

ON UNIT-ZERO BOOLEAN REPRESENTATIONS
OF OPERATIONS AND RELATIONS*

BY B. A. BERNSTEIN

1. *Introduction.* Consider an algebra $(K, +, \times)$, such as ordinary real algebra, in which there are two elements "0" and "1" having the properties that, for any element a ,

$$(1) \quad a + 0 = 0 + a = a, a1 = 1a = a.$$

Let

$$(x_1, x_2, \dots, x_m; a_1, a_2, \dots, a_m)$$

denote a *unit-zero function with respect to the sequence of m elements, a_1, \dots, a_m of K* , that is, a function $f(x_1, x_2, \dots, x_m)$ of m elements x_1, x_2, \dots, x_m such that $f = 1$ or 0 , according as the equalities, $x_i = a_i$, ($i = 1, 2, \dots, m$), all hold or do not all hold. Accordingly, $(x; a)$ will denote a *unit-zero function with respect to a* , that is, a function $f(x)$ such that $f(x) = 1$ or 0 , according as $x = a$ or $x \neq a$. Then the following propositions (2)–(4) evidently hold:

$$(2) \quad (x_1, x_2, \dots, x_m; a_1, a_2, \dots, a_m) = (x_1; a_1)(x_2; a_2) \dots (x_m; a_m);$$

$$(3) \quad a(x_1, x_2, \dots, x_m; a_1, a_2, \dots, a_m) = a \text{ or } 0,$$

according as $x_i = a_i$, ($i = 1, 2, \dots, m$), all hold or do not all hold;

$$(4) \quad a(x_1, x_2, \dots, x_m; a_1, a_2, \dots, a_m) \\ + b(x_1, x_2, \dots, x_m; b_1, b_2, \dots, b_m) = a, \text{ or } b, \text{ or } 0,$$

according as $x_i = a_i$ all hold, or $x_i = b_i$ all hold, or neither $x_i = a_i$ all hold nor $x_i = b_i$ all hold, ($i = 1, 2, \dots, m; a_i \neq b_i$ for some i). In a previous paper† propositions (1)–(4) were made the basis of a method of obtaining arithmetic representations of arbitrary operations and relations in a finite class of elements. Since

* Presented to the Society, April 11, 1931.

† B. A. Bernstein and N. Debely, *A practical method for the modular representation of finite operations and relations*, this Bulletin, vol. 38 (1932), pp. 110–114.

propositions (1)–(4) also hold when the symbols belong to Boolean algebra, the question naturally arises: To what extent can unit-zero functions be used analogously to obtain *Boolean* representations of arbitrary operations and relations? The object of the present paper is to answer this question.

2. *Determination of Boolean Unit-Zero Algebras.* The possibility of representing arbitrary operations and relations by unit-zero functions of an algebra hinges on the existence in this algebra of a unit-zero function for every sequence of m of its elements. Let us call an algebra which has a unit-zero function for every sequence of m of its elements a *unit-zero algebra*. I proceed first to determine all Boolean unit-zero algebras.

This determination is made easy by noting at the outset that a unit-zero Boolean function must satisfy proposition (2) above and also that it must be single-valued. We therefore need to look only for Boolean unit-zero functions $f(x)$ of a *single* variable x of the form*

$$(5) \quad (x; a) = (1; a)x + (0; a)x'.$$

From (5) we see, by putting $a=0, 1$, that in a Boolean algebra of *two* elements, x is the unit-zero function of x with respect to 1, and x' is the unit-zero function of x with respect to 0; in symbols,

$$(6) \quad (x; 1) = x, \quad (x; 0) = x'.$$

We have, then, that *a two-element Boolean algebra is a unit-zero algebra, the unit-zero functions of one variable x being given by (6).*

By (2) and (6), *all* the unit-zero functions of a two-element Boolean algebra can be readily written down. Thus, the unit-zero functions of two variables x, y are given by

$$(7) \quad \begin{aligned} (x, y; 1, 1) &= xy, & (x, y; 1, 0) &= xy', \\ (x, y; 0, 1) &= x'y, & (x, y; 0, 0) &= x'y'. \end{aligned}$$

In general, *the unit-zero functions of m variables are the 2^m constituents in the normal development of 1 with respect to the m variables.*

* The usual Boolean notations are employed: $a+b$, ab , a' , 0, 1 are respectively the *sum* of a and b , the *product* of a and b , the *negative* of a , the *zero* element, the *whole*.

Let us now consider a Boolean algebra A of more than two elements. A must have an element $e \neq 0, 1$. Suppose, first, that A has a unit-zero function $f(x)$, of form (5), with respect to e . Then

$$(i) \quad f(e) = 1, f(0) = 0, f(1) = 0, \quad (e \neq 0, 1).$$

But (i) is inconsistent with (5). Hence, *our algebra A has no unit-zero function with respect to a sequence containing the element e .*

Suppose, next, that the algebra A has a unit-zero function $f(x)$, of form (5) with respect to 0. Then

$$(ii) \quad f(0) = 1, f(1) = 0, f(e) = 0, \quad (e \neq 0, 1).$$

Hence, by (5),

$$(iii) \quad f(x) = x', f(e) = 0, \quad (e \neq 0, 1).$$

But equations (iii) are inconsistent. Hence, *our algebra A has no unit-zero function with respect to a sequence containing the element 0.*

Similarly, *our algebra A has no unit-zero function with respect to a sequence containing the element 1.* Hence, *a Boolean algebra of more than two elements has no unit-zero functions at all.*

Our main result is, then, the following theorem.

THEOREM A. *The only Boolean unit-zero algebra is a two-element Boolean algebra.*

3. *Dual Considerations.* By the Principle of Duality in Boolean algebras each of the foregoing propositions about unit-zero Boolean functions has a dual proposition corresponding to it. To state these duals, let me use the notion of *zero-unit function* (to be distinguished from *unit-zero function*). By a *zero-unit function of x_1, x_2, \dots, x_m with respect to the sequence a_1, a_2, \dots, a_m* , symbolized by

$$[x_1, x_2; \dots, x_m; a_1, a_2, \dots, a_m],$$

let us mean a function $f(x_1, x_2, \dots, x_m)$ such that $f=0$ or 1 , according as $x_i = a_i$, ($i = 1, 2, \dots, m$), all hold or do not all hold. The duals of (2), (3), and (4) are, then, respectively (2'), (3'), and (4') following:

$$(2') \quad [x_1, x_2, \dots, x_m; a_1, a_2, \dots, a_m] \\ = [x_1; a_1] + [x_2; a_2] + \dots + [x_m; a_m];$$

$$(3') \quad a + [x_1, x_2, \dots, x_m; a_1, a_2, \dots, a_m] = a \text{ or } 1,$$

according as $x_i = a_i$, ($i = 1, 2, \dots, m$), all hold or do not all hold;

$$(4') \quad \{a + [x_1, x_2, \dots, x_m; a_1, a_2, \dots, a_m]\} \\ \cdot \{b + [x_1, x_2, \dots, x_m; b_1, b_2, \dots, b_m]\} = a, \text{ or } b, \text{ or } 1,$$

according as $x_i = a_i$ all hold, or $x_i = b_i$ all hold, or neither $x_i = a_i$ all hold nor $x_i = b_i$ all hold, ($i = 1, 2, \dots, m$; $a_i \neq b_i$ for some i).

The dual of Theorem A is

THEOREM A'. *The only zero-unit Boolean algebra is a two-element Boolean algebra.*

For a *two-element* Boolean algebra we have, further:

$$(6') \quad [x; 0] = x, \quad [x; 1] = x';$$

$$(7') \quad [x, y; 0, 0] = x + y, \quad [x, y; 0, 1] = x + y',$$

$$[x, y; 1, 0] = x' + y, \quad [x, y; 1, 1] = x' + y'.$$

In general, *the zero-unit functions of m variables are the 2^m factor-constituents in the dual normal development of 0 with respect to the m variables.*

Propositions (2')–(7') will be used below in the representation of operations that do not satisfy the condition of closure.

4. *Representations.* It is now clear to what extent we can apply unit-zero Boolean functions in the representation of arbitrary operations and relations. From Theorem A, we have

THEOREM B. *A unit-zero Boolean representation of arbitrary operations and relations is possible when and only when the class consists of two elements.*

For a two-element class K , the theory of Boolean representation follows from propositions (2)–(7) and their duals. If we denote the two K -elements by the Boolean symbols 0, 1, the representations of all operations O and relations R in K are covered by the cases 1–3 following.

CASE 1. *O an m -ary operation satisfying the condition of closure.* There is a K -element, 0 or 1, for every sequence e_1, e_2, \dots, e_m taken from K . Let the sequences to which 1 corresponds be

(i) $\alpha_{11}, \alpha_{12}, \dots, \alpha_{1m}; \alpha_{21}, \alpha_{22}, \dots; \alpha_{2m}; \dots; \alpha_{k1}, \alpha_{k2}, \dots, \alpha_{km}$.

The representation of O is the Boolean function

$$(8) \quad \sum_{i=1}^k (x_1, x_2, \dots, x_m; \alpha_{i1}, \alpha_{i2}, \dots, \alpha_{im}).$$

CASE 2. O an m -ary operation not satisfying the closure condition. There are sequences in K to which no K -elements correspond. Let these sequences be

(ii) $\beta_{11}, \beta_{12}, \dots, \beta_{1m}; \beta_{21}, \beta_{22}, \dots, \beta_{2m}; \dots; \beta_{k1}, \beta_{k2}, \dots, \beta_{km}$.

Consider the operation O' obtained from O by assigning a K -element, 0 for convenience, to each of the sequences (ii). Let $\phi(x_1, x_2, \dots, x_m)$, obtained as in Case 1, be the representation of O' . Then the representation of O is the function

$$(9) \quad \phi(x_1, x_2, \dots, x_m) + \sum_{i=1}^k 0/[x_1, x_2, \dots, x_m; \beta_{i1}, \beta_{i2}, \dots, \beta_{im}],$$

where a/b means the unique K -element q satisfying the condition $bq = a$.*

CASE 3. R an m -adic relation. Let the sequences which do not satisfy R be

(iii) $\gamma_{11}, \gamma_{12}, \dots, \gamma_{1m}; \gamma_{21}, \gamma_{22}, \dots, \gamma_{2m}; \dots; \gamma_{k1}, \gamma_{k2}, \dots, \gamma_{km}$.

Then the representation of R is the Boolean equation

$$(10) \quad \sum_{i=1}^k (x_1, x_2, \dots, x_m; \gamma_{i1}, \gamma_{i2}, \dots, \gamma_{im}) = 0. \dagger$$

Of course, by the Duality Principle, the theory of representation can be stated primarily in terms of zero-unit functions instead of unit-zero functions.

5. *Illustrations.* The following illustrations, one for each of the above three cases, will make the theory of representation quite clear.

α . Let O be the operation defined by

* For a two-element Boolean algebra the *quotient* can be defined precisely as in ordinary algebra.

† Instead of 0, we can use 1 in (10), provided (i) are the sequences which do satisfy R .

$$(i) \quad \begin{array}{c|cc} & 0 & 1 \\ \hline 0 & 1 & 0 \\ 1 & 0 & 1 \end{array}$$

Its representation is

$$(ii) \quad (x, y; 0, 0) + (x, y; 1, 1) \equiv x'y' + xy.$$

β . Let O be the operation

$$(iii) \quad \begin{array}{c|cc} & 0 & 1 \\ \hline 0 & 0 & 1 \\ 1 & - & - \end{array},$$

where the blanks indicate that there are no K -elements corresponding to the sequences 1, 0; 1, 1.

Consider the operation O' defined by

$$(iv) \quad \begin{array}{c|cc} & 0 & 1 \\ \hline 0 & 0 & 1 \\ 1 & 0 & 0 \end{array}$$

By Case 1, the representation of O' is

$$(v) \quad x'y.$$

Hence, the representation of O is

$$(vi) \quad x'y + 0/[x, y; 1, 0] + 0/[x, y; 1, 1] \\ \equiv x'y + 0/(x' + y) + 0/(x' + y').$$

γ . Let R be a relation defined by

$$(vii) \quad \begin{array}{c|cc} & 0 & 1 \\ \hline 0 & - & + \\ 1 & + & - \end{array},$$

where “+” indicates that R holds and “-” indicates that R does not hold. Its representation is the equation

$$(viii) \quad (x, y; 0, 0) + (x, y; 1, 1) \equiv x'y' + xy = 0.*$$

THE UNIVERSITY OF CALIFORNIA

* For a complete set of Boolean representations of binary operations and dyadic relations in a two-element class, obtained from considerations other than the above, see my *Complete sets of representations of two-element algebras*, this Bulletin, vol. 30 (1924), pp. 24-30.