

### “Hasse principle” for $GL_2(D)$

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**1. Statement of a theorem.** Let  $D$  be a Euclidean domain and  $G = GL_2(D)$ , the group of invertible  $2 \times 2$  matrices over  $D$ .<sup>0)</sup> We shall prove that

(1.1) **Theorem.**  $\text{III}(G) = 1$ , i.e.,  $G$  enjoys the “Hasse principle”.<sup>1)</sup>

(1.2) **Remark.** Thanks to an excellent idea of M. Mazur, to prove (1.1) it is enough to verify that

$$(1.3) \quad \text{End}_c(G) = \text{Inn}(G),$$

where the left hand side is the set of all endomorphisms of  $G$  preserving conjugacy classes of  $G$ .<sup>1)</sup> Thus for each  $F \in \text{End}_c(G)$ , and  $A \in G$ , we have

$$(1.4) \quad F(A) \sim A, \quad \text{i.e.,} \quad F(A) = PAP^{-1},$$

$P$  depending on  $A$ .

Given an  $F \in \text{End}_c(G)$  we connect two elements  $A, B$  of  $G$  by a string according to the rule:

$$(1.5) \quad A - B \iff \exists P \in G \quad \text{so that}$$

$$F(A) = PAP^{-1} \quad \text{and} \quad F(B) = PBP^{-1}.$$

Note that  $A - B$  is not, a priori, an equivalence relation defined on  $G$ .<sup>2)</sup> Even so, this relation is very useful to prove the Hasse principle  $\text{III}(G) = 1$ . Note also that the relation (1.5) depends only on  $F$  modulo  $\text{Inn}(G)$ .

**2. Generators for  $G$ .** Before proving (1.1), let us gather some basic facts on  $G = GL_2(D)$ ,  $D$  being a Euclidean domain. Denote by  $D^*$  the group of invertible elements of  $D$ . Let  $N, M_\lambda$  ( $\lambda \in D, \lambda \neq 0$ ),  $D_\mu$  ( $\mu \in D^*, \mu \neq 1$ ) be elements of  $G$  defined by

$$(2.1) \quad N = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \quad M_\lambda = \begin{pmatrix} 1 & \lambda \\ 0 & 1 \end{pmatrix}, \quad D_\mu = \begin{pmatrix} 1 & 0 \\ 0 & \mu \end{pmatrix}.$$

<sup>0)</sup> Needless to say,  $D$  may be any commutative field.

<sup>1)</sup> As for unexplained notation and facts in this paper, see [1].

<sup>2)</sup> This reminds me somehow a children’s string game CAT’S CRADLE, or AYATORI in Japanese. One can play this game on any group  $G$  once an endomorphism  $F \in \text{End}_c(G)$  is chosen.

It is well-known and easy to prove that

$$(2.2) \quad G \text{ is generated by } N, M_\lambda, D_\mu:$$

$$G = \langle N, M_\lambda, D_\mu \rangle,$$

We will use repeatedly the following equalities on

$$P = \begin{pmatrix} x & y \\ z & t \end{pmatrix} \in G.$$

$$(2.3) \quad PNP^{-1} = (\det P)^{-1} \begin{pmatrix} yt - xz & x^2 - y^2 \\ t^2 - z^2 & xz - yt \end{pmatrix},$$

$$(2.4) \quad M_\lambda PNP^{-1} = (\det P)^{-1} \times \begin{pmatrix} yt - xz + \lambda(t^2 - z^2) & x^2 - y^2 + \lambda(xz - yt) \\ t^2 - z^2 & xz - yt \end{pmatrix},$$

$$(2.5) \quad D_\mu PNP^{-1} = (\det P)^{-1} \begin{pmatrix} yt - xz & x^2 - y^2 \\ \mu(t^2 - z^2) & \mu(xz - yt) \end{pmatrix}.$$

### 3. Proof of the theorem.

**Step (I).** To prove that  $N - M_\lambda$ . Since we can adjust a given  $F$  in  $\text{End}_c(G)$  by elements of  $\text{Inn}(G)$ , We may assume that

$$(3.1) \quad \begin{cases} F(M_\lambda) = M_\lambda, \\ F(N) = PNP^{-1}, \quad P \in G. \end{cases}$$

Our problem is to find  $P_0 \in G$  so that

$$(3.2) \quad \begin{cases} F(M_\lambda) = M_\lambda = P_0 M_\lambda P_0^{-1}, \\ F(N) = PNP^{-1} = P_0 N P_0^{-1}. \end{cases}$$

Put

$$(3.3) \quad P = \begin{pmatrix} x & y \\ z & t \end{pmatrix}, \quad P_0 = \begin{pmatrix} 1 & y_0 \\ 0 & 1 \end{pmatrix}.$$

Clearly  $P_0$ , with any  $y_0 \in D$ , meets the first equality of (3.2). As for the second equality of (3.2), in view of (2.3) for  $P$  and  $P_0$  we are forced to set  $y_0 = (yt - xz)/(xt - yz)$  and then we should verify the equality (3.2) which boils down to a single equality:

$$(3.4) \quad \det(P) = xt - yz = t^2 - z^2$$

as a little calculation shows. To get (3.4), we must use seriously the assumption that  $F$  is a homomor-

phism:  $F(M_\lambda N) = F(M_\lambda)F(N)$ . In other words, we have

$$(3.5) \quad Q(M_\lambda N)Q^{-1} = M_\lambda PNP^{-1},$$

with some  $Q \in G$ .

Take the trace of both sides of (3.5) and use (2.4). Then (3.4) follows miraculously.

**Step (II).** To Prove that  $N \sim D_\mu$ . As in Step (I), we may assume that

$$(3.6) \quad \begin{cases} F(D_\mu) = D_\mu, \\ F(N) = PNP^{-1}, \quad P \in G. \end{cases}$$

Again we must find  $P_0 \in G$  so that

$$(3.7) \quad \begin{cases} F(D_\mu) = D_\mu = P_0 D_\mu P_0^{-1} \\ F(N) = PNP^{-1} = P_0 N P_0^{-1}. \end{cases}$$

Put

$$(3.8) \quad P = \begin{pmatrix} x & y \\ z & t \end{pmatrix}, \quad P_0 = \begin{pmatrix} x_0 & 0 \\ 0 & 1 \end{pmatrix}.$$

Clearly  $P_0$ , with any  $x_0 \in D^*$ , meets the first equality of (3.7). As for the second equality of (3.7), in view of (2.3) for  $p$  and  $p_0$  we are forced to set  $x_0 = (x^2 - y^2)/(xt - yz)$  and then we should verify the equality (3.7) which boils down to a single equality:

$$(3.9) \quad xz - yt = 0$$

as a little calculation shows. To get (3.9), we must use again the property  $F(D_\mu N) = F(D_\mu)F(N)$ . We have then

$$(3.10) \quad Q(D_\mu N)Q^{-1} = D_\mu PNP^{-1}, \quad Q \in G.$$

Take the trace of both sides of (3.10) and use (2.5). Then (3.9) follows again.

**Step (III).** Combining (I), (II), we found, for any  $F \in \text{End}_c(G)$ ,  $\lambda(\neq 0) \in D$ ,  $\mu(\neq 1) \in D^*$ , matrices  $P, Q$  in  $G$  so that

$$(3.11) \quad \begin{cases} F(N) = PNP^{-1} = QNQ^{-1}, \\ F(M_\lambda) = PM_\lambda P^{-1}, \\ F(D_\mu) = QD_\mu Q^{-1}. \end{cases}$$

Here it is important to note that the matrix  $Q$  above does not depend on  $\mu$  up to scalars. In fact, let  $X = Q_\mu^{-1}Q_{\mu'}$ , with  $F(N) = Q_\mu N Q_\mu^{-1} = Q_{\mu'} N Q_{\mu'}^{-1}$ ,  $F(D_\mu) = Q_\mu D_\mu Q_\mu^{-1}$ ,  $F(D_{\mu'}) = Q_{\mu'} D_{\mu'} Q_{\mu'}^{-1}$ . Then

$$X = \begin{pmatrix} x & y \\ y & x \end{pmatrix}, \quad x^2 - y^2 \in D^*.$$

Then by comparing the traces of both sides of  $F(D_{\mu\mu'}) = F(D_\mu)F(D_{\mu'})$ , one verifies that  $X = xI$ , and so one can assume that  $Q_\mu = Q_{\mu'} = Q$ .

From the first line of (3.11), we infer that

$$(3.12) \quad R = Q^{-1}P = \begin{pmatrix} a & b \\ b & a \end{pmatrix}.$$

Adjusting  $F$  modulo  $\text{Inn}(G)$ , the last two lines of (3.11) imply that

$$(3.13) \quad \begin{cases} F(M_\lambda) = RM_\lambda R^{-1}, \\ F(D_\mu) = D_\mu. \end{cases}$$

Taking the traces of both sides of  $F(D_\mu M_\lambda) = F(D_\mu)F(M_\lambda)$ , we get, after a little calculation using  $\lambda \neq 0$ ,  $\mu \neq 1$  and  $a^2 - b^2 \neq 0$ ,

$$(3.14) \quad ab = 0.$$

In other words,

$$(3.15) \quad R = aI \quad \text{or} \quad bN, \quad I = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}.$$

So from (3.11), (3.13), (3.15), we may assume that either

$$(3.16) \quad \begin{cases} F(N) = N = {}^t N, \\ F(M_\lambda) = NM_\lambda N^{-1} = {}^t M_\lambda, \quad a = 0, \\ F(D_\mu) = D_\mu = {}^t D_\mu. \end{cases}$$

or

$$(3.17) \quad \begin{cases} F(N) = N, \\ F(M_\lambda) = M_\lambda, \quad b = 0, \\ F(D_\mu) = D_\mu. \end{cases}$$

If (3.16) was the case, we would have  $NM_\lambda D_\mu \sim F(NM_\lambda D_\mu) = F(N)F(M_\lambda)F(D_\mu) = N {}^t M_\lambda D_\mu = {}^t (D_\mu M_\lambda N)$  and, on taking the traces, we get  $\lambda\mu = \lambda$ , or  $\lambda(\mu - 1) = 0$ , contradicting our assumption on  $\lambda, \mu$ . So (3.17) shows that our original  $F$  is an inner automorphism, i.e., the Hasse principle  $\text{III}(G) = 1$  holds. □

### References

- [ 1 ] T. Ono: "Shafarevich-Tate sets" for profinite groups. Proc. Japan Acad., **75A**, 97-98 (1999).
- [ 2 ] H. Wada: "Hasse principle" for  $SL_n(D)$ . Proc. Japan Acad., **75A**, 67-69 (1999).