

THE FUNDAMENTAL LEMMA OF COMPLEXITY FOR ARBITRARY FINITE SEMIGROUPS¹

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1. **Statement of the results and some corollaries.** All semigroups considered are of finite order. In the recent paper [3] and in the recent book [2] the complexity of a semigroup was defined and definitive results were obtained for determining the complexity of a semigroup which was the union of groups. Herein we state generalizations, valid for arbitrary finite semigroups, of those previous results. All undefined notation is explained in [2].

We first recall the definition of complexity. See also [2] or [3]. One semigroup, S_1 , is said to *divide* another semigroup, S_2 , if and only if S_1 is a homomorphic image of a subsemigroup $S \leq S_2$. If S is a semigroup, $\text{Endo}(S)$ denotes the semigroup of endomorphisms of S under composition. If S_1 and S_2 are semigroups and Y is a homomorphism of S_1 into $\text{Endo}(S_2)$, the *semidirect product of S_2 by S_1 with connecting homomorphism Y* , denoted by $S_2 \times_Y S_1$, is the semigroup with elements $S_2 \times S_1$ and product defined by $(s_2, s_1) \cdot (s'_2, s'_1) = (s_2 \cdot Y(s_1)(s'_2), s_1 \cdot s'_1)$.

We can construct new semigroups from old ones by taking semidirect products and then divisors. $S_n \times_{Y_{n-1}} \cdots \times_{Y_2} S_2 \times_{Y_1} S_1$ denotes $(\cdots (S_n \times_{Y_{n-1}} S_{n-1}) \times_{Y_{n-2}} S_{n-2}) \cdots \times_{Y_1} S_1$ where Y_{n-2} is a homomorphism of S_{n-2} into $\text{Endo}(S_n \times_{Y_{n-1}} S_{n-1})$, etc. We say S is a *combinatorial* semigroup if and only if the subsemigroups of S which are groups are singletons. The main theorem of [1] (see also [2, Chapter 5]) implies that for each semigroup S there exist semigroups S_n, \cdots, S_1 and connecting homomorphisms Y_{n-1}, \cdots, Y_1 so that

$$(1.1) \quad S \text{ divides } S_n \times_{Y_{n-1}} \cdots \times_{Y_1} S_1$$

and S_k is either a simple nontrivial group dividing S or S_k is a combinatorial semigroup, for $k=1, \cdots, n$.

$\#_G(S)$, the (group) complexity of S , is by definition the smallest nonnegative integer n such that

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$$(1.2) \quad S \text{ divides } C_n \times_{Y_{n-1}} G_n \times_{Z_{n-1}} C_{n-1} \times_{Y_{n-2}} G_{n-1} \times_{Z_{n-2}} \cdots \times_{Y_0} G_1 \times_{Z_0} C_0$$

with C_n, \dots, C_1, C_0 combinatorial semigroups and G_n, \dots, G_1 nontrivial groups. For extensive background see [2].

Let \mathcal{S}_F denote the collection of all finite semigroups, \mathcal{S} the collection of all finite semigroups which are union of groups and N the non-negative integers. Then $\#_G: \mathcal{S}_F \rightarrow N$. In [3] and [2, Chapter 9], it was proved that $\#_G$ restricted to \mathcal{S} satisfies the following axioms:

AXIOM I. $\#_G(S) = \max \{ \#_G(S_i) : i = 1, \dots, n \}$ if $S \leq S_1 \times \cdots \times S_n$ where \leq denotes subdirect product. See [2].

AXIOM II. (FUNDAMENTAL LEMMA OF COMPLEXITY). Let I be a combinatorial ideal of S . Then

$$(1.3) \quad \#_G(S) = \#_G(S/I). \quad \text{Also } \#_G(\{0\}) = 0.$$

AXIOM III. Let $S \neq \{0\}$ and let S be a group mapping (GM) semigroup with RLM the right letter mapping homomorphic image of S .³ Then

$$(1.4) \quad \#_G(S) = \#_G(\text{RLM}(S)) + 1.$$

We ask which Axioms remain valid for $\#_G: \mathcal{S}_F \rightarrow N$?

It is trivial to verify that Axiom I remains valid for \mathcal{S}_F . It is easy to see that Axiom III is false for \mathcal{S}_F , e.g. the symmetric inverse semigroup on n letters has complexity 1. See [7]. In fact, no function from \mathcal{S}_F into N satisfies all three Axioms. In [2, Corollary 9.3.4], Axiom II is proved to be equivalent to Axiom II'.

AXIOM II'. Let the epimorphism $\theta: S \rightarrow T$ be one-to-one when restricted to each subgroup of S . Then $\#_G(S) = \#_G(T)$.

The epimorphisms of the hypothesis of Axiom II' are called γ -epimorphisms in [2]. Our main result is the following theorem.

THEOREM. *Axiom II, or equivalently, Axiom II' holds for all finite semigroups.*

It is well known (see [2, Proposition 8.2.17(b)]) that if S is a GM semigroup then either $\#_G(S)$ equals $\#_G(\text{RLM}(S)) + 1$ or $\#_G(\text{RLM}(S))$. We say S is a *pure group mapping* (PGM) semigroup if and only if S is a GM semigroup $\neq \{0\}$ and (1.4) holds for S .

³ S is a GM semigroup iff S has a 0-minimal noncombinatorial ideal I so that S acts faithfully on I by right multiplication and also by left multiplication. $\text{RLM}(S)$ is the action made faithful of S by right multiplication on the principal left ideals of I . See [2].

COROLLARY 1. $\#_G(S)$ equals the largest nonnegative integer $n = \#_1(S)$ such that there exists a series

$$(1.5) \quad S \twoheadrightarrow \text{PGM}_1 \twoheadrightarrow \text{RLM}(\text{PGM}_1) \twoheadrightarrow \dots \twoheadrightarrow \text{PGM}_n \twoheadrightarrow \text{RLM}(\text{PGM}_n)$$

where \twoheadrightarrow denotes epimorphism, and PGM_k denotes a PGM semigroup $\neq \{0\}$ for $k = 1, \dots, n$.

PROOF. First $\#_1(S) \leq \#_G(S)$ follows by the definition of PGM. The reverse inequality $\#_G(S) \leq \#_1(S)$ follows from [2, Lemma 8.2.19(b)], Axiom II and the definition of PGM. See the proof of [2, Theorem 9.2.5].

COROLLARY 2.

$$S_{\mathfrak{L}} \twoheadrightarrow T^A \text{ implies } \#G(T) \leq \#G(S) \leq \#G(T) + 1.$$

PROOF. $S \twoheadrightarrow T_{\mathfrak{L}^c}$ the minimal \mathfrak{L}' homomorphic image of S equals $S \twoheadrightarrow S^{\text{RLM}}$ by [2, Fact 8.3.9(c)]. Now apply Corollary 1.

COROLLARY 3. (CONTINUITY OF COMPLEXITY WITH RESPECT TO HOMOMORPHISMS) Let $\theta: S \twoheadrightarrow T$ be an epimorphism, and let $\#_G(S) = n$ and $\#_G(T) = k$. Then there exists epimorphisms $S = S_n \twoheadrightarrow S_{n-1} \twoheadrightarrow \dots \twoheadrightarrow S_k = T$, so that the composite epimorphism is θ , and $\#_G(S_j) = j$ for $j = k, \dots, n$.

PROOF. Apply [2, Theorem 8.1.14], the Theorem and Corollary 2.

COROLLARY 4. $\#_G(S)$ equals the maximum of the $\#_G(S')$ where S' ranges over the $\phi(S)$ where ϕ is an irreducible representation of S into $n \times n$ complex matrices.

PROOF. The direct sum of the ϕ 's give a γ -epimorphism by [6].

2. **Indication of the proof.** Complete details will appear in [4]. Unfortunately they are long and messy. However, we will try to make the philosophy of the proof clear by the following discussion.

Suppose for each semigroup S we can construct another semigroup $\alpha(S)$ such that

$$(2.1) \quad \alpha(S) \twoheadrightarrow S$$

and if I is a combinatorial ideal of S then

⁴ $\theta: S \twoheadrightarrow T$ is an \mathfrak{L} (resp. \mathfrak{L}') epimorphism iff $s_1, s_2 \in S$ (and s_1, s_2 regular elements) and $\theta_1(s_1) = \theta_1(s_2)$ implies $S^1 s_1 = S^1 s_2$. See [2].

$$(2.2) \quad \begin{aligned} &\alpha(S) \text{ divides } C \text{ w } \alpha[(S/I)] \text{ or at least} \\ &C(\alpha(S)) \leq (C, 1) \oplus C(\alpha(S/I)).^5 \end{aligned}$$

Then clearly to prove the Theorem it suffices to prove

$$(*) \quad C(\alpha(S)) \leq (C, 1) \oplus C(S)$$

or equivalently by (2.1)

$$(*) \quad C(\alpha(S)) \approx C(S)$$

where $C(S) \approx C(T)$ iff $C(S) \leq (C, 1) \oplus C(T)$ and $C(T) \leq (C, 1) \oplus C(S)$.

EXAMPLES OF α . Before continuing we list some good examples of α . With reference to [2, §5.4], we suppose that for each S we choose a system of subsemigroups S_n, \dots, S_1 and we let $\alpha(S)$ be the subsemigroup of

$$(S_n^I, (S_n^I) \sim) \text{ w } \dots \text{ w } (S_1^I, (S_1^I) \sim)$$

generated by $\{s: s \in S\}$ defined in the proof of Lemma 5.4.4 of [2]. Clearly (2.1) holds.

(2.3) If S is a union of groups and the system is chosen to be the \mathfrak{g} -classes of S as in Remark 5.4.14 of [2], then $\alpha(S)$ satisfies (2.2), as can be verified. See [3] or Chapter 9 of [2].

(2.4) If I is a combinatorial ideal of S , then the system S_n, \dots, S_1 can be chosen so that either $S_i \cap I = \emptyset$ or S_i is combinatorial and contains I . In this case (2.2) can be verified. See [4] for complete details.

Yet another way to construct α 's is the following.

(2.5) Consider the right regular representation (S^I, S) and apply the method of Zeiger (see [9] and Chapter 4 of [2]). Let $\alpha(S)$ be the subsemigroup of the wreath product of permutation-reset mapping semigroups so obtained which maps homomorphically onto S . Thus (2.1) holds and (2.2) can be verified. See [4] for complete details.

Now we give a method by which (*) can be proved. We first note that if S is a union of groups and α is given by (2.3), then (*) can be verified by brute force using the machine method of [1]. For the details see [3] or Chapter 9 of [2]. The general case seems difficult by direct methods and we proceed indirectly as follows.

⁵ $S_2 \text{ w } S_1$ denotes the wreath product of the right regular representation of S_1 by S_2 , i.e., $S_2 \text{ w } S_1 = (S_2^I, S_2) \wr (S_1^I, S_1)$. Let $n = \#(S)$ be as defined just before (2.10). Then by definition $C(S) = (C, n)$, resp. (G, n) , resp. $(C \vee G, n)$ if S satisfies (2.10) (b) and not (2.10) (a), resp. S satisfies (2.10) (a) and not (2.10) (b), resp. S satisfies both (2.10) (a) and (b). By definition, $(C, 1) \oplus (C, n) = (C, 1) \oplus (C \vee G, n) = (C, 1) \oplus (G, n-1) = (C, n)$. Finally, by definition $(\alpha, v) \leq (\beta, j)$ iff $v \leq j$, or $v=j$ and $\alpha=\beta$, or $v=j$ and $\alpha=C \vee G$. See [2].

Suppose one can show

(2.6), (2.6)' (ENLARGING LEMMA). *If $C(S) \approx C(T)$ and S divides T (resp. $T \twoheadrightarrow S$) and $(*)$ holds for T , then $(*)$ holds for S .*

(2.7) Let $\theta: S \twoheadrightarrow T$ be a $\gamma(\mathfrak{C})$ -epimorphism.⁶ Then $\alpha(S)$ divides C w $\alpha(T)$ with C combinatorial or at least

$$C(\alpha(S)) \leq (C, 1) \oplus C(\alpha(T)).$$

(2.8), (2.8)' Suppose $\theta: S \twoheadrightarrow T$ and θ is an \mathcal{L} (resp. \mathcal{L}') epimorphism. Then $\alpha(S)$ divides C_1 w G w C_2 w $\alpha(T)$ or at least $C(\alpha(S)) \leq (C, 3) \oplus C(\alpha(T))$.

Then

LEMMA (2.9). (2.1), (2.6)–(2.8) or (2.1), (2.6)', (2.7) and (2.8)' imply $(*)$.

PROOF. Suppose (2.9) is false and let S be a counter-example whose complexity number (defined next) $\#(S) = n$ is as small as possible. By the definition of complexity number $\#(S)$ either

$$(2.10)(a) \quad S \text{ divides } G_n \text{ w } C_{n-1} \text{ w } G_{n-2} \text{ w } C_{n-2} \text{ w } \dots = W$$

or

$$(2.10)(b) \quad S \text{ divides } C_n \text{ w } G_{n-1} \text{ w } C_{n-2} \text{ w } G_{n-2} \text{ w } \dots = W$$

where G_j 's are groups and the C_u 's are combinatorial monoids and for no smaller n is (2.10) (a) or (b) true. But $C(S) \approx C(W)$ so (2.6) implies $(*)$ is false for W . But in Case (2.10) (b)

$$(2.11)(b) \quad W \xrightarrow[\gamma(\mathfrak{C})]{} W_{-1} = p_{-1}(W)$$

where p_{-1} is the projection onto the first $n-1$ coordinates. In case (2.10) (a)

$$(2.11)(a) \quad W \xrightarrow{\mathcal{L}} W_{-1} = p_{-1}(W)$$

and in either case $\#(W_{-1}) = \#(W) - 1 = \#(S) - 1$. Thus by induction $(*)$ holds for W_{-1} and we have

$$(2.12)(a) \quad C(W_{-1}) = (C, n - 1)$$

$$(2.12)(b) \quad C(W_{-1}) = (G, n - 1)$$

respectively. But then (2.11), (2.12) and (2.7) and (2.8) implies $(*)$ holds for W , a contradiction. The other case with (2.6)', etc. proceeds similarly. This proves (2.9).

⁶ θ restricted to each \mathfrak{C} -class of S is one-to-one.

SKETCH OF THE PROOF OF THE THEOREM. Using α of (2.4) which we denote by α^* we verify (2.1) and (2.2) and further show that

$$\alpha^*(S) \xrightarrow{\gamma(\mathcal{C})} S.$$

We do not verify (2.6)–(2.8) directly for α^* of (2.4).

Then using α of (2.5) which we denote by Z , we verify (2.1) but not (2.2) for Z because I can contain large nonregular \mathcal{C} -classes of S . However, we can verify (2.6), (2.7) and (2.8) for Z by using the classification of maximal proper epimorphisms proved in [5]. Then Lemma (2.9) implies (*) for $Z(S)$. Then (*) for Z and (2.7) implies

$$C(S) \approx C(T) \quad \text{if } S \xrightarrow{\gamma(\mathcal{C})} T.$$

But from the first paragraph

$$\alpha^*(S) \xrightarrow{\gamma(\mathcal{C})} S$$

so (*) holds for α^* , so (2.1), (2.2) and (*) holds for α^* and the Theorem follows.

For further results on complexity see [7].

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