}_{s}/\delta_{\square}$ onto \mathcal{U}_{∞} .

This extends to the bipartite case. In this case, we define δ''_{\square} on \mathcal{W} as δ_{\square} is defined in [Borgs et al. 07b, Section 3.4], but allowing different measure-preserving mappings for the two coordinates. Then if we identify elements in \mathcal{W} with cut









distance $0, W \mapsto \Gamma_W''$ becomes a homeomorphism of $\mathcal{W}/\delta_{\square}''$ onto $\mathcal{B}_{\infty\infty}$. Instead of repeating and modifying the complicated proofs from [Borgs et al. 07b], one can use their result in the symmetric case and define an embedding $W \mapsto \widetilde{W}$ of \mathcal{W} into \mathcal{W}_s by

$$\widetilde{W}(x,y) = \begin{cases} 0, & x < \frac{1}{2}, \ y < \frac{1}{2}; \\ 1, & x > \frac{1}{2}, \ y > \frac{1}{2}; \\ \frac{1}{4} + \frac{1}{2}W(2x - 1, 2y), & x > \frac{1}{2}, \ y < \frac{1}{2}; \\ \frac{1}{4} + \frac{1}{2}W(2y - 1, 2x), & x < \frac{1}{2}, \ y > \frac{1}{2}. \end{cases}$$

It is easily seen that $\delta''_{\square}(W_1,W_2)$ and $\delta_{\square}(\widetilde{W}_1,\widetilde{W}_2)$ are equal within some constant factors, for $W_1,W_2\in\mathcal{W}$, and that for each graph F, $t(F,\widetilde{W})$ is a linear combination of $t(F_i,W)$ for a family of bipartite graphs F (obtained by partitioning V(F) and erasing edges within the two parts). This and the results in [Borgs et al. 07b], together with the simple fact that $W\mapsto t(F,W)$ is continuous for δ''_{\square} for every bipartite graph F, imply easily the result claimed.

3.4. A Reflection Involution

If G is a bipartite graph, let G^{\dagger} be the graph obtained by interchanging the order of the two vertex sets; thus $V_j(G^{\dagger}) = V_{3-j}(G)$ and $E(G^{\dagger}) = \{uv : vu \in E(G)\}$. We say that G^{\dagger} is the reflection of G. Obviously, $t(F, G^{\dagger}) = t(F^{\dagger}, G)$ for any bipartite graphs F and G. It follows that if $G_n \to \Gamma \in \overline{\mathcal{B}}$, then $G_n^{\dagger} \to \Gamma^{\dagger}$ for some $\Gamma^{\dagger} \in \overline{\mathcal{B}}$, and this defines a continuous map of $\overline{\mathcal{B}}$ onto itself that extends the map just defined for bipartite graphs. We have, by continuity,

$$t(F, \Gamma^{\dagger}) = t(F^{\dagger}, \Gamma), \quad F \in \mathcal{B}, G \in \overline{\mathcal{B}}.$$
 (3.3)

Furthermore, $\Gamma^{\dagger\dagger} = \Gamma$, so the map is an involution, and it maps $\mathcal{B}_{\infty\infty}$ onto itself. For a function W on $[0,1]^2$, let $W^{\dagger}(x,y) := W(y,x)$ be its reflection in the main diagonal. It follows from (3.2) and (3.3) that $\Gamma''(W^{\dagger}) = \Gamma''(W)^{\dagger}$.

3.5. Threshold Graph Limits

Let $\mathcal{T} := \bigcup_{n=1}^{\infty} \mathcal{T}_n$ be the family of all (unlabeled) threshold graphs. Thus \mathcal{T} is a subset of the family \mathcal{U} of all unlabeled graphs, and we define $\overline{\mathcal{T}}$ as the closure of \mathcal{T} in $\overline{\mathcal{U}}$, and $\mathcal{T}_{\infty} := \overline{\mathcal{T}} \setminus \mathcal{T} = \overline{\mathcal{T}} \cap \mathcal{U}_{\infty}$, i.e., the set of proper limits of sequences of threshold graphs; we call these threshold graph limits.

In the bipartite case, we similarly consider the set $T'' := \bigcup_{n_1,n_2 \geq 1} \mathcal{T}_{n_1,n_2} \subset \mathcal{B}$ of all bipartite threshold graphs, and let $\overline{T''} \subset \overline{\mathcal{B}}$ be its closure in $\overline{\mathcal{B}}$, and $T''_{\infty,\infty} := \overline{T''} \cap \mathcal{B}_{\infty\infty}$ the set of proper limits of sequences of bipartite threshold graphs; we call these bipartite threshold graph limits.







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Note that $\overline{\mathcal{T}}$, \mathcal{T}_{∞} , $\overline{\mathcal{T}''}$, $\mathcal{T}''_{\infty,\infty}$ are compact metric spaces, since they are closed subsets of $\overline{\mathcal{U}}$ or $\overline{\mathcal{B}}$.

We will give concrete representations of the threshold graph limits in Section 5. Here we give only a more abstract characterization.

Recall that t(F,G) is defined as the proportion of maps $V(F) \to V(G)$ that are graph homomorphisms. Since we are interested only in limits with $v(G) \to \infty$, it is equivalent to consider injective maps only. By inclusion–exclusion, it is further equivalent to consider $t_{\text{ind}}(F,G)$, defined as the probability that a random injective map $V(F) \to V(G)$ maps F isomorphically onto an induced copy of F in G; in other words, $t_{\text{ind}}(F,G)$ equals the number of labeled induced copies of F in G divided by the falling factorial $v(G) \cdots (v(G) - v(F) + 1)$. Then $t_{\text{ind}}(F,\cdot)$ extends by continuity to $\overline{\mathcal{U}}$, and by inclusion–exclusion, for graph limits $\Gamma \in \mathcal{U}_{\infty}$, $t_{\text{ind}}(F,\Gamma)$ can be written as a linear combination of $t(F_i,\Gamma)$ for subgraphs $F_i \subseteq F$. We can define t_{ind} for bipartite graphs in the same way; further details are in [Borgs et al. 07b] and [Diaconis and Janson 08].

Theorem 3.1.

- (i) Let $\Gamma \in \mathcal{U}_{\infty}$; i.e., Γ is a graph limit. Then $\Gamma \in \mathcal{T}_{\infty}$ if and only if $t_{\text{ind}}(P_4, \Gamma) = t_{\text{ind}}(C_4, \Gamma) = t_{\text{ind}}(2K_2, \Gamma) = 0$.
- (ii) Let $\Gamma \in \mathcal{B}_{\infty\infty}$; i.e., Γ is a bipartite graph limit. Then $\Gamma \in \mathcal{T}''_{\infty,\infty}$ if and only if $t_{\text{ind}}(2K_2,\Gamma) = 0$.

In view of (1.5) and (1.12), this is a special case of the following simple general statement.

Theorem 3.2. Let $\mathcal{F} = \{F_1, F_2, ...\}$ be a finite or infinite family of graphs, and let $\mathcal{U}_{\mathcal{F}} \subseteq \mathcal{U}$ be the set of all graphs that do not contain any graph from \mathcal{F} as an induced subgraph, i.e.,

$$\mathcal{U}_{\mathcal{F}} := \{ G \in \mathcal{U} : t_{\text{ind}}(F, G) = 0 \text{ for } F \in \mathcal{F} \}.$$

Let $\overline{\mathcal{U}}_{\mathcal{F}}$ be the closure of $\mathcal{U}_{\mathcal{F}}$ in $\overline{\mathcal{U}}$. Then

$$\overline{\mathcal{U}}_{\mathcal{F}} := \{ \Gamma \in \overline{\mathcal{U}} : t_{\text{ind}}(F, \Gamma) = 0 \text{ for } F \in \mathcal{F} \}.$$

In other words, if $\Gamma \in \mathcal{U}_{\infty}$ is a graph limit, then Γ is a limit of a sequence of graphs in $\mathcal{U}_{\mathcal{F}}$ if and only if $t_{\mathrm{ind}}(F,\Gamma) = 0$ for $F \in \mathcal{F}$.

Conversely, if $\Gamma \in \overline{\mathcal{U}}_{\mathcal{F}} \cap \mathcal{U}_{\infty}$ is represented by a function W, then the random graph G(n, W) is in $\mathcal{U}_{\mathcal{F}}$ (almost surely).

The same results hold in the bipartite case.









Proof. If $G_n \to \Gamma$ with $G \in \mathcal{U}_{\mathcal{F}}$, then $t(F,\Gamma) = \lim_{n \to \infty} t(F,G_n) = 0$ for every $F \in \mathcal{F}$, by the continuity of $t(F,\cdot)$.

Conversely, suppose that $\Gamma \in \mathcal{U}_{\infty}$ and $t(F,\Gamma) = 0$ for $F \in \mathcal{F}$, and let Γ be represented by a function W. It follows from (3.1) that if $F \in \mathcal{F}$ then $\mathbb{E} t(F,G(n,W)) = t(F,\Gamma) = 0$, and thus t(F,G(n,W)) = 0 almost surely (a.s.); consequently, $G(n,W) \in \mathcal{U}_{\mathcal{F}}$ a.s. This proves the second statement. Since $G(n,W) \to \Gamma$ a.s., it also shows that Γ is the limit of a sequence in $\mathcal{U}_{\mathcal{F}}$, and thus $\Gamma \in \overline{\mathcal{U}}_{\mathcal{F}}$, which completes the proof of the first part.

4. Degree Distributions

The results in this section hold for general graphs. They are applied to threshold graphs in Section 5.

Let \mathcal{P} be the set of probability measures on [0, 1], equipped with the standard topology of weak convergence, which makes \mathcal{P} a compact metric space (see, for example, [Billingsley 68]).

If G is a graph, let $d(v) = d_G(v)$ denote the degree of vertex $v \in V(G)$, and let D_G denote the random variable defined as the degree $d_G(v)$ of a randomly chosen vertex v (with the uniform distribution on V(G)). Thus $0 \leq D_G \leq v(G) - 1$. For a bipartite graph we similarly define $D_{G;j}$ as the degree $d_G(v)$ of a randomly chosen vertex $v \in V_j(G)$, j = 1, 2. Note that $0 \leq D_{G;1} \leq v_2(G)$ and $0 \leq D_{G;2} \leq v_1(G)$. Since we are interested in dense graphs, we will normalize these random degrees to $D_G/v(G)$ and, in the bipartite case, $D_{G;1}/v_2(G)$ and $D_{G;2}/v_1(G)$; these are random variables in [0,1].

The distribution of $D_G/v(G)$ will be called the *(normalized) degree distribution* of G and denoted by $\nu(G) \in \mathcal{P}$; in other words, $\nu(G)$ is the empirical distribution function of $\{d_G(v)/v(G) : v \in V(G)\}$. In the bipartite case we similarly have two (normalized) degree distributions: $\nu_1(G)$ for $V_1(G)$ and $\nu_2(G)$ for $V_2(G)$.

The moments of the degree distribution(s) are given by the functional $t(F, \cdot)$ for stars F, as stated in the following lemma. We omit the proof, which is a straightforward consequence of the definitions.

Lemma 4.1. The moments of $\nu(G)$ are given by

$$\int_0^1 t^k \, d\nu(G)(t) = t(K_{1,k}, G), \quad k \ge 1,$$
(4.1)

where $K_{1,k}$ is a star with k edges.







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In the bipartite case, similarly, for $k \geq 1$,

$$\int_0^1 t^k \, \mathrm{d}\nu_1(G)(t) = t(K_{1,k}, G), \quad \int_0^1 t^k \, \mathrm{d}\nu_2(G)(t) = t(K_{k,1}, G). \tag{4.2}$$

This enables us to extend the definition of the (normalized) degree distribution to the limit objects by continuity.

Theorem 4.2. If G_n are graphs with $v(G_n) \to \infty$ and $G_n \to \Gamma$ for some $\Gamma \in \overline{\mathcal{U}}$ as $n \to \infty$, then $\nu(G_n) \to \nu(\Gamma)$ for some distribution $\nu(\Gamma) \in \mathcal{P}$. This defines the "degree distribution" $\nu(\Gamma)$ (uniquely) for every graph limit $\Gamma \in \mathcal{U}_{\infty}$, and $\Gamma \mapsto \nu(\Gamma)$ is a continuous map $\mathcal{U}_{\infty} \to \mathcal{P}$. Furthermore, (4.1) holds for all $G \in \overline{\mathcal{U}}$.

Similarly, in the bipartite case, ν_1 and ν_2 extend to continuous maps $\overline{\mathcal{B}} \to \mathcal{P}$ such that (4.2) holds for all $G \in \overline{\mathcal{B}}$. Furthermore, $\nu_2(\Gamma) = \nu_1(\Gamma^{\dagger})$ for $\Gamma \in \overline{\mathcal{B}}$.

Proof. The result is an immediate consequence of Lemma 4.1 and the method of moments. The last sentence follows from (4.2) and (3.3).

Remark 4.3. Theorem 4.2 says that the degree distribution ν is a *testable* graph parameter in the sense of [Borgs et al. 07b]; see in particular [Borgs et al. 07b, Section 6] (except that ν takes values in \mathcal{P} instead of \mathbb{R}).

If Γ is represented by a function W on $[0,1]^2$, we can easily find its degree distribution from W.

Theorem 4.4. If $W \in \mathcal{W}_s$, then $\nu(\Gamma_W)$ equals the distribution of $\int_0^1 W(U,y) \, \mathrm{d}y$, where $U \sim U(0,1)$.

Similarly, in the bipartite case, if $W \in \mathcal{W}$, then $\nu_1(\Gamma_W'')$ equals the distribution of $\int_0^1 W(U,y) \, dy$, and $\nu_2(\Gamma_W'')$ equals the distribution of $\int_0^1 W(x,U) \, dx$.

Proof. By (4.1) and (3.1),

$$\int_{0}^{1} t^{k} d\nu(\Gamma_{W})(t) = t(K_{1,k}, \Gamma_{W}) = \int_{[0,1]} \left(\int_{[0,1]} W(x,y) dy \right)^{k} dx$$
$$= \mathbb{E} \left(\int_{[0,1]} W(U,y) dy \right)^{k}$$

for every $k \ge 1$, and the result follows. The bipartite case is similar, using (3.2).









If a graph G has n vertices, its number of edges is

$$|E(G)| = \frac{1}{2} \sum_{v \in V(G)} d(v) = \frac{n}{2} \mathbb{E} D_G = \frac{n^2}{2} \mathbb{E} (D_G/n) = \frac{n^2}{2} \int_0^1 t \, d\nu(G)(t).$$

Hence the edge density of G is

$$\frac{|E(G)|}{\binom{n}{2}} = \frac{n}{n-1} \int_0^1 t \, d\nu(G)(t). \tag{4.3}$$

If (G_n) is a sequence of graphs with $v(G_n) \to \infty$ and $G_n \to \Gamma \in \mathcal{U}_{\infty}$, we see from (4.3) and Theorem 4.2 that the graph densities converge to $\int_0^1 t \, \mathrm{d}\nu(\Gamma)(t)$, the mean of the distribution $\nu(\Gamma)$, which thus may be called the (edge) density of $\Gamma \in \mathcal{U}_{\infty}$.

If Γ is represented by a function W on $[0,1]^2$, Theorem 4.4 yields the following.

Corollary 4.5. The graph Γ_W has edge density $\iint_{[0,1]^2} W(x,y) \, dx \, dy$ for every $W \in \mathcal{W}_s$.

Proof. By Theorem 4.4, the mean of $\mu(\Gamma_W)$ equals

$$\mathbb{E} \int_0^1 W(U, y) \, \mathrm{d}y = \int_0^1 \int_0^1 W(x, y) \, \mathrm{d}x \, \mathrm{d}y.$$

5. Limits of Threshold Graphs

Recall from Section 3.5 that \mathcal{T}_{∞} is the set of limits of threshold graphs, and $\mathcal{T}''_{\infty,\infty}$, is the set of limits of bipartite threshold graphs. Our purpose in this section is to characterize the threshold graph limits, i.e., the elements of \mathcal{T}_{∞} and $\mathcal{T}''_{\infty,\infty}$, and give simple criteria for the convergence of a sequence of threshold graphs to one of these limits. We begin with some definitions.

A function $W: [0,1]^2 \to \mathbb{R}$ is increasing if $W(x,y) \leq W(x',y')$ whenever $0 \leq x \leq x' \leq 1$ and $0 \leq y \leq y' \leq 1$. A set $S \subseteq [0,1]^2$ is increasing if its indicator $\mathbf{1}_S$ is an increasing function on $[0,1]^2$, i.e., if $(x,y) \in S$ implies $(x',y') \in S$ whenever $0 \leq x \leq x' \leq 1$ and $0 \leq y \leq y' \leq 1$.

If $\mu \in \mathcal{P}$, let F_{μ} be its distribution function $F_{\mu}(x) := \mu([0,x])$, and let $F_{\mu}(x-) := \mu([0,x])$ be its left-continuous version. Thus $F_{\mu}(0-) = 0 \le F_{\mu}(0)$ and $F_{\mu}(1-) \le 1 = F_{\mu}(1)$. Further, let $F_{\mu}^{-1} : [0,1] \to [0,1]$ be the right-continuous inverse defined by

$$F_{\mu}^{-1}(x) := \sup\{ t \le 1 : F_{\mu}(t) \le x \}. \tag{5.1}$$







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Note that $F_{\mu}^{-1}(0) \geq 0$ and $F_{\mu}^{-1}(1) = 1$. Finally, define

$$S_{\mu} := \left\{ (x, y) \in [0, 1]^2 : x \ge F_{\mu} ((1 - y) -) \right\}. \tag{5.2}$$

It is easily seen that S_{μ} is a closed increasing subset of $[0,1]^2$ and that it contains the upper and right edges $\{(x,1)\}$ and $\{(1,y)\}$. Since $x \geq F_{\mu}((1-y)-) \iff F_{\mu}^{-1}(x) \geq 1-y$, we also have

$$S_{\mu} = \left\{ (x, y) \in [0, 1]^2 : F_{\mu}^{-1}(x) + y \ge 1 \right\}. \tag{5.3}$$

We further write $W_{\mu} := \mathbf{1}_{S_{\mu}}$ and let $\Gamma''_{\mu} := \Gamma''(W_{\mu})$, and when W is symmetric, $\Gamma_{\mu} := \Gamma(W_{\mu})$. We denote the interior of a set S by S° . It is easily verified from (5.2) that

$$S_{\mu}^{\circ} = \{ (x, y) \in (0, 1)^{2} : x > F_{\mu}(1 - y) \}.$$
 (5.4)

Recall that the *Hausdorff distance* between two nonempty compact subsets K_1 and K_2 of some metric space S is defined by

$$d_H(K_1, K_2) := \max \left(\max_{x \in K_1} d(x, K_2), \max_{y \in K_2} d(y, K_1) \right). \tag{5.5}$$

This defines a metric on the set of all nonempty compact subsets of S. If S is compact, the resulting topology on the set of compact subsets of S (with the empty set as an isolated point) is compact and equals the *Fell topology* (see, for example, [Kallenberg 02, Appendix A.2]) on the set of all closed subsets of S.

Let λ_d denote the Lebesgue measure in \mathbb{R}^d . For measurable subsets S_1, S_2 of $[0,1]^2$, we also consider their measure distance $\lambda_2(S_1\Delta S_2)$. This equals the L^1 -distance of their indicator functions, and is thus a metric modulo null sets.

For functions in W we also use two different metrics: the L^1 -distance

$$\int_{[0,1]^2} |W_1(x,y) - W_2(x,y)| \, \mathrm{d}x \, \mathrm{d}y$$

and, in the symmetric case, the *cut distance* δ_{\square} defined in [Borgs et al. 07b], and in the bipartite case its analogue δ''_{\square} ; see Section 3. Note that the cut distance is only a pseudometric, since the distance of two different functions may be 0. Note further that the cut distance is less than or equal to the L^1 -distance.

We can now prove one of our main results, giving several related characterizations of threshold graph limits. There are two versions, since we treat the bipartite case in parallel.









5.1. The Bipartite Case

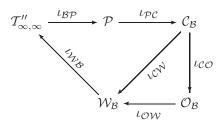
It is convenient to begin with the bipartite case.

Theorem 5.1. There are bijections between the set $\mathcal{T}''_{\infty,\infty}$ of graph limits of bipartite threshold graphs and each of the following sets:

- (i) The set \mathcal{P} of probability distributions on [0,1].
- (ii) The set C_B of increasing closed sets $S \subseteq [0,1]^2$ that contain the upper and right edges $[0,1] \times \{1\} \cup \{1\} \times [0,1]$.
- (iii) The set $\mathcal{O}_{\mathcal{B}}$ of increasing open sets $S \subseteq (0,1)^2$.
- (iv) The set W_B of increasing 0-1-valued functions $W : [0,1]^2 \to \{0,1\}$ modulo a.e. equality.

More precisely, there are commuting bijections between these sets given by the following mappings and their compositions:

$$\iota_{\mathcal{BP}} : \mathcal{T}''_{\infty,\infty} \to \mathcal{P}, \qquad \iota_{\mathcal{BP}}(\Gamma) := \nu_{1}(\Gamma);
\iota_{\mathcal{PC}} : \mathcal{P} \to \mathcal{C}_{\mathcal{B}}, \qquad \iota_{\mathcal{PC}}(\mu) := S_{\mu};
\iota_{\mathcal{CO}} : \mathcal{C}_{\mathcal{B}} \to \mathcal{O}_{\mathcal{B}}, \qquad \iota_{\mathcal{CO}}(S) := S^{\circ};
\iota_{\mathcal{CW}} : \mathcal{C}_{\mathcal{B}} \to \mathcal{W}_{\mathcal{B}}, \qquad \iota_{\mathcal{CW}}(S) := \mathbf{1}_{S};
\iota_{\mathcal{OW}} : \mathcal{O}_{\mathcal{B}} \to \mathcal{W}_{\mathcal{B}}, \qquad \iota_{\mathcal{OW}}(S) := \mathbf{1}_{S};
\iota_{\mathcal{WB}} : \mathcal{W}_{\mathcal{B}} \to \mathcal{T}''_{\infty,\infty}, \qquad \iota_{\mathcal{WB}}(W) := \Gamma''_{W};$$
(5.6)



In particular, a probability distribution $\mu \in \mathcal{P}$ corresponds to $\Gamma''_{\mu} \in \mathcal{T}''_{\infty,\infty}$ and to $S_{\mu} \in \mathcal{C}_{\mathcal{B}}$, $S_{\mu}^{\circ} \in \mathcal{O}_{\mathcal{B}}$, and $W_{\mu} \in \mathcal{W}_{\mathcal{B}}$. Conversely, $\Gamma \in \mathcal{T}''_{\infty,\infty}$ corresponds to $\nu_1(\Gamma) \in \mathcal{P}$. Thus, the mappings $\Gamma \mapsto \nu_1(\Gamma)$ and $\mu \mapsto \Gamma''_{\mu}$ are the inverses of each other.

Moreover, these bijections are homeomorphisms, with any of the following topologies or metrics: the standard (weak) topology on \mathcal{P} ; the Hausdorff metric or the Fell topology or the measure distance on $\mathcal{C}_{\mathcal{B}}$; the measure distance on $\mathcal{O}_{\mathcal{B}}$; the L^1 -distance or the cut distance on the set $\mathcal{W}_{\mathcal{B}}$.







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Proof. The mappings in (5.6) are all well defined, except that we do not yet know that ι_{WB} maps $W_{\mathcal{B}}$ into $T''_{\infty,\infty}$. We thus regard ι_{WB} as a map $W_{\mathcal{B}} \to \mathcal{B}_{\infty\infty}$ and let $\widetilde{\mathcal{B}} := \iota_{WB}(W_{\mathcal{B}})$ be its image; we will identify this as $T''_{\infty,\infty}$ later. For the time being, we also regard $\iota_{\mathcal{BP}}$ as defined on $\widetilde{\mathcal{B}}$ (or on all of $\mathcal{B}_{\infty\infty}$).

Consider first $\iota_{\mathcal{PC}}: \mathcal{P} \to \mathcal{C}_{\mathcal{B}}$. By (5.2), S_{μ} determines F_{μ} at all continuity points, and thus it determines μ . Consequently, $\iota_{\mathcal{PC}}$ is injective.

If $S \in \mathcal{C}_{\mathcal{B}}$ and $y \in [0,1]$, then $\{x : (x,y) \in S\}$ is a closed subinterval of [0,1] that contains 1, and thus $S = \{(x,y) \in [0,1]^2 : x \geq g(y)\}$ for some function $g : [0,1] \to [0,1]$. Moreover, g(1) = 0, g is decreasing, i.e., $g(y_2) \leq g(y_1)$ if $y_1 \leq y_2$, and since S is closed, g is right-continuous. Thus g(1-x) is increasing and left-continuous, and hence there exists a probability measure $\mu \in \mathcal{P}$ such that $F_{\mu}(x-) = g(1-x)$, $x \in [0,1]$. By (5.2), then

$$\iota_{\mathcal{PC}}(\mu) = S_{\mu} = \{ (x, y) \in [0, 1]^2 : x \ge g(y) \} = S.$$

Hence $\iota_{\mathcal{PC}}$ is onto. Consequently, $\iota_{\mathcal{PC}}$ is a bijection of \mathcal{P} onto $\mathcal{C}_{\mathcal{B}}$.

If S_1 and S_2 are two different sets in $\mathcal{C}_{\mathcal{B}}$, then there exists a point $(x,y) \in S_1 \setminus S_2$, say. There is a small open disk with center in (x,y) that does not intersect S_2 , and since S_1 is increasing, at least a quarter of the disk is contained in $S_1 \setminus S_2$. Hence $\lambda_2(S_1 \Delta S_2) > 0$. Similarly, if S_1 and S_2 are two different sets in $\mathcal{O}_{\mathcal{B}}$ and $(x,y) \in S_1 \setminus S_2$, then there is a small open disk with center in (x,y) that is contained in S_1 , and since S_2 is increasing, at least a quarter of the disk is contained in $S_1 \setminus S_2$, whence $\lambda_2(S_1 \Delta S_2) > 0$. This shows that the measure distance is a metric on $\mathcal{C}_{\mathcal{B}}$ and on $\mathcal{O}_{\mathcal{B}}$, and that the mappings $\iota_{\mathcal{CW}}$ and $\iota_{\mathcal{OW}}$ into $\mathcal{W}_{\mathcal{B}}$ are injective (recall that a.e. equal functions are identified in $\mathcal{W}_{\mathcal{B}}$).

Next, let $S \subseteq [0,1]^2$ be increasing. If $(x,y) \in \overline{S}$ with x < 1 and y < 1, it is easily seen that $(x,x+\delta) \times (y,y+\delta) \subseteq S$ for $\delta = \min\{1-x,1-y\}$, and thus $(x,x+\delta) \times (y,y+\delta) \subseteq S^{\circ}$. It follows that for any real a, the intersection of the boundary $\partial S := \overline{S} \setminus S^{\circ}$ with the diagonal line $L_a := \{(x,x+a) : x \in \mathbb{R}\}$ consists of at most two points (of which one is on the boundary of $[0,1]^2$). In particular, $\lambda_1(\partial S \cap L_a) = 0$ and thus

$$\lambda_2(\partial S) = 2^{-1/2} \int_{-1}^1 \lambda_1(\partial S \cap L_a) \, \mathrm{d}a = 0.$$
 (5.7)

Consequently, ∂S is a null set for every increasing S. Among other things, this shows that if $S \in \mathcal{C}_{\mathcal{B}}$, then $\iota_{\mathcal{OW}}\iota_{\mathcal{CO}}(S) = \mathbf{1}_{S^{\circ}} = \mathbf{1}_{S}$ a.e. Since elements of $\mathcal{W}_{\mathcal{B}}$ are defined modulo a.e. equality, this shows that $\iota_{\mathcal{OW}}\iota_{\mathcal{CO}} = \iota_{\mathcal{CW}} : \mathcal{C}_{\mathcal{B}} \to \mathcal{W}_{\mathcal{B}}$.

If $W \in \mathcal{W}_{\mathcal{B}}$, and thus $W = \mathbf{1}_S$ for some increasing $S \subseteq [0,1]^2$, let

$$\widetilde{S} := \overline{S \cup [0,1] \times \{1\} \cup \{1\} \times [0,1]}.$$
 (5.8)









Then $\widetilde{S} \in \mathcal{C}_{\mathcal{B}}$ and (5.7) imply that $\iota_{\mathcal{CW}}(\widetilde{S}) = \mathbf{1}_{\widetilde{S}} = \mathbf{1}_{S} = W$ a.e. Similarly, $S^{\circ} \in \mathcal{O}_{\mathcal{B}}$ and $\iota_{\mathcal{OW}}(S^{\circ}) = \mathbf{1}_{S} = W$ a.e. Consequently, $\iota_{\mathcal{CW}}$ and $\iota_{\mathcal{OW}}$ are onto, and thus bijections. Similarly (or as a consequence), $\iota_{\mathcal{CO}}$ is a bijection of $\mathcal{C}_{\mathcal{B}}$ onto $\mathcal{O}_{\mathcal{B}}$, with inverse $S \mapsto \widetilde{S}$ given by (5.8).

Note that the composition $\iota_{\mathcal{CW}}\iota_{\mathcal{PC}}$ maps $\mu \mapsto \mathbf{1}_{S_{\mu}} = W_{\mu}$, and let $\iota_{\mathcal{PB}}$ be the composition $\iota_{\mathcal{WB}}\iota_{\mathcal{CW}}\iota_{\mathcal{PC}} : \mu \mapsto \Gamma''(W_{\mu}) = \Gamma''_{\mu}$ mapping \mathcal{P} into $\mathcal{B}_{\infty\infty}$. Since $\iota_{\mathcal{PC}}$ and $\iota_{\mathcal{CW}}$ are bijections, its image satisfies $\iota_{\mathcal{BP}}(\mathcal{P}) = \iota_{\mathcal{WB}}(\mathcal{W}_{\mathcal{B}}) = \widetilde{\mathcal{B}} \subseteq \mathcal{B}_{\infty\infty}$.

If $\mu \in \mathcal{P}$, then the composition $\iota_{\mathcal{BP}}\iota_{\mathcal{PB}}(\mu) = \nu_1(\Gamma''_{\mu})$ equals by Theorem 4.4 and (5.3) the distribution of

$$\int_0^1 \mathbf{1}_{S_{\mu}}(U, y) \, \mathrm{d}y = F_{\mu}^{-1}(U). \tag{5.9}$$

As is well known, and easy to verify using (5.1), this distribution equals μ . Hence, the composition $\iota_{\mathcal{BP}}\iota_{\mathcal{PB}}$ is the identity. It follows that $\iota_{\mathcal{PB}}$ is injective and thus a bijection of \mathcal{P} onto its image $\widetilde{\mathcal{B}}$, and that $\iota_{\mathcal{BP}}$ (restricted to $\widetilde{\mathcal{B}}$) is its inverse.

We have shown that all mappings in (5.6) are bijections, except that we have not yet shown that $\widetilde{\mathcal{B}} = \mathcal{T}''_{\infty,\infty}$. We next show that the mappings are homeomorphisms.

Recall that the topology on \mathcal{P} can be defined by the Lévy metric defined by (see, e.g., [Gut 05, Problem 5.25])

$$d_L(\mu_1, \mu_2) := \inf\{\varepsilon > 0 : F_{\mu_1}(x - \varepsilon) - \varepsilon \le F_{\mu_2}(x) \le F_{\mu_1}(x + \varepsilon) + \varepsilon \text{ for all } x\}.$$
(5.10)

If $\mu_1, \mu_2 \in \mathcal{P}$ with $d_L(\mu_1, \mu_2) < \varepsilon$, it follows from (5.2) and (5.10) that if $(x, y) \in S_{\mu_1}$ and $x, y < 1 - \varepsilon$, then

$$F_{\mu_2}((1-y-\varepsilon)-) \le F_{\mu_1}((1-y)-) + \varepsilon \le x + \varepsilon,$$

and thus $(x + \varepsilon, y + \varepsilon) \in S_{\mu_2}$. Considering also the simple cases $x \in [1 - \varepsilon, 1]$ and $y \in [1 - \varepsilon, 1]$, it follows that if $(x, y) \in S_{\mu_1}$, then $d((x, y), S_{\mu_2}) \leq \sqrt{2} \varepsilon$. Consequently, by (5.5) and symmetry,

$$d_H(S_{\mu_1}, S_{\mu_2}) \le \sqrt{2} d_L(\mu_1, \mu_2),$$

which shows that $\iota_{\mathcal{PC}}$ is continuous if $\mathcal{C}_{\mathcal{B}}$ is given the topology given by the Hausdorff metric.

The same argument shows that for any (x_0, y_0) , the intersection of the difference $S_{\mu_1} \Delta S_{\mu_2}$ with the diagonal line L_a defined above is an interval of length at most $\sqrt{2} d_L(\mu_1, \mu_2)$, and thus by integration over a as in (5.7),

$$\lambda_2 \left(S_{\mu_1} \Delta S_{\mu_2} \right) \le 2d_L(\mu_1, \mu_2).$$









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Hence, $\iota_{\mathcal{PC}}$ is continuous also if $\mathcal{C}_{\mathcal{B}}$ is given the topology given by the measure distance.

Since \mathcal{P} is compact and $\iota_{\mathcal{PC}}$ is a bijection, it follows that $\iota_{\mathcal{PC}}$ is a homeomorphism for both these topologies on $\mathcal{C}_{\mathcal{B}}$. In particular, these topologies coincide on $\mathcal{C}_{\mathcal{B}}$. As remarked before the theorem, since $[0,1]^2$ is compact, also the Fell topologies coincide with these on $\mathcal{C}_{\mathcal{B}}$.

The bijections $\iota_{\mathcal{CO}}$, $\iota_{\mathcal{CW}}$, and $\iota_{\mathcal{OW}}$ are isometries for the measure distance on $\mathcal{C}_{\mathcal{B}}$ and $\mathcal{O}_{\mathcal{B}}$ and the L^1 -distance on $\mathcal{W}_{\mathcal{B}}$, and thus homeomorphisms. Furthermore, still using the L^1 -distance on $\mathcal{W}_{\mathcal{B}}$, it is easily seen from (3.2), as for the symmetric case in [Lovász and Szegedy 06, Borgs et al. 07b], that for every fixed bipartite graph F, the mapping $W \mapsto t(F, \Gamma_W'')$ is continuous, which by definition of the topology in $\mathcal{B}_{\infty\infty}$ means that $\iota_{\mathcal{W}\mathcal{B}}: W \mapsto \Gamma_W''$ is continuous. Hence, the bijection $\iota_{\mathcal{W}\mathcal{B}}$ is a homeomorphism of the compact space $\mathcal{W}_{\mathcal{B}}$ onto its image $\widetilde{\mathcal{B}}$.

As stated above, the cut distance is only a pseudometric on W. But two functions in W with cut distance 0 are mapped onto the same element in $\mathcal{B}_{\infty\infty}$, and since we have shown that $\iota_{W\mathcal{B}}$ is injective on $W_{\mathcal{B}}$, it follows that the restriction of the cut distance to $W_{\mathcal{B}}$ is a metric. Moreover, the identity map on $W_{\mathcal{B}}$ is continuous from the L^1 -metric to the cut metric, and since the space is compact under the former metric, the two metrics are equivalent on $W_{\mathcal{B}}$.

We have shown that all mappings are homeomorphisms. It remains only to show that $\widetilde{\mathcal{B}} = \mathcal{T}''_{\infty,\infty}$. To do this, observe first that if G is a bipartite threshold graph, and we order its vertices in each of the two vertex sets with increasing vertex degrees, then the function W_G defined in Section 3 is increasing and belongs thus to $\mathcal{W}_{\mathcal{B}}$. Consequently, $\pi(G) = \iota_{\mathcal{W}\mathcal{B}}(W_G) \in \widetilde{\mathcal{B}}$. If $\Gamma \in \mathcal{T}''_{\infty,\infty}$, then by definition there exists a sequence G_n of bipartite threshold graphs with $v_1(G_n), v_2(G_n) \to \infty$ such that $G_n \to \Gamma$ in $\overline{\mathcal{B}}$. This implies that $\pi(G_n) \to \Gamma$ in $\mathcal{B}_{\infty\infty}$, and since $\pi(G_n) \in \widetilde{\mathcal{B}}$ and $\widetilde{\mathcal{B}}$ is compact and thus a closed subset of $\mathcal{B}_{\infty\infty}$, we find that $\Gamma \in \widetilde{\mathcal{B}}$.

Conversely, if $\Gamma \in \widetilde{\mathcal{B}}$, then $\Gamma = \iota_{\mathcal{WB}}\iota_{\mathcal{CW}}(S)$ for some set $S \in \mathcal{C}_{\mathcal{B}}$. For each n, partition $[0,1]^2$ into n^2 closed squares Q_{ij} of side 1/n, and let S_n be the union of all Q_{ij} that intersect S. Then $S_n \in \mathcal{C}_{\mathcal{B}}$, $S \subseteq S_n$, and $d_H(S_n, S) \leq \sqrt{2}/n$. Let $W_n := \mathbf{1}_{S_n} = \iota_{\mathcal{CW}}(S_n)$ and let $\Gamma_n := \iota_{\mathcal{WB}}(W_n) \in \widetilde{\mathcal{B}}$. Since $\iota_{\mathcal{CW}}$ and $\iota_{\mathcal{WB}}$ are continuous, $W_n \to W := \mathbf{1}_S$ in $\mathcal{W}_{\mathcal{B}}$ and $\Gamma_n \to \iota_{\mathcal{WB}}(W) = \Gamma$ in $\widetilde{\mathcal{B}} \subset \mathcal{B}_{\infty\infty}$. However, W_n is a step function of the form $W(G_n)$ for some bipartite graph G_n with $v_1(G_n) = v_2(G_n) = n$, and thus $\pi(G_n) = \Gamma''_{W_n} = \Gamma_n$. Moreover, each S_n , and thus each W_n , is increasing, and hence G_n is a bipartite threshold graph. Since $\pi(G_n) = \Gamma_n \to \Gamma$ in $\mathcal{B}_{\infty\infty}$, it follows that $G_n \to \Gamma$ in $\overline{\mathcal{B}}$, and thus $\Gamma \in \mathcal{T}''_{\infty,\infty}$.

Consequently, $\widetilde{\mathcal{B}} = \mathcal{T}''_{\infty,\infty}$, which completes the proof.









Remark 5.2. Another unique representation by increasing closed sets is given by the family $\mathcal{C}'_{\mathcal{B}}$ of closed increasing subsets S of $[0,1]^2$ that satisfy $S = \overline{S^{\circ}}$; there are bijections $\mathcal{C}'_{\mathcal{B}} \to \mathcal{O}_{\mathcal{B}}$ and $\mathcal{O}_{\mathcal{B}} \to \mathcal{C}'_{\mathcal{B}}$ given by $S \mapsto S^{\circ}$ and $S \mapsto \overline{S}$. We can again use the measure distance on $\mathcal{C}'_{\mathcal{B}}$, but not the Hausdorff distance. (For example, $[0,1] \times [1-\varepsilon,1] \to \emptyset$ in $\mathcal{C}'_{\mathcal{B}}$ as $\varepsilon \to 0$.)

Corollary 5.3. The degree distribution yields a homeomorphism $\Gamma \mapsto \nu_1(\Gamma)$ of $\mathcal{T}''_{\infty,\infty}$ onto \mathcal{P} .

Of course, $\Gamma \mapsto \nu_2(\Gamma) = \nu_1(\Gamma^{\dagger})$ yields another homeomorphism of $\mathcal{T}''_{\infty,\infty}$ onto \mathcal{P} . To see the connection between these, and (more importantly) to prepare for the corresponding result in the nonbipartite case, we investigate further the reflection involution.

If $S \subseteq [0,1]^2$, let $S^{\dagger} := \{(x,y) : (y,x) \in S\}$ be the set S reflected in the main diagonal. Thus $\mathbf{1}_{S^{\dagger}} = \mathbf{1}_{S}^{\dagger}$. We have defined the reflection map \dagger for bipartite graphs and graph limits, and for the sets and functions in Theorem 5.1(ii)(iii)(iv), and it is easily seen that these correspond to each other by the bijections in Theorem 5.1. Consequently, there is a corresponding map (involution) $\mu \mapsto \mu^{\dagger}$ of \mathcal{P} onto itself too. This map is less intuitive than the others; to find it explicitly, we find from (5.2), (5.3), and $S_{\mu^{\dagger}} = S_{\mu}^{\dagger}$ that

$$x \ge F_{\mu^{\dagger}}((1-y)-) \iff (y,x) \in S_{\mu} \iff F_{\mu}^{-1}(y)+x \ge 1$$

and thus $F_{\mu^{\dagger}}((1-y)-) = 1 - F_{\mu}^{-1}(y)$ and

$$F_{\mu^{\dagger}}(t) = 1 - F_{\mu}^{-1}((1-t)), \quad 0 \le t \le 1.$$
 (5.11)

This means that the graph of the distribution function is reflected about the diagonal between (0,1) and (1,0) (and adjusted at the jumps).

The map \dagger is continuous on \mathcal{P} , by Theorem 5.1 and the obvious fact that $S \mapsto S^{\dagger}$ is continuous on, for example, $\mathcal{C}_{\mathcal{B}}$.

We let $\mathcal{P}_s := \{ \mu \in \mathcal{P} : \mu^{\dagger} = \mu \} = \{ \mu \in \mathcal{P} : S_{\mu} = S_{\mu}^{\dagger} \}$ be the set of probability distributions invariant under the involution \dagger . Since \dagger is continuous, \mathcal{P}_s is a closed and thus compact subset of \mathcal{P} .

Remark 5.4. If $\mu \in \mathcal{P}_s$, let $x_0 := 1 - \inf\{x : (x, x) \in S_\mu\}$. Then (5.2) and (5.4) imply that $F_\mu(x_0-) \le 1 - x_0 \le F_\mu(x_0)$, and the restriction of F_μ to $[0, x_0)$ is an increasing right-continuous function with values in $[0, 1-x_0]$, and this restriction determines $F_\mu(t)$ for $x \ge x_0$ too by (5.11).

Conversely, given any $x_0 \in [0, 1]$ and increasing right-continuous $F : [0, x_0) \to [0, 1 - x_0]$, there is a unique $\mu \in \mathcal{P}_s$ with $F_{\mu}(x) = F(x)$ for $x < x_0$ and $F_{\mu}(x_0) \ge 1 - x_0$.









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5.2. Nonbipartite Case

We can now state our main theorem for (nonbipartite) threshold graph limits.

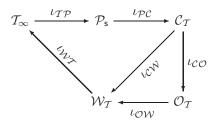
Theorem 5.5. There are bijections between the set \mathcal{T}_{∞} of graph limits of threshold graphs and each of the following sets:

- (i) The set \mathcal{P}_s of probability distributions on [0,1] symmetric with respect to \dagger .
- (ii) The set C_T of symmetric increasing closed sets $S \subseteq [0,1]^2$ that contain the upper and right edges $[0,1] \times \{1\} \cup \{1\} \times [0,1]$.
- (iii) The set $\mathcal{O}_{\mathcal{T}}$ of symmetric increasing open sets $S \subseteq (0,1)^2$.
- (iv) The set W_T of symmetric increasing 0-1-valued functions $W:[0,1]^2 \to \{0,1\}$ modulo a.e. equality.

More precisely, there are commuting bijections between these sets given by the following mappings and their compositions:

$$\iota_{\mathcal{T}\mathcal{P}}: \mathcal{T}_{\infty} \to \mathcal{P}_{s}, \qquad \iota_{\mathcal{T}\mathcal{P}}(\Gamma) := \nu(\Gamma);
\iota_{\mathcal{P}\mathcal{C}}: \mathcal{P}_{s} \to \mathcal{C}_{\mathcal{T}}, \qquad \iota_{\mathcal{P}\mathcal{C}}(\mu) := S_{\mu};
\iota_{\mathcal{C}\mathcal{O}}: \mathcal{C}_{\mathcal{T}} \to \mathcal{O}_{\mathcal{T}}, \qquad \iota_{\mathcal{C}\mathcal{O}}(S) := S^{\circ};
\iota_{\mathcal{C}\mathcal{W}}: \mathcal{C}_{\mathcal{T}} \to \mathcal{W}_{\mathcal{T}}, \qquad \iota_{\mathcal{C}\mathcal{W}}(S) := \mathbf{1}_{S};
\iota_{\mathcal{O}\mathcal{W}}: \mathcal{O}_{\mathcal{T}} \to \mathcal{W}_{\mathcal{T}}, \qquad \iota_{\mathcal{O}\mathcal{W}}(S) := \mathbf{1}_{S};
\iota_{\mathcal{W}\mathcal{T}}: \mathcal{W}_{\mathcal{T}} \to \mathcal{T}_{\infty}, \qquad \iota_{\mathcal{W}\mathcal{T}}(W) := \Gamma_{W}.$$
(5.12)

In particular, a probability distribution $\mu \in \mathcal{P}_s$ corresponds to $\Gamma_{\mu} \in \mathcal{T}_{\infty}$ and to $S_{\mu} \in \mathcal{C}_{\mathcal{T}}$, $S_{\mu}^{\circ} \in \mathcal{O}_{\mathcal{T}}$, and $W_{\mu} \in \mathcal{W}_{\mathcal{T}}$. Conversely, $\Gamma \in \mathcal{T}_{\infty}$ corresponds to $\nu(\Gamma) \in \mathcal{P}_s$. Thus, the mappings $\Gamma \mapsto \nu(\Gamma)$ and $\mu \mapsto \Gamma_{\mu}$ are the inverses of each other.



Moreover, these bijections are homeomorphisms, with any of the following topologies or metrics: the standard (weak) topology on $\mathcal{P}_s \subset \mathcal{P}$; the Hausdorff









metric or the Fell topology or the measure distance on C_T ; the measure distance on O_T ; the L^1 -distance or the cut distance on the set W_T . These homeomorphic topological spaces are compact metric spaces.

Proof. The mappings $\iota_{\mathcal{PC}}$, $\iota_{\mathcal{CO}}$, $\iota_{\mathcal{CW}}$, $\iota_{\mathcal{CW}}$ are restrictions of the corresponding mappings in Theorem 5.1, and it follows from Theorem 5.1 and the definitions that these mappings are bijections and homeomorphisms for the given topologies. The spaces are closed subspaces of the corresponding spaces in Theorem 5.1, since \dagger is continuous on these spaces, and are thus compact metric spaces.

The rest is as in the proof of Theorem 5.1, and we omit some details. It follows from Theorem 4.4 that the composition $\iota_{\mathcal{WB}}\iota_{\mathcal{CW}}\iota_{\mathcal{PC}}: \mu \mapsto \Gamma(W_{\mu}) = \Gamma_{\mu}$ is a bijection of \mathcal{P}_s onto a subset T' of $T''_{\infty,\infty}$, with $\iota_{\mathcal{TP}}$ as its inverse. It follows that these mappings too are homeomorphisms, and that the L^1 -distance and cut distance are equivalent on $\mathcal{W}_{\mathcal{T}}$.

To see that $T' = T_{\infty}$, we also follow the proof of Theorem 5.1. A minor complication is that if $G \in \mathcal{T}$ is a threshold graph, and we order the vertices with increasing degrees, then W_G is not increasing, because $W_G(x,x) = 0$ for all x, since we consider loopless graphs only. However, we can define $W^*(G)$ by changing W_G to be 1 on some squares on the diagonal so that $W^*(G)$ is increasing and thus $W^*(G) \in \mathcal{W}_{\mathcal{T}}$, and the error $\|W_G - W^*(G)\|_{L^1}$ is at most 1/v(G). If we define $\pi^*(G) := \Gamma(W^*(G)) \in \mathcal{T}'$, we see that if (G_n) is a sequence of threshold graphs with $v(G_n) \to \infty$, then for every graph F, by a simple estimate, see, e.g., [Lovász and Szegedy 06, Lemma 4.1], we have

$$|t(F, \pi^*(G_n)) - t(F, G_n)| \le e(F) ||W(G_n) - W^*(G_n)||_{L^1} \le e(F)/v(G_n) \to 0.$$
(5.13)

It follows that $G_n \to \Gamma$ in $\overline{\mathcal{U}}$ if and only if $\pi^*(G_n) \to \Gamma$ in \mathcal{U}_{∞} . If $\Gamma \in \mathcal{T}_{\infty}$, then there exists such a sequence $G_n \to \Gamma$, and thus $\pi^*(G_n) \to \Gamma$ in \mathcal{U}_{∞} , and since $\pi^*(G_n) \in \mathcal{T}'$ and \mathcal{T}' is compact, we find that $\Gamma \in \mathcal{T}'$.

The converse follows in the same way. If $\Gamma \in \mathcal{T}'$, then $\Gamma = \iota_{\mathcal{WT}}(W)$ for some function $W \in \mathcal{W}_{\mathcal{T}}$. The approximating step functions W_n constructed in the proof of Theorem 5.1 are symmetric, and if we let W_n^* be the modification that vanishes on all diagonal squares, then $W_n^* = W_{G_n}$ for some threshold graph G_n , and for every graph F,

$$t(F, G_n) = t(F, W_n^*) = t(F, W_n) + o(1) = t(F, W) + o(1).$$

Hence, $G_n \to \Gamma_W = \Gamma$ in $\overline{\mathcal{U}}$, and thus $\Gamma \in \mathcal{T}_{\infty}$. Consequently, $\mathcal{T}' = \mathcal{T}_{\infty}$.

Corollary 5.6. The degree distribution yields a homeomorphism $\Gamma \mapsto \nu(\Gamma)$ of \mathcal{T}_{∞} onto the closed subspace \mathcal{P}_s of \mathcal{P} .









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Remark 5.7. The fact that a graph limit Γ can be represented by a function $W \in \mathcal{W}_{\mathcal{T}}$ if and only if $t_{\text{ind}}(P_4, \Gamma) = t_{\text{ind}}(C_4, \Gamma) = t_{\text{ind}}(2K_2, \Gamma) = 0$, which by Theorem 3.1 is equivalent to the bijection $\mathcal{T}_{\infty} \leftrightarrow \mathcal{W}_{\mathcal{T}}$ in Theorem 5.5, is also proved in [Lovász and Szegedy 09].

We have described the possible limits of sequences of threshold graphs; this makes it easy to see when such sequences converge.

Theorem 5.8. Let G_n be a sequence of threshold graphs such that $v(G_n) \to \infty$. Then G_n converges in $\overline{\mathcal{U}}$ as $n \to \infty$ if and only if the degree distributions $\nu(G_n)$ converge to some distribution μ . In this case, $\mu \in \mathcal{P}_s$ and $G_n \to \Gamma_{\mu}$.

Proof. As in the proof of Theorem 5.5, $G_n \to \Gamma$ if and only if $\pi^*(G_n) \to \Gamma$ in $\mathcal{T}' = \mathcal{T}_{\infty}$, which by Theorem 5.5 holds if and only if $\nu(\pi^*(G_n)) \to \nu(\Gamma)$. By Theorem 4.4, $\nu(\pi^*(G_n))$ equals the distribution of $\int_0^1 W^*(G_n)(U,y) \, \mathrm{d}y$, but this random variable differs by at most $1/\nu(G_n) = o(1)$ from the random variable $\int_0^1 W_{G_n}(U,y) \, \mathrm{d}y$, which has degree distribution $\nu(G_n)$. The result follows. \square

Theorem 5.9. Let G_n be a sequence of bipartite threshold graphs such that

$$v_1(G_n), v_2(G_n) \to \infty.$$

Then G_n converges in $\overline{\mathcal{B}}$ as $n \to \infty$ if and only if the degree distributions $\nu_1(G_n)$ converge to some distribution μ . In this case, $\nu_2(G_n) \to \mu^{\dagger}$ and $G_n \to \Gamma''_{\mu}$.

Proof. $G_n \to \Gamma$ if and only if $\pi(G_n) \to \Gamma$ in $\widetilde{\mathcal{B}} = \mathcal{B}_{\infty\infty}$, which by Theorem 5.1 holds if and only if $\nu_1(\pi(G_n)) \to \nu_1(\Gamma)$. It follows from Theorem 4.4 that $\nu_1(\pi(G_n)) = \nu_1(G_n)$, and the result follows from Theorem 5.1.

Remark 5.10. A threshold graph limit Γ is, by Theorem 5.5, determined by its degree distribution and the fact that it is a threshold graph limit. By Theorem 3.2 and Lemma 4.1, Γ is thus determined by $t(F,\Gamma)$ for F in the set $\{P_4,C_4,2K_2,K_{1,1},K_{1,2},\ldots\}$. It is shown in [Lovász and Szegedy 09] that in some special cases, a finite set of F is enough; for example, the limit defined by the function $W(x,y)=\mathbf{1}[x+y\geq 1]$ (see Example 1.3 and Figure 4) is the unique graph limit with $t(P_4,\Gamma)=t(C_4,\Gamma)=t(2K_2,\Gamma)=0$, $t(K_2,\Gamma)=1/2$, $t(P_3,\Gamma)=1/3$.









6. Random Threshold Graphs

We consider several ways to define random threshold graphs. We will consider only constructions with a fixed number n of vertices; in fact, we take the vertex set to be $[n] = \{1, \ldots, n\}$, where $n \geq 1$ is a given parameter. By a random threshold graph we thus mean a random element of $\mathcal{T}_n := \{G \in \mathcal{T} : V(G) = [n]\}$ for some n; we do not imply any particular construction or distribution unless otherwise stated. (We can regard these graphs as either labeled or unlabeled.)

This section treats four classes of examples: a canonical example based on increasing sets, random-weights examples, random-attachment examples, and uniform random-threshold graphs.

6.1. Increasing Set

For any symmetric increasing $S \subseteq [0,1]^2$, we let $W = \mathbf{1}_S$ and define $T_{n;S} := G(n,W)$ as in Section 3. In other words, we take i.i.d. random variables

$$U_1,\ldots,U_n \sim U(0,1)$$

and draw an edge ij if $(U_i, U_j) \in S$.

As stated in Section 3, $G(n,W) \xrightarrow{\text{a.s.}} \Gamma_W$, which in this case means that $T_{n;S} \xrightarrow{\text{a.s.}} \Gamma(\mathbf{1}_S) \in \mathcal{T}_{\infty}$. We denote $\Gamma(\mathbf{1}_S)$ by Γ_S and have thus the following result, using also Theorem 4.4.

Theorem 6.1. As $n \to \infty$, $T_{n;S} \xrightarrow{\text{a.s.}} \Gamma_S$. In particular, the degree distribution $\nu(T_{n;S})$ converges a.s. to $\nu(\Gamma_S)$, which equals the distribution of

$$\varphi_S(U) := |\{ y : (U, y) \in S \}| = \mathbb{P}((U, U') \in S \mid U),$$
(6.1)

with $U, U' \sim U(0,1)$ independent.

By Theorem 5.5, this construction gives a canonical representation of the limit objects in \mathcal{T}_{∞} , and we may restrict ourselves to closed or open sets as in Theorem 5.5(ii)(iii) to get a unique representation. We can obtain any desired degree distribution $\mu \in \mathcal{P}_s$ for the limit by choosing $S = S_{\mu}$. This construction further gives a canonical representation of random threshold graphs for finite n, provided we make two natural additional assumptions.

Theorem 6.2. Suppose that $(G_n)_1^{\infty}$ is a sequence of random threshold graphs with $V(G_n) = [n]$ such that the distribution of each G_n is invariant under permutations of [n] and that the restriction (induced subgraph) of G_{n+1} to [n] has the







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same distribution as G_n , for every $n \ge 1$. If further $\nu(G_n) \xrightarrow{p} \mu$ as $n \to \infty$, for some $\mu \in \mathcal{P}$, then $\mu \in \mathcal{P}_s$, and for every n, $G_n \stackrel{d}{=} T_{n;S_{\mu}}$.

Proof. It follows from Theorem 5.8 that $G_n \stackrel{p}{\longrightarrow} \Gamma_{\mu}$. (To apply Theorem 5.8 to convergence in probability, we can use the standard trick of considering subsequences that converge a.e., since every subsequence has such a subsubsequence [Kallenberg 02, Lemma 4.2].)

If we represent a graph by its edge indicators, the random graph G_n can be regarded as a family of 0–1-valued random variables indexed by pairs (i,j), $1 \le i < j \le n$. By assumption, these families for different n are consistent, so by the Kolmogorov extension theorem [Kallenberg 02, Theorem 6.16], they can be defined for all n together, which means that there exists a random infinite graph G_{∞} with vertex set $\mathbb N$ whose restriction to [n] coincides (in distribution) with G_n . Moreover, since each G_n is invariant under permutations of the vertices, so is G_{∞} , i.e., G_{∞} is exchangeable.

By the representation theorem by Aldous and Hoover (see [Aldous 85], [Kallenberg 05], and [Diaconis and Janson 08]), every exchangeable random infinite graph can be obtained as a mixture of $G(\infty, W)$; in other words, as $G(\infty, W)$ for some random function $W \in \mathcal{W}_s$. In this case, the subgraphs G_n converge in probability to the corresponding random Γ_W ; see [Diaconis and Janson 08]. Since we have shown that the G_n converge to a deterministic graph limit Γ_μ , we can take W deterministic, so it follows that $G_\infty \stackrel{\mathrm{d}}{=} G(\infty, W)$ for some $W \in \mathcal{W}_s$; moreover, $\Gamma_\mu = \Gamma_W$, and thus by Theorem 5.5 we can choose $W = W_\mu$. (Recall that in general, W is not unique.) Consequently,

$$G_n \stackrel{\mathrm{d}}{=} G(n, W_\mu) = T_{n; S_\mu},$$

which completes the proof.

6.2. Random Weights

Definition (1.1) suggests immediately the construction (1.6): Let X_1, X_2, \ldots , be i.i.d. copies of a random variable X, let $t \in \mathbb{R}$, and let $T_{n;X,t}$ be the threshold graph with vertex set [n] and edges ij for all pairs ij such that $X_i + X_j > t$. (We can without loss of generality let t = 0, by replacing X by X - t/2.) Examples 1.2 and 1.3 are in this mode.

Let $F(x) := \mathbb{P}(X \leq x)$ be the distribution function of X, and let F^{-1} be its right-continuous inverse defined by

$$F^{-1}(u) := \sup\{ x \in \mathbb{R} : F(x) \le u \}$$
 (6.2)









(cf. (5.1), where we consider distributions on [0, 1] only). Thus $-\infty < F^{-1}(u) < \infty$ if 0 < u < 1, while $F^{-1}(1) = \infty$. It is well known that the random variables X_i can be constructed as $F^{-1}(U_i)$ with U_i independent uniformly distributed random variables on (0,1), which leads to the following theorem, showing that this construction is equivalent to the one in Section 6.1 for a suitable set S. Parts of this theorem were found earlier; see [Konno et al. 05, Masuda et al. 05].

Theorem 6.3. Let S be the symmetric increasing set

$$S := \{ (x,y) \in (0,1]^2 : F^{-1}(x) + F^{-1}(y) > t \}.$$
 (6.3)

Then $T_{n;X,t} \stackrel{\mathrm{d}}{=} T_{n;S}$ for every n.

Furthermore, as $n \to \infty$, the degree distribution $\nu(T_{n;X,t})$ converges a.s. to μ , and thus $T_{n;X,t} \xrightarrow{\text{a.s.}} \Gamma_{\mu}$, where $\mu \in \mathcal{P}_s$ is the distribution of the random variable 1 - F(t - X), i.e.,

$$\mu[0, s] = \mathbb{P}(1 - F(t - X) \le s), \quad s \in [0, 1].$$
 (6.4)

Proof. Taking $X_i = F^{-1}(U_i)$, we see that

there is an edge
$$ij \iff F^{-1}(U_i) + F^{-1}(U_j) > t \iff (U_i, U_j) \in S$$
,

which shows that $T_{n;X,t} = T_{n;S}$.

The remaining assertions now follow from Theorem 6.1 together with the calculation, with $U, U' \sim U(0, 1)$ independent and $X = F^{-1}(U), X' = F^{-1}(U'),$

$$\varphi_S(U) = \mathbb{P}((U, U') \in S \mid U) = \mathbb{P}(F^{-1}(U) + F^{-1}(U') > t \mid U)$$

= $\mathbb{P}(X + X' > t \mid X) = \mathbb{P}(X' > t - X \mid X) = 1 - F(t - X).$

The set S defined in (6.3) is in general neither open nor closed; the corresponding open set is

$$S^{\circ} = \left\{ (x, y) \in (0, 1)^2 : F^{-1}(x -) + F^{-1}(y -) > t \right\},\,$$

and the corresponding closed set S_{μ} in Theorem 5.5 can be found as \tilde{S}° from (5.8). If we assume for simplicity that the distribution of X is continuous, then, as is easily verified,

$$S_{\mu} = \{ (x, y) \in [0, 1]^2 : F^{-1}(x) + F^{-1}(y) \ge t \},$$

where we define $F^{-1}(1) = \infty$ (and interpret $\infty + (-\infty) = \infty$ in case $F^{-1}(0) = -\infty$). We can use these sets instead of S in (6.3), since they differ by null sets only.







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6.3. Random Addition of Vertices

Preferential attachment graphs are a rich topic of research in modern graph theory. See the monograph [Lu and Chung 06], along with the survey [Mitzenmacher 04]. The versions in this section are natural because of (1.2) and the construction (1.7).

Let $T_{n,p}$ be the random threshold graph with n vertices obtained by adding vertices one by one with the new vertices chosen as isolated or dominant at random, independently of each other and with a given probability $p \in [0,1]$ of being dominant. (Starting with a single vertex, there are thus n-1 vertex additions.)

The vertices are not equivalent (for example, note that the edges 1i, $i \neq 1$, appear independently, but not the edges ni, $i \neq n$), so we also define the random threshold graph $\widehat{T}_{n,p}$ obtained by a random permutation of the vertices in $T_{n,p}$. (In considering unlabeled graphs, there is no difference between $T_{n,p}$ and $\widehat{T}_{n,p}$.)

Remark 6.4. We may, as stated in (1.7), use different probabilities p_i for different vertices. We leave it to the reader to explore this case, for example with $p_i = f(i/n)$ for some given continuous function $f: [0,1] \to [0,1]$.

Theorem 6.5. The degree distribution $\nu(T_{n,p})$ converges a.s. as $n \to \infty$ to a distribution μ_p that for 0 , has constant density <math>(1-p)/p on (0,p) and p/(1-p) on (p,1); μ_0 is a point mass at 0, and μ_1 is a point mass at 1. In particular, $\mu_{1/2}$ is the uniform distribution on [0,1]. Consequently, $T_{n,p} \xrightarrow{\text{a.s.}} \Gamma_{\mu_p} \in \mathcal{T}_{\infty}$.

Proof. Let $Z_n(t)$ be the number of vertices in $\{1, \ldots, \lfloor nt \rfloor\}$ that are added as dominant. It follows from the law of large numbers that $n^{-1}Z_n(t) \xrightarrow{\text{a.s.}} pt$, uniformly on [0, 1], and we assume this in the sequel of the proof.

If vertex k was added as isolated, it has degree $Z_n(1) - Z_n(k/n)$, since its neighbors are the vertices that later are added as dominant. Similarly, if vertex k was added as dominant, it has degree $k-1+Z_n(1)-Z_n(k/n)$. Consequently, if μ_n is the (normalized) degree distribution of $T_{n,p}$, and ϕ is any continuous function on [0,1], then

$$\int_0^1 \phi(t) \, \mathrm{d}\mu_n(t) = \frac{1}{n} \sum_{k=1}^n \phi(d(k)/n)$$

$$= \frac{1}{n} \sum_{k=1}^n \left(\phi\left(n^{-1} Z_n(1) - n^{-1} Z_n(k/n)\right) \mathbf{1} [\Delta Z_n(k/n) = 0] + \phi\left(n^{-1} Z_n(1) - n^{-1} Z_n(k/n) + (k-1)/n\right) \mathbf{1} [\Delta Z_n(k/n) = 1] \right).$$









Since $n^{-1}Z_n(t) \to pt$ uniformly, and ϕ is uniformly continuous, it follows that as $n \to \infty$,

$$\int_{0}^{1} \phi(t) d\mu_{n}(t) = \frac{1}{n} \sum_{k=1}^{n} \left(\phi \left(p(1 - k/n) \right) \mathbf{1} [\Delta Z_{n}(k/n) = 0] + \phi \left(p(1 - k/n) + k/n \right) \mathbf{1} [\Delta Z_{n}(k/n) = 1] \right) + o(1)$$

$$= \int_{0}^{1} \phi \left(p(1 - t) \right) d \left(n^{-1} \lfloor nt \rfloor - n^{-1} Z_{n}(t) \right)$$

$$+ \int_{0}^{1} \phi \left(p(1 - t) + t \right) d \left(n^{-1} Z_{n}(t) \right) + o(1).$$

Since the convergence $n^{-1}Z_n(t) \to pt$ implies (weak) convergence of the corresponding measures, we finally obtain, as $n \to \infty$,

$$\int_{0}^{1} \phi(t) d\mu_{n}(t) \to \int_{0}^{1} \phi(p(1-t))(1-p) dt + \int_{0}^{1} \phi(p(1-t)+t) p dt$$

$$= \frac{1-p}{p} \int_{0}^{p} \phi(x) dx + \frac{p}{1-p} \int_{p}^{1} \phi(x) dx$$

$$= \int_{0}^{1} \phi(x) d\mu_{p}(x),$$

with obvious modifications if p = 0 or p = 1.

Let $S_p := S_{\mu_p}$ be the corresponding subset of $[0,1]^2$. If $0 , <math>\mu_p$ has the distribution function

$$F_{\mu_p}(x) = \begin{cases} \frac{1-p}{p}x, & 0 \le x \le p, \\ 1 - \frac{p}{1-p}(1-x), & p \le x \le 1, \end{cases}$$
 (6.5)

then it follows from (5.2) that S_p is the quadrilateral with vertices (0,1), (1-p, 1-p), (1,0), and (1,1); see Figure 8. In the special case p=1/2, μ_p is the uniform distribution on [0,1], and S_p is the triangle $\{(x,y) \in [0,1]^2 : x+y \ge 1\}$ pictured in Figure 4 with vertices (0,1), (1,0), and (1,1). Finally, S_0 consists of the upper and right edges only, and $S_1 = [0,1]^2$.

Removing any vertex from $T_{n,p}$ (and relabeling the remaining ones) yields $T_{n-1,p}$. It follows that the same property holds for $\widehat{T}_{n,p}$, so $\widehat{T}_{n,p}$ satisfies the assumptions of Theorem 6.2. Since $\widehat{T}_{n,p}$ has the same degree distribution as $T_{n,p}$, Theorems 6.2 and 6.5 yield the following equality.

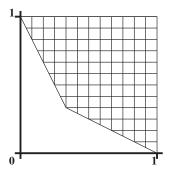
Corollary 6.6. If
$$0 \le p \le 1$$
 and $n \ge 1$, then $\widehat{T}_{n,p} \stackrel{\mathrm{d}}{=} T_{n;S_p}$.











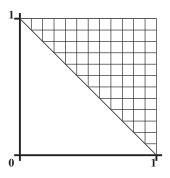


Figure 8. Two examples of the sets S_p ; the one on the right shows the special case p = 0.5.

Hence the random threshold graphs in this subsection are special cases of the general construction in Section 6.1. We can also construct them using random weights as in Section 6.2.

Corollary 6.7. If $0 \le p \le 1$ and $n \ge 1$, then $\widehat{T}_{n,p} \stackrel{d}{=} T_{n;X,0}$, where X has the density 1-p on (-1,0) and p on (0,1).

Proof. A simple calculation shows that the set S given by (6.3) is the quadrilateral S_p .

We may transform X by a linear map; for example, we may equivalently take X with density 2(1-p) on (0,1/2) and 2p on (1/2,1), with the threshold t=1. In particular, $\widehat{T}_{n,1/2} \stackrel{\mathrm{d}}{=} T_{n;U,1}$, where $U \sim U(0,1)$ as in Example 1.3.

6.4. Uniform Random Threshold Graphs

Let T_n be a random unlabeled threshold graph of order n with the uniform distribution studied in Section 2. Similarly, let T_n^L be a random labeled threshold graph of order n with the uniform distribution. Although T_n and T_n^L have different distributions, see Section 2, the next theorem shows that they have the same limit as $n \to \infty$.

Theorem 6.8. The degree distributions $\nu(T_n)$ and $\nu(T_n^L)$ both converge in probability to the uniform distribution λ on [0,1]. Hence $T_n \stackrel{\mathrm{p}}{\longrightarrow} \Gamma_{\lambda}$ and $T_n^L \stackrel{\mathrm{p}}{\longrightarrow} \Gamma_{\lambda}$.

By Section 2.1, $T_n \stackrel{\text{d}}{=} T_{n,1/2}$; hence the result for unlabeled graphs follows from Theorem 6.5.









Proof. We use Theorem 2.3; in fact, the proof works for random threshold graphs generated by Algorithm 3 for any i.i.d. random variables B_2, B_3, \ldots with finite mean, and any B_1 . (In the case that B_2 is always a multiple of some d > 1, there is a trivial modification.) Let $\beta := \mathbb{E} B_2$.

The algorithm starts by choosing (random) block lengths B_1, B_2, \ldots until their sum is at least n, and then rejects them and restarts (step 3) unless the sum is exactly n. It is simpler to ignore this check, so we consider the following modified algorithm: Take B_1, B_2, \ldots as above. Let $S_k := \sum_{j=1}^k B_j$ be their partial sums and let $\tau(n) := \min\{k : S_k \ge n\}$. Toss a coin to determine whether the first block is isolated or dominant, and construct a random threshold graph by adding $\tau(n)$ blocks of vertices with $B_1, \ldots, B_{\tau(n)}$ elements, alternately isolated and dominant.

This gives a random graph \widetilde{G}_n with $S_{\tau(n)}$ vertices, but conditioned on $S_{\tau(n)} = n$; we obtain the desired random threshold graph (cf. Theorem 2.3). Since $\mathbb{P}(S_{\tau(n)} = n)$ converges to $1/\beta > 0$ by renewal theory, it suffices to prove that $\nu(\widetilde{G}_n) \stackrel{\mathrm{p}}{\longrightarrow} \lambda$ as $n \to \infty$. In fact, we will show that $\nu(\widetilde{G}_n) \stackrel{\mathrm{a.s.}}{\longrightarrow} \lambda$ if we first choose an infinite sequence B_1, B_2, \ldots and then let $n \to \infty$.

Let $S_m^O := \sum_{2k+1 \le m} B_{2k+1}$ and $S_m^E := \sum_{2k \le m} B_{2k}$ be the partial sums of the odd and even terms. By the law of large numbers, a.s. $S_n/n \to \beta$ and $S_n^O/n \to \frac{1}{2}\beta$, $S_n^E/n \to \frac{1}{2}\beta$. We now consider a fixed sequence $(B_j)_1^\infty$ such that these limits hold. Since $S_{\tau(n)-1} < n \le S_{\tau(n)}$, it follows, as is well known, that $n/\tau(n) \to \beta$, so $\tau(n) = n/\beta + o(n)$.

Suppose for definiteness that the first block is chosen to be isolated; then every odd block is isolated and every even block is dominant. (In the opposite case, interchange even and odd below.) If $i \in (S_{2k}, S_{2k+1}]$, then i belongs to block 2k+1, so i is added as isolated, and the neighbors of i will be only the vertices added after i as dominant, i.e., $\bigcup_{k<\ell<\tau(n)/2}(S_{2\ell-1}, S_{2\ell}]$, and

$$d(i) = \sum_{2k < 2\ell < \tau(n)} B_{2\ell} = S_{\tau(n)}^E - S_{\tau(i)}^E.$$

If instead $i \in (S_{2k-1}, S_{2k}]$, then i is also joined to all vertices up to S_{2k} , and thus

$$d(i) = \sum_{2\ell \le \tau(n)} B_{2\ell} + \sum_{2\ell+1 \le \tau(i)} B_{2\ell+1} = S_{\tau(n)}^E + S_{\tau(i)}^O.$$

Hence if i is in an odd block,

$$\frac{d(i)}{n} = \frac{1}{n} \Big(\tau(n) \frac{\beta}{2} - \tau(i) \frac{\beta}{2} + o(n) \Big) = \frac{n - i + o(n)}{2n} = \frac{1}{2} - \frac{i}{2n} + o(1),$$

and if i is in an even block, similarly,

$$\frac{d(i)}{n} = \frac{1}{2} + \frac{i}{2n} + o(1).$$







Now fix $t \in (0, 1/2)$ and let $\varepsilon > 0$. Then the following holds if n is large enough:

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If i is in an even block, then $d(i)/n \ge 1/2 + o(1) > t$. If i is in an odd block and $i \le i_1 := (1 - 2t - 2\varepsilon)n$, then $d(i)/n = \frac{1}{2}(n-i)/n + o(1) \ge t + \varepsilon + o(1) > t$. If i is in an odd block and $i \geq i_2 := (1 - 2t + 2\varepsilon)n$, then $d(i)/n = \frac{1}{2}(n-i)/n + o(1) \leq i$ $t-\varepsilon+o(1) < t$. Consequently, for large $n, d(i)/n \le t$ only if i is in an odd block $(S_{2k}, S_{2k+1}]$, and in this case $2k+1 > \tau(i_1)$ is necessary and $2k+1 > \tau(i_2)$ is sufficient. Hence

$$S_{\tau(n)}^O - S_{\tau(i_2)}^O \leq |\{\,i: d(i)/n \leq t\,\}| \leq S_{\tau(n)}^O - S_{\tau(i_1)}^O.$$

Since $\nu(\widetilde{G}_n)[0,t] = \frac{1}{n} |\{i: d(i)/n \leq t\}|$ and

$$\frac{1}{n} \left(S_{\tau(n)}^O - S_{\tau(i_j)}^O \right) = \frac{\beta(\tau(n) - \tau(i_j)) + o(n)}{2n} = \frac{n - i_j + o(n)}{2n} = t \pm \varepsilon + o(1),$$

it follows that

$$t - \varepsilon + o(1) \le \nu(\widetilde{G}_n)[0, t] \le t + \varepsilon + o(1).$$

Since ε is arbitrary, this shows that $\nu(\widetilde{G}_n)[0,t] \to t$, for every $t \in (0,\frac{1}{2})$. We clearly obtain the same result if the first block is dominant.

For $t \in (\frac{1}{2}, 1)$ we can argue similarly, now analyzing the dominant blocks. Alternatively, we may apply the result just obtained to the complement of G_n , which is obtained from the same B_j by switching the types of the blocks. This shows that $\nu(\widetilde{G}_n)[0,t] \to t$ for $t \in (\frac{1}{2},1)$ too.

Hence, $\nu(\tilde{G}_n)[0,t] \to t$ for every $t \in (0,1)$ except possibly $\frac{1}{2}$, which shows that $\nu(G_n) \to \lambda$.

Vertex Degrees in Uniform Random Threshold Graphs

We have seen in Theorem 6.8 that the normalized degree distributions $\nu(T_n)$ and $\nu(T_n^L)$ for uniform unlabeled and labeled random threshold graphs both converge to the uniform distribution on [0,1]. This is for weak convergence of distributions in \mathcal{P} , which is equivalent to averaging over degrees in intervals (an, bn); we here refine this by studying individual degrees.

Let $N_d(G)$ be the number of vertices of degree d in the graph G. Thus D_G , the degree of a random vertex in G, has distribution $\mathbb{P}(D_G = d) = N_d/v(G)$. (Recall that $\nu(G)$ is the distribution of $D_G/\nu(G)$; see Section 4.)

We will study the random variables $N_d(T_n)$ and $N_d(T_n^L)$ describing the numbers of vertices of a given degree d in a uniform random unlabeled or labeled threshold graph, and in particular their expectations $\mathbb{E} N_d(T_n)$ and $\mathbb{E} N_d(T_n^L)$;









note that $\mathbb{E} N_d(T_n)/n$ and $\mathbb{E} N_d(T_n^L)/n$ are the probabilities that a given (or random) vertex in the random graph T_n or T_n^L has degree d. By symmetry under complementation,

$$N_d(T_n) \stackrel{\mathrm{d}}{=} N_{n-1-d}(T_n)$$
 and $N_d(T_n^L) \stackrel{\mathrm{d}}{=} N_{n-1-d}(T_n^L)$.

Let us first look at N_0 , the number of isolated vertices. (By symmetry, we have the same results for N_{n-1} , the number of dominant vertices). Note that for every $n \geq 2$, $\mathbb{P}(N_0(T_n) = 0) = \mathbb{P}(N_0(T_n^L) = 0) = 1/2$ by symmetry.

Theorem 7.1.

(i) For any $n \geq 1$,

$$\mathbb{P}(N_0(T_n) = j) = \begin{cases}
2^{-j-1}, & 0 \le j \le n-2, \\
0, & j = n-1, \\
2^{-n+1}, & j = n.
\end{cases}$$
(7.1)

In other words, if $X \sim \text{Ge}(1/2)$, then $N_0(T_n) \stackrel{\text{d}}{=} X'_n$, where $X'_n := X_n$ if x < n-1 and $X'_n := n$ if $X_n \ge n-1$. Furthermore, $\mathbb{E} N_0(T_n) = 1$, and $N_0(T_n) \stackrel{\text{d}}{\longrightarrow} \text{Ge}(1/2)$ as $n \to \infty$, with convergence of all moments.

(ii) $\mathbb{P}(N_0(T_n^L) = j) = t(n, j)/t(n)$, where t(n, j) is given by (2.6); in particular, if $0 \le j \le n-2$, then

$$\mathbb{P}(N_0(T_n^L) = j) = \frac{1}{2j!} \frac{t(n-j)/(n-j)!}{t(n)/n!} = \frac{1}{2j!} (\log 2)^j (1 + O(\rho^{n-j}))$$

with $\rho = \log 2/(2\pi) \approx 0.11$. Hence $N_0(T_n^L) \stackrel{d}{\longrightarrow} \operatorname{Po}(\log 2)$ as $n \to \infty$ with convergence of all moments; in particular, $\mathbb{E} N_0(T_n^L) \to \log 2$.

- **Proof.** (i): A threshold graph has j isolated vertices if and only if the extended binary code $\alpha_1 \cdots \alpha_n$ in Section 2 ends with exactly j 0's. For a random unlabeled threshold graph T_n , the binary code $\alpha_2 \cdots \alpha_n$ is uniformly distributed, and thus (7.1) follows. The remaining assertions follow directly.
- (ii): In the labeled case, the exact distribution is given by (2.6), and the asymptotics follow by (2.3). Uniform integrability of any power $N_0(T_n^L)^m$ follows by the same estimates, and thus moment convergence holds.

For higher degrees, we begin with an exact result for the unlabeled case.

Theorem 7.2.
$$\mathbb{E} N_d(T_n) = 1$$
 for every $d = 0, \ldots, n-1$.







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Actually, this is the special case p = 1/2 of a more general theorem for the random threshold graph $T_{n,p}$ defined in Section 6.3 (cf. Theorem 6.5, which is for weak convergence, but on the other hand yields an a.s. limit while we here study the expectations).

Theorem 7.3. Let 0 . If <math>q = 1 - p and $X \sim \text{Bin}(n, p)$, then for $0 \le d \le n - 1$, $\mathbb{E} \, N_d(T_{n,p}) = \frac{q}{p} + \left(\frac{p}{q} - \frac{q}{p}\right) \mathbb{P}(X \le d).$

Proof. We use the definition in Section 6.3. (For the uniform case p = 1/2, this is Algorithm 1.) Let d_i be the degree of vertex i. Then if $\alpha_1 \cdots \alpha_n$ is the extended binary code of the graph, we have

$$d_i = (i-1)\alpha_i + \sum_{j=i+1}^n \alpha_j.$$

Since the α_i are i.i.d. Be(p) for $i=2,\ldots,n$, the probability-generating function of d_i is

$$\mathbb{E} x^{d_i} = \mathbb{E} x^{(i-1)\alpha_i} \prod_{j=i+1}^n \mathbb{E} x^{\alpha_j} = (px^{i-1} + q)(px + q)^{n-i}.$$

Consequently,

$$\sum_{d} \mathbb{E} N_{d}(T_{n,p}) x^{d} = \sum_{i=1}^{n} \mathbb{E} x^{d_{i}} = \sum_{i=1}^{n} p x^{i-1} (px+q)^{n-i} + \sum_{i=1}^{n} q (px+q)^{n-i}$$

$$= p \frac{x^{n} - (px+q)^{n}}{x - (px+q)} + q \frac{1 - (px+q)^{n}}{1 - (px+q)}$$

$$= \frac{(q/p) + (p/q - q/p)(px+q)^{n} - (p/q)x^{n}}{1 - x}.$$

In the special case p=1/2, this is $(1-x^n)/(1-x)=\sum_{d=0}^{n-1}x^d$, which proves Theorem 7.2 by identifying coefficients. For general p, Theorem 7.3 follows in the same way.

Recall that R_d denotes the number of preferential arrangements, or surjection numbers, given in (2.4).

Theorem 7.4.

(i) In the unlabeled case, for any sequence d = d(n) with $0 \le d \le n - 1$, $N_d(T_n) \stackrel{d}{\longrightarrow} Ge(1/2)$ with convergence of all moments.









(ii) In the labeled case, let X_d , $0 \le d \le \infty$, have the modified Poisson distribution given by

$$\mathbb{P}(X_d = \ell) = \begin{cases} \frac{\gamma_d}{\log 2} \, \mathbb{P}\big(\operatorname{Po}(\log 2) = \ell \big) = \gamma_d \frac{(\log 2)^{\ell - 1}}{2 \cdot \ell!}, & \ell \ge 1, \\ 1 - \frac{\gamma_d}{2 \log 2}, & \ell = 0, \end{cases}$$

where $\gamma_0 := \log 2$, $\gamma_d := 2R_d(\log 2)^{d+1}/d!$ for $d \ge 1$, and $\gamma_\infty := 1$. Then for every fixed $d \ge 0$, $N_d(T_n^L) \stackrel{\mathrm{d}}{=} N_{n-1-d}(T_n^L) \stackrel{\mathrm{d}}{\longrightarrow} X_d$, and for every sequence $d = d(n) \to \infty$ with $n - d \to \infty$, $N_d(T_n^L) \stackrel{\mathrm{d}}{=} N_{n-1-d}(T_n^L) \stackrel{\mathrm{d}}{\longrightarrow} X_\infty$ as $n \to \infty$, in both cases with convergence of all moments.

In particular, $\mathbb{E} N_d(T_n^L) = \mathbb{E} N_{n-1-d}(T_n^L)$ converges to γ_d for every fixed d, and to $\gamma_{\infty} = 1$ if $d \to \infty$ and $n - d \to \infty$.

In the labeled case we thus have, in particular,

$$\mathbb{E} N_0(T_n^L) \to \log 2 \approx 0.69315,$$

$$\mathbb{E} N_1(T_n^L) \to 2(\log 2)^2 \approx 0.96091,$$

$$\mathbb{E} N_2(T_n^L) \to 3(\log 2)^3 \approx 0.99907,$$

$$\mathbb{E} N_3(T_n^L) \to \frac{13}{3}(\log 2)^4 \approx 1.00028.$$

The values for degrees 0 and 1 (and symmetrically n-1 and n-2) are thus substantially smaller than 1, which is clearly seen in Figure 7. (We can regard this as an edge effect; the vertices with degrees close to 0 or n-1 are those added last in Algorithm 3. Figure 7 also shows an edge effect at the other side; there is a small bump for degrees around n/2, which correspond to the vertices added very early in the algorithm; this bump vanishes asymptotically, as shown by Theorem 7.4; we believe that it has height of order $n^{-1/2}$ and width of order $n^{1/2}$, but we have not analyzed it in detail.)

Proof. The cases d=0 and d=n-1 follow from Theorem 7.1. We may thus suppose $1 \le d \le n-2$. We use Algorithm 3. We know that vertices in each block have the same degree, while different blocks have different degrees; thus there is at most one block with degree d.

Let $p_d(\ell)$ be the probability that there is such a block of length $\ell \geq 1$, and that this block is added as isolated. By symmetry, the probability that there is a dominant block of length ℓ with degree d is p_{n-1-d} , and thus

$$\mathbb{P}(N_d = \ell) = p_d(\ell) + p_{n-1-d}(\ell), \quad \ell \ge 1.$$
 (7.2)







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If block j is an isolated block, then the degree of the vertices in it equals the number of vertices added as dominant after it, i.e., $B_{j+1} + B_{j+3} + \cdots + B_{j+2k-1}$, if the total number τ of blocks is j+2k-1 or j+2k. Consequently, there is an isolated block of length ℓ with vertices of degree d if and only if there exist $j \geq 1$ and $k \geq 1$ with

- $B_j = \ell$,
- block j is isolated,
- $\sum_{i=1}^{k} B_{j+2i-1} = d$,
- $\sum_{i=1}^{j+2k-1} B_i = n$ or $\sum_{i=1}^{j+2k} B_i = n$.

Recall that B_1, B_2, \ldots are independent and that B_2, B_3, \ldots have the same distribution, while B_1 has a different one. (The distributions differ between the unlabeled and labeled cases.) Let

$$\hat{S}_n := \sum_{i=1}^m B_i$$
 and $S_n := \sum_{i=1}^m B_{i+1}$.

Further, let

$$\hat{u}(n) = \sum_{m=0}^{\infty} \mathbb{P}(\hat{S}_m = n) = \mathbb{P}(B_{\tau} = n) = \mathbb{P}\left(\sum_{i=1}^{\tau} B_i = n\right),$$

$$u(n) = \sum_{m=0}^{\infty} \mathbb{P}(S_m = n),$$

and recall that $u(n), \hat{u}(n) \to 1/\mu := 1/\mathbb{E}B_2$ (exponentially fast) by standard renewal theory (for example by considering generating functions). For any $m \ge j + 2k - 1$,

$$\sum_{i=1}^{m} B_i - B_j - \sum_{i=1}^{k} B_{j+2i-1} \stackrel{\mathrm{d}}{=} \begin{cases} S_{m-1-k}, & j = 1, \\ \hat{S}_{m-1-k}, & j \ge 2, \end{cases}$$

and it follows that since B_1, B_2, \ldots are independent and we condition on $\hat{S}_{\tau} = n$,

$$p_{d}(\ell) = \frac{1}{2\hat{u}(n)} \left\{ \sum_{k=1}^{\infty} \mathbb{P}(B_{1} = \ell) \, \mathbb{P}(S_{k} = d) \right.$$

$$\cdot \left(\mathbb{P}(S_{k-1} = n - \ell - d) + \mathbb{P}(S_{k} = n - \ell - d) \right)$$

$$+ \sum_{j=2}^{\infty} \sum_{k=1}^{\infty} \mathbb{P}(B_{j} = \ell) \, \mathbb{P}(S_{k} = d)$$

$$\cdot \left(\mathbb{P}(\hat{S}_{j+k-2} = n - \ell - d) + \mathbb{P}(\hat{S}_{j+k-1} = n - \ell - d) \right) \right\}$$

$$(7.3)$$









In the double sum, $\mathbb{P}(B_j = \ell) = \mathbb{P}(B_2 = \ell)$ does not depend on j, so the sum is at most

$$\mathbb{P}(B_2 = \ell) \sum_k \mathbb{P}(S_k = d) 2\hat{u}(n - \ell - d) = 2 \,\mathbb{P}(B_2 = \ell) u(d) \hat{u}(n - \ell - d)$$
$$= O(\mathbb{P}(B_2 = \ell)).$$

Similarly, the first sum is $O(\mathbb{P}(B_1 = \ell)) = O(\mathbb{P}(B_2 = \ell))$, and it follows that $p_d(\ell) = O(\mathbb{P}(B_2 = \ell))$ and thus, by (7.2),

$$\mathbb{P}(N_d = \ell) = O(\mathbb{P}(B_2 = \ell)),\tag{7.4}$$

uniformly in n, d, and ℓ . This shows tightness, so convergence $\mathbb{P}(N_d = \ell) \to \mathbb{P}(X = \ell)$ for some nonnegative-integer-valued random variable X and each fixed $\ell \geq 1$ implies convergence in distribution (i.e., for $\ell = 0$ too). Further, since all moments of B_2 are finite, (7.4) implies that all moments $\mathbb{E}[N_d^m]$ are bounded, uniformly in d and n; hence convergence in distribution implies that all moments converge too. In the rest of the proof we thus let $\ell \geq 1$ be fixed.

If $d \leq n/2$, it is easy to see that

$$\mathbb{P}(S_{k-1} = n - \ell - d) + \mathbb{P}(S_k = n - \ell - d) = O((n - \ell - d)^{-1/2}) = O(n^{-1/2}),$$

uniformly in k, so the first sum in (7.3) is $O(n^{-1/2}u(d)) = O(n^{-1/2})$. If d > n/2, we similarly have $\mathbb{P}(S_k = d) = O(d^{-1/2}) = O(n^{-1/2})$, and thus the sum is $O(n^{-1/2}u(n-\ell-d)) = O(n^{-1/2})$. Hence (7.3) yields

$$p_{d}(\ell) = O(n^{-1/2}) + \frac{\mathbb{P}(B_{2} = \ell)}{2\hat{u}(n)} \sum_{k=1}^{\infty} \mathbb{P}(S_{k} = d)$$

$$\cdot \left(\sum_{i=k} \mathbb{P}(\hat{S}_{i} = n - \ell - d) + \sum_{i=k+1} \mathbb{P}(\hat{S}_{i} = n - \ell - d) \right).$$
(7.5)

The term with i = k can be taken twice, just as those with i > k, since $\sum_{k} \mathbb{P}(S_{k} = d) \mathbb{P}(\hat{S}_{k} = n - \ell - d) = O(n^{-1/2})$ by the same argument as for the first sum in (7.3). Further, for $i \geq k$, $\hat{S}_{i} - \hat{S}_{k} \stackrel{\text{d}}{=} S_{i-k}$, and this is independent of \hat{S}_{k} ; thus

$$\mathbb{P}(\hat{S}_{i} = n - \ell - d) = \mathbb{P}(S_{i-k} = n - \ell - d - \hat{S}_{k}) = \mathbb{E}\,\mathbb{P}(S_{i-k} = n - \ell - d - \hat{S}_{k} \mid \hat{S}_{k})$$

and

$$\sum_{i=k}^{\infty} \mathbb{P}(\hat{S}_i = n - \ell - d) = \mathbb{E} u(n - \ell - d - \hat{S}_k).$$







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Hence (7.5) yields

$$p_d(\ell) = \frac{\mathbb{P}(B_2 = \ell)}{\hat{u}(n)} \sum_{k=1}^{\infty} \mathbb{P}(S_k = d) \, \mathbb{E} \, u(n - \ell - d - \hat{S}_k) + O(n^{-1/2}). \tag{7.6}$$

If d is fixed, then $\mathbb{E}u(n-\ell-d-\hat{S}_k)\to\mu^{-1}$ by dominated convergence as $n\to\infty$ for each k, and thus (7.6) yields, by dominated convergence again,

$$p_d(\ell) \to \mathbb{P}(B_2 = \ell) \sum_{k=1}^{\infty} \mathbb{P}(S_k = d) = u(d) \, \mathbb{P}(B_2 = \ell). \tag{7.7}$$

If $d \to \infty$, we use the fact that $u(m) - \mathbf{1}[m \ge 0]\mu^{-1}$ is summable over \mathbb{Z} to see that

$$\mathbb{E} u(n-\ell-d-\hat{S}_k) - \mu^{-1} \mathbb{P}(n-\ell-d-\hat{S}_k \ge 0) = O(\max_m \mathbb{P}(\hat{S}_k = m)),$$

which tends to 0 as $k \to \infty$; on the other hand, $\mathbb{P}(S_k = d) \to 0$ for every fixed k. It follows that (7.6) yields

$$p_d(\ell) = \mathbb{P}(B_2 = \ell) \sum_{k=1}^{\infty} \mathbb{P}(S_k = d) \, \mathbb{P}(\hat{S}_k \le n - \ell - d) + o(1).$$

If $\tau_d := \min\{k : S_k \geq d\}$ and \hat{S}'_k denotes a copy of \hat{S}_k independent of $\{S_j\}_1^{\infty}$, then

$$\sum_{k=1}^{\infty} \mathbb{P}(S_k = d) \, \mathbb{P}(\hat{S}_k \le n - \ell - d) = u(d) \, \mathbb{P}(\hat{S}'_{\tau_d} \le n - \ell - d \mid S_{\tau_d} = d).$$

It is easy to see that with $\sigma^2 := \operatorname{Var}(B_2)$, as $d \to \infty$,

$$\left((\hat{S}'_{\tau_d} - d) / \sqrt{d} \mid S_{\tau_d} = d \right) = \left((\hat{S}'_{\tau_d} - S_{\tau_d}) / \sqrt{d} \mid S_{\tau_d} = d \right) \stackrel{\mathrm{d}}{\longrightarrow} N(0, 2\sigma^2/\mu);$$

cf. [Gut and Janson 83] (the extra conditioning on $S_{\tau_d} = d$ makes no difference). Hence when $d \to \infty$,

$$p_d(\ell) = \mathbb{P}(B_2 = \ell)u(d)\Phi((n - \ell - 2d)/\sqrt{d}) + o(1).$$

(By (7.7), this holds for fixed d too.)

We next observe that $\Phi((n-\ell-2d)/\sqrt{d}) = \Phi((n-2d)/\sqrt{n/2}) + o(1)$; this is easily seen by considering separately the three cases $d/n \to a \in [0,1/2)$, $d/n \to a \in (1/2,1]$, and $d/n \to 1/2$ and $(n-2d)/\sqrt{n/2} \to b \in [-\infty,\infty]$ (the









general case follows by considering suitable subsequences). Hence, we have when $d \to \infty$, recalling that then $u(d) \to \mu^{-1}$,

$$p_d(\ell) = \mu^{-1} \mathbb{P}(B_2 = \ell) \Phi((n - 2d) / \sqrt{n/2}) + o(1).$$

For fixed d, this implies that $p_{n-d-1}(\ell) \to 0$, and thus (7.2) and (7.7) yield

$$\mathbb{P}(N_d = \ell) = p_d(\ell) + p_{n-1-d}(\ell) = u(d)\,\mathbb{P}(B_2 = \ell) + o(1).$$

Similarly, if $d \to \infty$ and $n - d \to \infty$, then

$$\mathbb{P}(N_d = \ell) = p_d(\ell) + p_{n-1-d}(\ell)$$

$$= \mu^{-1} \mathbb{P}(B_2 = \ell) \left(\Phi\left((n - 2d) / \sqrt{n/2} \right) + \Phi\left((2d + 2 - n) / \sqrt{n/2} \right) \right) + o(1)$$

$$= \mu^{-1} \mathbb{P}(B_2 = \ell) + o(1).$$

We have thus proven convergence as $n \to \infty$, with all moments, $N_d \stackrel{\mathrm{d}}{\longrightarrow} X_d$ for fixed d and $N_d \stackrel{\mathrm{d}}{\longrightarrow} X_\infty$ for $d = d(n) \to \infty$ with $n - d \to \infty$, where

$$\mathbb{P}(X_d = \ell) = u(d)\,\mathbb{P}(B_2 = \ell) = 2u(d)\,\mathbb{P}(B^* = \ell), \quad \ell \ge 1,\tag{7.8}$$

$$\mathbb{P}(X_d = 0) = 1 - \mathbb{P}(X_d \ge 1) = 1 - u(d), \tag{7.9}$$

for $1 \le d \le \infty$, with $u(\infty) := \mu^{-1}$.

In the unlabeled case, $B_2 = (B^* \mid B^* \geq 1) \stackrel{\mathrm{d}}{=} B^* + 1$ with $B^* \sim \mathrm{Ge}(1/2)$. Consider a random infinite string $\alpha_1\alpha_2\cdots$ of i.i.d. $\mathrm{Be}(1/2)$ binary digits, and define a block as a string of $m \geq 0$ 0's followed by a single 1. Then $B_{j+1}, j \geq 1$, can be interpreted as the successive block lengths in $\alpha_1\alpha_2\cdots$, and thus u(d) is the probability that some block ends at d, i.e., $u(d) = \mathbb{P}(\alpha_d = 1) = 1/2$, for every $d \geq 1$. It follows from (7.8)–(7.9) that $X_d \stackrel{\mathrm{d}}{=} B^* \sim \mathrm{Ge}(1/2)$ for every $d \geq 1$, and (i) follows.

In the labeled case, when $B^* \sim \text{Po}(\log 2)$, we use generating functions:

$$\sum_{d=0}^{\infty} u(d)x^d = \sum_{k=0}^{\infty} \mathbb{E} x^{S_k} = \sum_{k=0}^{\infty} (\mathbb{E} x^{B_2})^k = \frac{1}{1 - \mathbb{E} x^{B_2}} = \frac{\mathbb{P}(B^* \ge 1)}{1 - \mathbb{E} x^{B^*}}$$
$$= \frac{1/2}{1 - e^{(x-1)\log 2}} = \frac{1}{2 - e^{x\log 2}} = \sum_{d=0}^{\infty} \frac{R_d}{d!} (x \log 2)^d,$$

where we recognize the generating function (2.5). Thus, $u(d) = R_d(\log 2)^d/d!$. (A direct combinatorial proof of this is also easy.)

We let, using
$$\mu := \mathbb{E} B_2 = \mathbb{E} B^* / \mathbb{P}(B^* \ge 1) = 2 \log 2$$
,

$$\gamma_d := \mathbb{E} X_d = u(d) \mathbb{E} B_2 = \mu u(d) = 2 \log 2u(d) = 2R_d(\log 2)^{d+1}/d!$$









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and note that $\gamma_d \to \gamma_\infty = 1$ as $d \to \infty$, since $u(d) \to \mu^{-1}$, or by the known asymptotics of R_d [Flajolet and Sedgewick 09, (II.16)]. The description of X_d in the statement now follows from (7.8)–(7.9).

8. Random Bipartite Threshold Graphs

The constructions and results in Section 6 have analogues for bipartite threshold graphs. The proofs are simple modifications of those above and are omitted.

8.1. Increasing Set

For any increasing $S \subseteq [0,1]^2$, define $T_{n_1,n_2;S} := G(n_1,n_2,\mathbf{1}_S)$. In other words, take i.i.d. random variables $U_1',\ldots,U_{n_1}',U_1'',\ldots,U_{n_2}'' \sim U(0,1)$ and draw an edge ij if $(U_i',U_j'') \in S$.

Theorem 8.1. As $n_1, n_2 \to \infty$, $T_{n_1, n_2; S} \xrightarrow{\text{a.s.}} \Gamma_S''$. In particular, the degree distribution $\nu_1(T_{n;S})$ converges a.s. to $\nu_1(\Gamma_S'')$, which equals the distribution of $\varphi_S(U)$ defined by (6.1).

As in Section 6, this gives a canonical representation of random bipartite threshold graphs under natural assumptions.

Theorem 8.2. Suppose that $(G_{n_1,n_2})_{n_1,n_2\geq 1}$ are random bipartite threshold graphs with $V_1(G_{n_1,n_2})=[n_1]$ and $V_2(G_{n_1,n_2})=[n_2]$ such that the distribution of each G_{n_1,n_2} is invariant under permutations of V_1 and V_2 and that the restrictions (induced subgraphs) of G_{n_1+1,n_2} and G_{n_1,n_2+1} to V(G) both have the same distribution as G_{n_1,n_2} , for every $n_1,n_2\geq 1$. If further $\nu_1(G_{n_1,n_2})\stackrel{\mathrm{P}}{\longrightarrow} \mu$ as $n_1,n_2\to\infty$, for some $\mu\in\mathcal{P}$, then for every n_1,n_2 , $G_{n_1,n_2}\stackrel{\mathrm{d}}{=} T_{n_1,n_2;S_\mu}$.

8.2. Random Weights

Definition (1.9) suggests the following construction:

(8.1) Let X and Y be two random variables and let $t \in \mathbb{R}$. Let X_1, X_2, \ldots , be copies of X, let Y_1, Y_2, \ldots , be copies of Y, all independent, and let $T_{n_1, n_2; X, Y, t}$ be the bipartite threshold graph with vertex sets $[n_1]$ and $[n_2]$ and edges ij for all pairs ij such that $X_i + Y_j > t$.

Theorem 8.3. Let S be the increasing set

$$S := \{ (x, y) \in (0, 1]^2 : F_X^{-1}(x) + F_Y^{-1}(y) > t \}.$$
 (8.2)









Then $T_{n_1,n_2;X,Y,t} \stackrel{d}{=} T_{n_1,n_2;S}$ for every $n_1, n_2 \ge 1$.

Furthermore, as $n_1, n_2 \to \infty$, the degree distribution $\nu_1(T_{n_1,n_2;X,Y,t})$ converges a.s. to μ and thus $T_{n_1,n_2;X,Y,t} \xrightarrow{\text{a.s.}} \Gamma''_{\mu}$, where $\mu \in \mathcal{P}$ is the distribution of the random variable $1 - F_Y(t - X)$, i.e.,

$$\mu[0, s] = \mathbb{P}(1 - F_Y(t - X) \le s), \quad s \in [0, 1].$$
 (8.3)

In the special case that $\mathbb{P}(X \in [0,1]) = 1$, $Y \sim U(0,1)$, and t = 1, (8.3) yields $\mu[0,s] = \mathbb{P}(X \leq s)$, so μ is the distribution of X; further, the set S in (8.2) is a.e. equal to S_{μ} in (5.3).

Corollary 8.4. If $\mu \in \mathcal{P}_s$, let X have distribution μ and let $Y \sim U(0,1)$. Then $T_{n_1,n_2;X,Y,t} \stackrel{\mathrm{d}}{=} T_{n_1,n_2;S_{\mu}}$ for every $n_1,n_2 \geq 1$. Furthermore, as $n_1,n_2 \to \infty$, $\nu_1(T_{n_1,n_2;X,Y,t}) \stackrel{\mathrm{p}}{\longrightarrow} \mu$ and $T_{n_1,n_2;X,Y,t} \stackrel{\mathrm{p}}{\longrightarrow} \Gamma''_{\mu}$.

This yields another canonical construction for every $\mu \in \mathcal{P}$. (We claim only convergence in probability in Corollary 8.4; convergence a.s. holds at least along every increasing subsequence $(n_1(m), n_2(m))$; see [Diaconis and Janson 08, Remark 8.2].)

8.3. Random Addition of Vertices

Definition (1.10) suggests the following construction:

(8.4) Let $T_{n_1,n_2;p_1,p_2}$ be the random bipartite threshold graph with $n_1 + n_2$ vertices obtained as follows: Take n_1 "white" vertices and n_2 "black" vertices, and arrange them in random order. Then join each white vertex with probability p_1 to all earlier black vertices, and join each black vertex with probability p_2 to all earlier white vertices (otherwise, the vertex is joined to no earlier vertex), the decisions being made independently by tossing a biased coin once for each white vertex, and another biased coin once for each black vertex.

For $p_1, p_2 \in [0, 1]$, let μ_{p_1, p_2} be the probability measure in \mathcal{P} with distribution function

$$F_{\mu_{p_1,p_2}}(x) = \begin{cases} \frac{1-p_1}{p_2} x, & 0 \le x < p_2, \\ 1 - \frac{p_1}{1-p_2} (1-x), & p_2 \le x < 1. \end{cases}$$
(8.5)

Hence μ_{p_1,p_2} has density $(1-p_1)/p_2$ on $(0,p_2)$, and $p_1/(1-p_2)$ on $(p_2,1)$; if $p_2=0$ there is also a point mass $1-p_1$ at 0, and if $p_2=1$ there is also a point mass p_1 at 1. It follows from (5.2) that the corresponding subset $S_{p_1,p_2}:=S_{\mu_{p_1,p_2}}$ of









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 $[0,1]^2$ is the quadrilateral with vertices (0,1), $(1-p_1,1-p_2)$, (1,0), and (1,1) (including degenerate cases in which p_1 or p_2 is 0 or 1).

This is an extension of the definitions in Section 6.3; we have $\mu_{p,p} = \mu_p$ and $S_{p,p} = S_p$. Note also that $\mu^{\dagger}_{p_1,p_2} = \mu_{p_2,p_1}$. In particular, $\mu_{p_1,p_2} \in \mathcal{P}_s$ only if $p_1 = p_2$.

Theorem 8.5. As $n_1, n_2 \to \infty$, the degree distributions

$$\nu_1(T_{n_1,n_2;p_1,p_2})$$
 and $\nu_2(T_{n_1,n_2;p_1,p_2})$

converge in probability to μ_{p_1,p_2} and μ_{p_2,p_1} , respectively. Consequently,

$$T_{n_1,n_2;p_1,p_2} \stackrel{\mathbf{p}}{\longrightarrow} \Gamma_{p_1,p_2}'' := \Gamma_{\mu_{p_1,p_2}}'' \in \mathcal{T}_{\infty,\infty}''.$$

Corollary 8.6. If $p_1, p_2 \in [0, 1]$ and $n_1, n_2 \ge 1$, then

$$T_{n_1,n_2;p_1,p_2} \stackrel{\mathrm{d}}{=} T_{n_1,n_2;S_{p_1,p_2}} \stackrel{\mathrm{d}}{=} T_{n_1,n_2;X_1,X_2,0},$$

where X_j has the density $1 - p_j$ on (-1,0) and p_j on (0,1), j = 1, 2.

Note that if $p_1 + p_2 = 1$, then S_{p_1,p_2} is the upper triangle $S_{1/2} := \{(x,y) : x + y \ge 1\}$. Hence the distribution of $T_{n_1,n_2;p_1,p_2}$ does not depend on p_1 as long as $p_2 = 1 - p_1$. In particular, we may then choose $p_1 = 1$ and $p_2 = 0$. In this case, Definition (8.4) simplifies as follows:

(8.6) Let T_{n_1,n_2} be the random bipartite threshold graph with $n_1 + n_2$ vertices obtained as follows: Take n_1 "white" vertices and n_2 "black" vertices, and arrange them in random order. Join every white vertex to every earlier black vertex.

If $p_1 = 1$ and $p_2 = 0$, then further $X_1 \stackrel{\text{d}}{=} U \sim U(0,1)$ and $X_2 \stackrel{\text{d}}{=} U - 1$ in Corollary 8.6. Hence we have found a number of natural constructions that yield the same random bipartite threshold graph.

Corollary 8.7. If $p_1 \in [0, 1]$ and $n_1, n_2 \ge 1$, then

$$T_{n_1,n_2;p_1,1-p_1} \stackrel{\mathrm{d}}{=} T_{n_1,n_2;1,0} = T_{n_1,n_2} \stackrel{\mathrm{d}}{=} T_{n_1,n_2;S_{1/2}} \stackrel{\mathrm{d}}{=} T_{n_1,n_2;U,U,1},$$

with $U \sim U(0,1)$.

We will see in the next subsection that this random bipartite threshold graph is uniformly distributed as an unlabeled bipartite threshold graph.









8.4. Uniform Random Bipartite Threshold Graphs

It is easy to see that for every bipartite threshold graph, if we color the vertices in V_1 white and the vertices in V_2 black, then there is an ordering of the vertices such that a white vertex is joined to every earlier black vertex but not to any later one. (For example, if there are weights as in (1.9), order the vertices according to w_i' and w_j'' , taking the white vertices first in case of a tie.) This yields a one-to-one correspondence between unlabeled bipartite threshold graphs on $n_1 + n_2$ vertices and sequences of n_1 white and n_2 black balls. Consequently, the number of unlabeled bipartite threshold graphs is

$$|\mathcal{T}_{n_1,n_2}| = \binom{n_1 + n_2}{n_1}, \quad n_1, n_2 \ge 1.$$

Moreover, it follows that T_{n_1,n_2} is uniformly distributed in \mathcal{T}_{n_1,n_2} ; hence Corollary 8.7 yields the following:

Theorem 8.8. The random bipartite threshold graphs T_{n_1,n_2} , $T_{n_1,n_2;p_1,1-p_1}$ $(0 \le p_1 \le 1)$, $T_{n_1,n_2;S_{1/2}}$, $T_{n_1,n_2;U,U,1}$ are all uniformly distributed, regarded as unlabeled bipartite threshold graphs.

We have not studied uniform random labeled bipartite threshold graphs.

Spectrum of Threshold Graphs

There is a healthy literature on the eigenvalue distribution of the adjacency matrix for various classes of random graphs. Much of this is focused on the spectral gap (e.g., most k-regular graphs are Ramanujan [Davidoff et al. 03]). See [Jakobson et al. 99] for evidence showing that random k-regular graphs have the same limiting eigenvalue distribution as the Gaussian orthogonal ensemble. The following results show that random threshold graphs give a family of examples with highly controlled limiting spectrum.

There is a tight connection between the degree distribution of a threshold graph and the spectrum of its Laplacian; see [Merris 94, Hammer and Kelmans 96, Merris and Roby 05]. Recall that the Laplacian of a graph G, with V(G) = [n], say, is the $n \times n$ matrix $\mathcal{L} = D - A$, where A is the adjacency matrix of G and D is the diagonal matrix with entries $d_{ii} = d_G(i)$. (Thus \mathcal{L} is symmetric and has row sums 0.) It is easily seen that $\langle \mathcal{L}x, y \rangle = \sum_{ij \in E(G)} (x_i - x_j)(y_i - y_j)$ for $x, y \in \mathbb{R}^n$. The eigenvalues λ_i of \mathcal{L} satisfy $0 \le \lambda_i \le n$, $i = 1, \ldots, n$, and we define the normalized spectral distribution $\nu_L \in \mathcal{P}$ as the empirical distribution of $\{\lambda_i/n\}_{i=1}^n$.









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For a threshold graph, it is easily seen that if we order the vertices as in (1.2) and Section 2.1, then for each i = 2, ..., n the function

$$\varphi_i(j) := \begin{cases} -1, & j < i, \\ i - 1, & j = i, \\ 0, & j > i, \end{cases}$$

is an eigenfunction of \mathcal{L} with eigenvalue d(i) or d(i)+1, depending on whether i is added as isolated or dominant, i.e., whether $\alpha_i=0$ or 1 in the binary code of the graph. Together with $\varphi_1:=1$ (which is an eigenfunction with eigenvalue 0 for any graph), these form an orthogonal basis of eigenfunctions. The Laplacian spectrum thus can be written

$$\{0\} \cup \{d(i) + \alpha_i : i = 2, \dots, n\}.$$
 (9.1)

In particular, the eigenvalues are all integers.

Moreover, (9.1) shows that the spectrum $\{\lambda_i\}_1^n$ is closely related to the degree sequence; in particular, asymptotically they are the same in the sense that if G_n is a sequence of threshold graphs with $v(G_n) \to \infty$ and $\mu \in \mathcal{P}$, then

$$\nu_L(G_n) \to \mu \iff \nu(G_n) \to \mu.$$
 (9.2)

(See [Hammer and Kelmans 96] for a detailed comparison of the Laplacian spectrum and the degree sequence for threshold graphs.) In particular, Theorem 5.8 can be restated using the spectral distribution:

Theorem 9.1. Let G_n be a sequence of threshold graphs such that $v(G_n) \to \infty$. Then G_n converges in $\overline{\mathcal{U}}$ as $n \to \infty$ if and only if the spectral distributions $\nu_L(G_n)$ converge to some distribution μ . In this case, $\mu \in \mathcal{P}_s$ and $G_n \to \Gamma_{\mu}$.

Remark 9.2. It can be shown that the spectrum and the degree sequence are asymptotically close in the sense that (9.2) holds for any graphs G_n with $v(G_n) \to \infty$, even though in general there is no simple relation like (9.1).

Another relation between the spectrum and the degree sequence for a threshold graph is that their Ferrers diagrams are transposes of each other; see [Merris 94, Merris and Roby 05]. This is easily verified from (9.1) by induction. If we scale the Ferrers diagrams by n, so that they fit in the unit square $[0,1]^2$ with a corner at (0,1), then the lower boundary is the graph of the empirical distribution function of the corresponding normalized values, i.e., the distribution function of $\nu(G)$ or $\nu_L(G)$. Hence these distribution functions are related by reflection in









the diagonal between (0,1) and (1,0), so by (5.11) (and the comment after it), for any threshold graph G,

$$\nu_L(G) = \nu(G)^{\dagger}$$
.

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