

Algorithms for b -Functions, Induced Systems, and Algebraic Local Cohomology of D -Modules

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1. Introduction. Let K be an algebraically closed field of characteristic zero and let X be a Zariski open set of K^n with a positive integer n . We fix a coordinate system $x = (x_1, \dots, x_n)$ of X and write $\partial = (\partial_1, \dots, \partial_n)$ with $\partial_i := \partial / \partial x_i$. We denote by \mathcal{D}_X the sheaf of algebraic differential operators on X (cf. [2], [3]).

We assume that (a presentation of) a coherent left \mathcal{D}_X -module \mathcal{M} is given. Let u be a section of \mathcal{M} and let $f = f(x)$ be an arbitrary polynomial of n variables. Let s be an indeterminate. If \mathcal{M} is holonomic, then for each point p of $Y := \{x \in X \mid f(x) = 0\}$, there exist a germ $P(x, \partial, s)$ of $\mathcal{D}_X[s]$ at p and a polynomial $b(s) \in K[s]$ of one variable so that

$$(1.1) \quad P(x, \partial, s)(f^{s+1}u) = b(s)f^s u$$

holds (cf. [11]). More precisely, (1.1) means that there exists a nonnegative integer m so that

$$Q := f^{m-s}(b(s) - P(x, \partial, s)f)f^s \in \mathcal{D}_X[s]$$

satisfies $Qu = 0$ in $\mathcal{M}[s] := K[s] \otimes_K \mathcal{M}$. A monic polynomial $b(s)$ of the least degree that satisfies (1.1) is called the (generalized) b -function for f and u . When \mathcal{M} coincides with the sheaf \mathcal{O}_X of regular functions and $u = 1$, we get the classical b -function (or the Bernstein-Sato polynomial) of f . Algorithms for computing the Bernstein-Sato polynomial have been given by several authors ([21], [25], [4], [16]) but not for an arbitrary f .

One of the main purposes of the present paper is to give algorithms for computing the b -function for u and f and for computing the algebraic local cohomology groups $\mathcal{H}_{[Y]}^j(\mathcal{M})$ ($j = 0, 1$) as left \mathcal{D}_X -modules (cf. [11] for the definition). The algorithm for the local cohomology groups needs some information on the b -function.

These algorithms are actually obtained as byproducts of the solution of more general problems as follows:

Let \mathcal{M} be a left coherent $\mathcal{D}_{K \times X}$ -module. For the sake of simplicity, let us assume here that a

section u of \mathcal{M} generates \mathcal{M} . We identify X with the subset $\{(t, x) \in K \times X \mid t = 0\}$ of $K \times X$. Then the b -function of u along X at $p \in X$ is a nonzero polynomial $b(s) \in K[s]$ of the least degree that satisfies

$$(b(t\partial_t) + tP(t, x, t\partial_t, \partial))u = 0$$

with a germ $P(t, x, t\partial_t, \partial)$ of $\mathcal{D}_{K \times X}$ at p , where we write $\partial_t := \partial / \partial t$. \mathcal{M} is called *specializable* along X at p if such $b(s)$ exists.

We first present an algorithm which computes $b(s)$, or determines that there is none, by using a kind of Gröbner basis for the Weyl algebra related to a filtration introduced by Kashiwara [12]. Such Gröbner bases were used in [18], [19], [20].

If \mathcal{M} is specializable, then its induced system to X is the complex of left \mathcal{D}_X -modules \mathcal{M}_X^\bullet whose cohomology groups are coherent \mathcal{D}_X -modules. We also obtain an algorithm of computing the cohomology groups of \mathcal{M}_X^\bullet by using an FW-Gröbner basis. These algorithms for the b -function and the induced system, combined with a viewpoint of Malgrange [17], provide algorithms for the b -function for a polynomial (and a section of a holonomic system), and for the algebraic local cohomology groups.

When K coincides with the field \mathbf{C} of complex numbers, we can consider the problems explained so far with \mathcal{D}_X replaced by the sheaf $\mathcal{D}_X^{\text{an}}$ of *analytic* differential operators. Then our algorithms yield correct solutions also in this analytic case if the left $\mathcal{D}_X^{\text{an}}$ -module \mathcal{M}^{an} in question is written in the form $\mathcal{M}^{\text{an}} = \mathcal{D}_X^{\text{an}} \otimes_{\mathcal{D}_X} \mathcal{M}$ with a coherent \mathcal{D}_X -module \mathcal{M} whose presentation is given explicitly.

We have implemented the algorithms by using a computer algebra system Kan [24]. Details of the present paper will appear elsewhere.

2. Gröbner bases. Let us denote by A_n and by A_{n+1} the Weyl algebra on n variables x , and the Weyl algebra on $n + 1$ variables (t, x) re-

spectively with coefficients in K . Let r be a positive integer and put $L := \mathbf{N}^{2+2n} = \mathbf{N} \times \mathbf{N} \times \mathbf{N}^n \times \mathbf{N}^n$ with $\mathbf{N} := \{0, 1, 2, \dots\}$. An element P of $(A_{n+1})^r$ is written in a finite sum

$$(2.1) \quad P = \sum_{i=1}^r \sum_{(\mu, \nu, \alpha, \beta) \in L} a_{\mu\nu\alpha\beta i} t^\mu x^\alpha \partial_t^\nu \partial^\beta e_i$$

with $a_{\mu\nu\alpha\beta i} \in K$, $e_1 := (1, 0, \dots, 0), \dots, e_r := (0, \dots, 0, 1)$, $x^\alpha := x_1^{\alpha_1} \cdots x_n^{\alpha_n}$, $\partial^\beta := \partial_1^{\beta_1} \cdots \partial_n^{\beta_n}$ for $\alpha = (\alpha_1, \dots, \alpha_n)$, $\beta = (\beta_1, \dots, \beta_n) \in \mathbf{N}^n$.

For each integer m , we set

$$F_m((A_{n+1})^r) := \{P = \sum_{i=1}^r \sum_{(\mu, \nu, \alpha, \beta) \in L} a_{\mu\nu\alpha\beta i} t^\mu x^\alpha \partial_t^\nu \partial^\beta e_i \mid a_{\mu\nu\alpha\beta i} = 0 \text{ if } \nu - \mu > m\}.$$

Then $\{F_m((A_{n+1})^r)\}_{m \in \mathbf{Z}}$ constitutes a filtration of $(A_{n+1})^r$. For a nonzero element P of $(A_{n+1})^r$, the F -order $\text{ord}_F(P)$ of P is defined as the least $m \in \mathbf{Z}$ such that $P \in F_m((A_{n+1})^r)$.

Let $<_F$ be a total order on $L \times \{1, \dots, r\}$ which satisfies

- (A-1) $(\alpha, i) <_F (\beta, j)$ implies $(\alpha + \gamma, i) <_F (\beta + \gamma, j)$ for any $\alpha, \beta, \gamma \in L$ and $i, j \in \{1, \dots, r\}$;
- (A-2) if $\nu - \mu < \nu' - \mu'$, then $(\mu, \nu, \alpha, \beta, i) <_F (\mu', \nu', \alpha', \beta', j)$ for any $\alpha, \beta, \alpha', \beta' \in \mathbf{N}^n$, $\mu, \nu, \mu', \nu' \in \mathbf{N}$ and any $i, j \in \{1, \dots, r\}$;
- (A-3) $(\mu, \mu, \alpha, \beta, i) \geq_F (0, 0, 0, 0, i)$ for any $\mu \in \mathbf{N}$, $\alpha, \beta \in \mathbf{N}^n$, $i \in \{1, \dots, r\}$.

Let P be a nonzero element of $(A_{n+1})^r$ which is written in the form (2.1). Then the *leading exponent* $\text{lexp}_F(P) \in L \times \{1, \dots, r\}$ of P with respect to $<_F$ is defined as the maximum element

$$\max\{(\mu, \nu, \alpha, \beta, i) \mid a_{\mu\nu\alpha\beta i} \neq 0\}$$

with respect to the order $<_F$. The set of leading exponents $E_F(N)$ of a subset N of $(A_{n+1})^r$ is defined by

$$E_F(N) := \{\text{lexp}_F(P) \mid P \in N \setminus \{0\}\}.$$

Definition 2.1. A finite set \mathbf{G} of generators of a left A_{n+1} -submodule N of $(A_{n+1})^r$ is called an *FW-Gröbner basis* of N if we have

$$E_F(N) = \bigcup_{P \in \mathbf{G}} (\text{lexp}_F(P) + L),$$

where we write

$$(\alpha, i) + L = \{(\alpha + \beta, i) \mid \beta \in L\}$$

for $\alpha \in L$ and $i \in \{1, \dots, r\}$.

Since the order $<_F$ is not a well-order, the Buchberger algorithm ([5], [9], [6], [22]) for computing Gröbner bases does not work directly. In order to bypass this difficulty to obtain an algorithm of computing FW-Gröbner bases, we use

the homogenization technique.

Definition 2.2. For $\lambda, \mu, \nu, \lambda', \mu', \nu' \in \mathbf{N}$ and $\alpha, \beta, \alpha', \beta' \in \mathbf{N}^n$, an order $<_H$ on $L_1 \times \{1, \dots, r\}$ with $L_1 := \mathbf{N} \times L$ is defined so that we have $(\lambda, \mu, \nu, \alpha, \beta, i) <_H (\lambda', \mu', \nu', \alpha', \beta', j)$ if and only if one of the following conditions holds:

- (1) $\lambda < \lambda'$;
- (2) $\lambda = \lambda'$, $(\mu + l, \nu, \alpha, \beta, i) <_F (\mu' + l', \nu', \alpha', \beta', j)$ with $l, l' \in \mathbf{N}$ such that $\nu - \mu - l = \nu' - \mu' - l'$;
- (3) $(\lambda, \nu, \alpha, \beta, i) = (\lambda', \nu', \alpha', \beta', j)$, $\mu < \mu'$

This definition is independent of the choice of l, l' in view of the condition (A-1).

Lemma 2.3. (1) $<_H$ is a well-order.

- (2) If $\nu - \mu - \lambda = \nu' - \mu' - \lambda'$, then $(\lambda, \mu, \nu, \alpha, \beta, i) <_H (\lambda', \mu', \nu', \alpha', \beta', j)$ if and only if $(\mu, \nu, \alpha, \beta, i) <_F (\mu', \nu', \alpha', \beta', j)$.

Definition 2.4. An element P of $(A_{n+1}[x_0])^r$

of the form

$$P = \sum_{i=1}^r \sum_{\lambda, \mu, \nu, \alpha, \beta} a_{\lambda\mu\nu\alpha\beta i} x_0^\lambda t^\mu x^\alpha \partial_t^\nu \partial^\beta e_i$$

is said to be *F-homogeneous* of order m if $a_{\lambda\mu\nu\alpha\beta i} = 0$ whenever $\nu - \mu - \lambda \neq m$.

Definition 2.5. For an element P of $(A_{n+1})^r$ of the form (2.1), put $m := \min\{\nu - \mu \mid a_{\mu\nu\alpha\beta i} \neq 0 \text{ for some } \alpha, \beta \in \mathbf{N}^n \text{ and } i \in \{1, \dots, r\}\}$. Then the *F-homogenization* $P^h \in (A_{n+1}[x_0])^r$ of P is defined by

$$P^h := \sum_{i=1}^r \sum_{\mu, \nu, \alpha, \beta} a_{\mu\nu\alpha\beta i} x_0^{\nu-\mu-m} t^\mu x^\alpha \partial_t^\nu \partial^\beta e_i.$$

P^h is F -homogeneous of order m .

Proposition 2.6. Let N be a left $A_{n+1}[x_0]$ -submodule of $(A_{n+1}[x_0])^r$ generated by F -homogeneous operators. Then there exists an H -Gröbner basis (i.e. a Gröbner basis with respect to $<_H$) consisting of F -homogeneous operators. Moreover, such an H -Gröbner basis can be computed by the Buchberger algorithm.

Theorem 2.7. Let N be a left A_{n+1} -submodule of $(A_{n+1})^r$ generated by $P_1, \dots, P_d \in (A_{n+1})^r$. Let us denote by N^h the left $A_{n+1}[x_0]$ -submodule of $(A_{n+1}[x_0])^r$ generated by $(P_1)^h, \dots, (P_d)^h$. Let $\mathbf{G} = \{Q_1(x_0), \dots, Q_k(x_0)\}$ be an H -Gröbner basis of N^h consisting of F -homogeneous operators. Then $\mathbf{G}(1) := \{Q_1(1), \dots, Q_k(1)\}$ is an FW-Gröbner basis of N .

Let us denote by $\mathcal{D}_{K \times X} \big|_X$ the sheaf theoretic restriction of $\mathcal{D}_{K \times X}$ to $X = X \times \{0\}$. Then for a germ Q of $(\mathcal{D}_{K \times X} \big|_X)^r$ at $p \in X$, there exist $P \in$

$(A_{n+1})^r$ and $a(t, x) \in K[t, x]$ with $a(0, p) \neq 0$ so that $Q = a(t, x)^{-1}P$. For each integer m , we put

$$F_m((\mathcal{D}_{K \times X}|_X)^r)_p := \{a^{-1}P \mid P \in F_m((A_{n+1})^r), a = a(t, x) \in K[t, x], a(0, p) \neq 0\}.$$

For a germ Q of $(\mathcal{D}_{K \times X}|_X)^r$ at p , its F-order $\text{ord}_F(Q)$ is defined as the minimum $m \in \mathbf{Z}$ so that $P \in F_m((\mathcal{D}_{K \times X}|_X)^r)$. Put $m := \text{ord}_F(Q)$ and let $a(t, x) \in K[t, x]$ and $P \in (A_{n+1})^r$ be as above. Suppose that P is written in the form (2.1). Then the formal symbol $\bar{\sigma}(Q)$ of Q is defined by

$$\bar{\sigma}(Q) = \bar{\sigma}_m(Q) := a(0, x)^{-1} \sum_{i=1}^r \sum_{\nu-\mu=m} a_{\mu\nu\alpha\beta i} t^\mu x^\alpha \partial_i^\nu \partial^\beta e_i.$$

Definition 2.8. Let P be a nonzero element of $(A_{n+1})^r$ (resp. $(\mathcal{D}_{K \times X}|_X)^r$) of F-order m . Then we define $\phi(P)(s) \in (A_n[s])^r$ (resp. $(\mathcal{D}_X[s])^r$), by

$$\begin{aligned} \bar{\sigma}_0(t^m P) &= \phi(P)(t\partial_i) \text{ if } m \geq 0, \\ \bar{\sigma}_0(\partial_i^{-m} P) &= \phi(P)(t\partial_i) \text{ if } m < 0. \end{aligned}$$

Theorem 2.9. We use the same notation as in Theorem 2.7. Let \mathcal{N} be the left $\mathcal{D}_{K \times X}|_X$ -submodule of $(\mathcal{D}_{K \times X}|_X)^r$ generated by N . Let $\phi(\mathcal{N})$ be the left $\mathcal{D}_X[s]$ -submodule of $(\mathcal{D}_X[s])^r$ generated by the set $\{\phi(P)(s) \mid P \in \mathcal{N}, \text{ord}_F(P) = 0\}$. Then $\phi(\mathcal{N})$ is generated by $\phi(Q_1(1)), \dots, \phi(Q_k(1))$.

3. b -function of a D -module. Let \mathcal{M} be a left coherent $\mathcal{D}_{K \times X}|_X$ -module on X . We assume that a left A_{n+1} -submodule N of $(A_{n+1})^r$ is given explicitly so that $\mathcal{M} = \mathcal{D}_{K \times X}|_X \otimes_{A_{n+1}} M$ holds with $M := (A_{n+1})^r / N$. Set $\mathcal{N} := \mathcal{D}_{K \times X}|_X \otimes_{A_{n+1}} N \subset (\mathcal{D}_{K \times X}|_X)^r$. For each integer m , put

$$\begin{aligned} F_m(\mathcal{N}) &:= \mathcal{N} \cap F_m((\mathcal{D}_{K \times X}|_X)^r), \\ F_m(\mathcal{M}) &:= F_m((\mathcal{D}_{K \times X}|_X)^r) / F_m(\mathcal{N}). \end{aligned}$$

Then $\{F_m(\mathcal{M})\}_{m \in \mathbf{Z}}$ is a filtration of \mathcal{M} satisfying

$$F_k(\mathcal{D}_{K \times X}|_X)F_m(\mathcal{M}) = F_{k+m}(\mathcal{M})$$

for any $k, m \in \mathbf{Z}$. The b -function $b(s, p) \in K[s]$ of \mathcal{M} (with respect to the filtration $\{F_m(\mathcal{M})\}$) at $p \in X$ is the monic polynomial $b(s, p)$ of s of the least degree, if any, that satisfies

$$(3.1) \quad b(t\partial_i, p)(F_0(\mathcal{M})/F_{-1}(\mathcal{M}))_p = 0.$$

If such $b(s, p)$ exists, \mathcal{M} is called specializable along X at p . It is known that if \mathcal{M} is holonomic, then \mathcal{M} is specializable at any $p \in X$ ([13], [14]).

Let \mathbf{G} be an FW-Gröbner basis of N , which can be computed by the homogenization and the Buchberger algorithm with a set of generators as input (Theorem 2.7). Put $\phi(\mathbf{G}) := \{\phi(P) \mid P \in$

$\mathbf{G}\}$ and let $\psi(N)$ be the left $A_n[s]$ -submodule of $(A_n[s])^r$ generated by $\phi(\mathbf{G})$. Let $<_D$ be a total order on $L_0 \times \{1, \dots, r\}$ with $L_0 := \mathbf{N}^{1+2n}$ which satisfies (A-1) with L replaced by L_0 and (A-4) $(\alpha, i) >_D (0, i)$ for any $\alpha \in L_0 \setminus \{0\}$ and $i \in \{1, \dots, r\}$; (A-5) $|\beta| < |\beta'|$ implies $(\mu, \alpha, \beta, i) <_D (\mu', \alpha', \beta', j)$ for any $\mu, \mu' \in \mathbf{N}$, $\alpha, \alpha', \beta, \beta' \in \mathbf{N}^n$, $i, j \in \{1, \dots, r\}$.

Theorem 3.1. Under the above assumptions, let \mathbf{G}_1 be a Gröbner basis of $\psi(N)$ with respect to $<_D$ and put $\mathbf{G}_0 := \mathbf{G}_1 \cap K[s, x]^r$. Let \mathcal{T} be the $\mathcal{O}_X[s]$ -submodule of $(\mathcal{O}_X[s])^r$ generated by \mathbf{G}_0 . Then \mathcal{M} is specializable at p if and only if $\mathcal{T}_p \cap K[s]e_i \neq \{0\}$. If \mathcal{M} is specializable, then its b -function $b(s, p)$ is the monic polynomial of s of the least degree that satisfies $b(s, p)e_i \in \mathcal{T}_p \cap K[s]^r$ for any $i = 1, \dots, r$.

Since we have a set of generators of \mathcal{T} , it is easy to compute $\mathcal{T} \cap K[s]^r$. This can be done, e.g., by primary decomposition of the $K[s, x]$ -submodule of $K[s, x]^r$ which is generated by \mathbf{G}_0 (cf.[8]). Thus we obtain an algorithm of determining if \mathcal{M} is specializable at each point of X and of computing the b -function if that is the case.

4. Induced system. We retain the notation of the preceding section. The induced system of \mathcal{M} to X is the complex

$$\mathcal{M}_X^* : 0 \rightarrow \mathcal{M} \xrightarrow{t} \mathcal{M} \rightarrow 0$$

of left \mathcal{D}_X -modules, where the homomorphism t denotes the one defined by $t(u) = tu$ for $u \in \mathcal{M}$. Let us write $\mathcal{M}_X := \mathcal{M}/t\mathcal{M}$. For each integer m , we put

$$\text{gr}_m^F(\mathcal{M}) := F_m(\mathcal{M})/F_{m-1}(\mathcal{M}).$$

Lemma 4.1. Assume that $b(s) \in K[s]$ satisfies $b(t\partial_i)\text{gr}_0^F(\mathcal{M}) = 0$. Then the homomorphism $t : \text{gr}_{k+1}^F(\mathcal{M}) \rightarrow \text{gr}_k^F(\mathcal{M})$ is bijective if $b(k) \neq 0$.

Proposition 4.2. Assume that $b(s) \in K[s]$ satisfies $b(t\partial_i)\text{gr}_0^F(\mathcal{M}) = 0$. Put

$$\begin{aligned} k_1 &:= \max\{k \in \mathbf{Z} \mid b(k) = 0\}, \\ k_0 &:= \min\{k \in \mathbf{Z} \mid b(k) = 0\}. \end{aligned}$$

Then \mathcal{M}_X^* is quasi-isomorphic to the complex

$$0 \rightarrow F_{k_1+1}(\mathcal{M})/F_{k_0}(\mathcal{M}) \xrightarrow{t} F_{k_1}(\mathcal{M})/F_{k_0-1}(\mathcal{M}) \rightarrow 0$$

of left \mathcal{D}_X -modules. In particular, $t : \mathcal{M} \rightarrow \mathcal{M}$ is bijective if $b(k) \neq 0$ for any $k \in \mathbf{Z}$.

Proposition 4.3. Assume that there exists $b(s) \in K[s]$ and $m \in \mathbf{N}$ so that

$$b(t\partial_i)\partial_i^m \text{gr}_0^F(\mathcal{M}) = 0.$$

Assume, moreover, $b(k) \neq 0$ for any $k \in \mathbf{Z}$. Then the homomorphism $t : \mathcal{M} \rightarrow \mathcal{M}$ is injective.

Let P be an element of $F_m((A_{n+1})^r)$. Then we can write P in the form

$$P = \sum_{i=1}^r \sum_{k=0}^m P_{ik}(t\partial_t, x, \partial)\partial_t^k e_i + R$$

uniquely with $P_{ik} \in A_n[t\partial_t]$ and $R \in F_{-1}((A_{n+1})^r)$. Then we put

$$\rho(P, k_0) := \sum_{i=1}^r \sum_{k=k_0}^m P_{ik}(0, x, \partial)\partial_t^k e_i$$

for each integer k_0 with $0 \leq k_0 \leq m$.

Theorem 4.4. *Assume that $b(s) \in K[s]$ satisfies $b(t\partial_t)\text{gr}_0^F(\mathcal{M}) = 0$. Put*

$$k_1 := \max\{k \in \mathbb{Z} \mid b(k) = 0\},$$

$$k_0 := \max\{0, \min\{k \in \mathbb{Z} \mid b(k) = 0\}\}.$$

(We have $k_1 = m - 1$ and $k_0 = 0$ under the assumption of Proposition 4.3.) Let \mathbf{G} be an FW-Gröbner basis of N . Then we have an isomorphism

$$M_X \simeq \left(\bigoplus_{i=1}^r \bigoplus_{k=k_0}^{k_1} \mathcal{D}_X \partial_t^k e_i\right) / N_X$$

of left \mathcal{D}_X -modules, where N_X is the left \mathcal{D}_X -module generated by a finite set

$$\{\rho(\partial_t^j P, k_0) \mid P \in \mathbf{G}, j \in N, k_0 \leq j + \text{ord}_F(P) \leq k_1\}.$$

Our next aim is to give an algorithm for computing the structure of the kernel $\mathcal{H}^{-1}(M_X)$ of $t: \mathcal{M} \rightarrow \mathcal{M}$ as a left \mathcal{D}_X -module. For two integers $k_0 \leq k_1$, put

$$\tilde{\mathcal{D}}^{(k_0, k_1)} := \bigoplus_{i=1}^r \bigoplus_{k=k_0}^{k_1} \mathcal{D}_X[t\partial_t] S_k e_i,$$

where we put $S_k := \partial_t^k$ if $k \geq 0$, and $S_k := t^{-k}$ if $k < 0$. Let P be a section of $(\mathcal{D}_{K \times X}|_X)^r$ of F-order m . Then we can write P uniquely in the form

$$P = \sum_{i=1}^r \sum_{k=-\infty}^m P_{ik}(t\partial_t, x, \partial) S_k e_i$$

with $P_{ik} \in (\mathcal{D}_X[t\partial_t])^r$. Then we define

$$\tau(P, k_0) := \sum_{i=1}^r \sum_{k=k_0}^m P_{ik}(t\partial_t, x, \partial) S_k e_i.$$

Proposition 4.5. *Let \mathbf{G} be an FW-Gröbner basis of N . Then, for any integers $k_0 \leq k_1$, we have an isomorphism*

$$F_{k_1}(\mathcal{M})/F_{k_0-1}(\mathcal{M}) \simeq \tilde{\mathcal{D}}^{(k_0, k_1)} / \mathcal{G}^{(k_0, k_1)}$$

of left $\mathcal{D}_X[t\partial_t]$ -modules, where $\mathcal{G}^{(k_0, k_1)}$ is a left $\mathcal{D}_X[t\partial_t]$ -module generated by a finite set

$$\{\tau(S_j P, k_0) \mid P \in \mathbf{G}, j \in \mathbb{Z}, k_0 \leq j + \text{ord}_F(P) \leq k_1\}.$$

Let $\varphi: \tilde{\mathcal{D}}^{(k_0+1, k_1+1)} \rightarrow \tilde{\mathcal{D}}^{(k_0, k_1)}$ be a left \mathcal{D}_X -module homomorphism defined by

$$\varphi\left(\sum_{i=1}^r \sum_{k=k_0}^{k_1} P_{i, k+1}(t\partial_t, x, \partial) S_{k+1} e_i\right)$$

$$= \sum_{i=1}^r \sum_{k=k_0}^{k_1} P_{i, k+1}(t\partial_t - 1, x, \partial) T_k e_i$$

with

$$T_k := \begin{cases} S_k & (k \leq -1) \\ t\partial_t S_k & (k \geq 0). \end{cases}$$

Theorem 4.6. *Under the same assumptions as in Proposition 4.2, we have an isomorphism $\mathcal{H}^{-1}(M_X) \simeq \varphi^{-1}(\mathcal{G}^{(k_0, k_1)}) / \mathcal{G}^{(k_0+1, k_1+1)}$ as left $\mathcal{D}_X[t\partial_t]$ -modules. Moreover, $\varphi^{-1}(\mathcal{G}^{(k_0, k_1)}) / \mathcal{G}^{(k_0+1, k_1+1)}$ is finitely generated as left \mathcal{D}_X -module.*

The left $\mathcal{D}_X[t\partial_t]$ -module $\varphi^{-1}(\mathcal{G}^{(k_0, k_1)})$ can be easily computed by the same method as for computing ideal intersection and quotient in the polynomial ring by means of Gröbner basis (cf. [7]). Then by eliminating $t\partial_t$, also by means of a Gröbner basis, we get an algorithm of computing a presentation of $\mathcal{H}^{-1}(M_X)$ as a left \mathcal{D}_X -module.

5. Algebraic local cohomology. Let N be a left A_n -submodule of $(A_n)^r$ and put $M := (A_n)^r / N$ and $\mathcal{M} := \mathcal{D}_X \otimes_{A_n} M$. Let $f = f(x) \in K[x]$ be a polynomial and put $Y := \{x \in X \mid f(x) = 0\}$. Then the algebraic local cohomology group $\mathcal{H}_{[Y]}^j(\mathcal{M})$ has a structure of left \mathcal{D}_X -module and vanishes for $j \neq 0, 1$ ([11]). Our purpose is to give an algorithm of computing $\mathcal{H}_{[Y]}^j(\mathcal{M})$ as a left \mathcal{D}_X -module.

Let \mathcal{I} be a left ideal of $\mathcal{D}_{K \times X}$ generated by operators $t - f(x), \partial_1 + (\partial f / \partial x_1)\partial_t, \dots, \partial_n + (\partial f / \partial x_n)\partial_t$, and put $\mathcal{L} := \mathcal{D}_{K \times X} / \mathcal{I}$. Then by a method similar to that used by Malgrange [17], we get the following.

Theorem 5.1. *We have isomorphisms*

$$\mathcal{H}^j((\mathcal{M} \otimes_{\mathcal{O}_X} \mathcal{L})_X) \simeq \mathcal{H}_{[Y]}^{j+1}(\mathcal{M})$$

of left \mathcal{D}_X -modules for $j = -1, 0$.

Let p_1 and p_2 be the projections of $X \times K \times X$ to X and to $K \times X$ respectively and put

$$\Delta := \{(x, t, y) \in X \times K \times X \mid x = y\}.$$

Then we have by [11]

$$\mathcal{M} \overset{\mathbf{L}}{\otimes}_{\mathcal{O}_X} \mathcal{L} \simeq \mathcal{O}_\Delta \overset{\mathbf{L}}{\otimes}_{\mathcal{O}_{X \times K \times X}} (\mathcal{M} \overset{\mathbf{L}}{\otimes} \mathcal{L})$$

with

$$\mathcal{M} \overset{\mathbf{L}}{\otimes} \mathcal{L} := \mathcal{D}_{X \times K \times X} \otimes_{p_1^{-1}\mathcal{D}_X \otimes p_2^{-1}\mathcal{D}_{K \times X}} (p_1^{-1}\mathcal{M} \otimes_{\mathbb{K}} p_2^{-1}\mathcal{L}),$$

where $\overset{\mathbf{L}}{\otimes}$ denotes the left derived functor of \otimes in the derived category. In other words, $\mathcal{M} \otimes_{\mathcal{O}_X} \mathcal{L}$ coincides with the induced system of $\mathcal{M} \overset{\mathbf{L}}{\otimes} \mathcal{L}$ along Δ . It is easy to see that $\mathcal{M} \overset{\mathbf{L}}{\otimes} \mathcal{L}$ is specializable along Δ (in fact, Δ is non-characteristic for this module). Hence we can compute $\mathcal{M} \otimes_{\mathcal{O}_X} \mathcal{L}$ by applying Theorem 4.4 repeatedly with $k_0 = k_1 =$

0. Combining this fact with Theorems 4.4, 4.6, 5.1, we obtain an algorithm of computing $\mathcal{H}_{[Y]}^j(\mathcal{M})$ for $j = 0, 1$.

Theorem 5.2. *Assume $r = 1$ and let $u \in \mathcal{M}$ be the residue class of $1 \in \mathcal{D}_X$. Let $\tilde{b}(s)$ be the b -function of $\mathcal{M} \otimes_{\mathcal{O}_X} \mathcal{L}$ along X in the sense of Section 3 and let $b(s)$ be the b -function for f and u defined by (1.1) both at a point p of Y . Then we have the following:*

- (1) $b(s)$ divides $\tilde{b}(-s - 1)$;
- (2) if the homomorphism $f : \mathcal{M} \rightarrow \mathcal{M}$ defined by $f(v) = fv$ for $v \in \mathcal{M}$ is injective at p , then we have $b(s) = \pm \tilde{b}(-s - 1)$;
- (3) the homomorphism $f : \mathcal{M} \rightarrow \mathcal{M}$ is injective if and only if

$$\mathcal{H}^{-1}((\mathcal{M} \otimes_{\mathcal{O}_X} \mathcal{L})_X) = 0.$$

Thus the algorithm for $\tilde{b}(s)$ provides an algorithm to compute the b -function for f and u in generic cases. Since $f : \mathcal{O}_X \rightarrow \mathcal{O}_X$ is injective, we have an algorithm to compute the Bernstein-Sato polynomial of an arbitrary f .

It is also possible (in generic cases) to compute $\mathcal{H}_{[Y]}^j(\mathcal{M})$ for algebraic set Y of codimension greater than one. For example, let $f_1(x), f_2(x)$ be two polynomials and put

$$Y_i := \{x \in X \mid f_i(x) = 0\} \quad (i = 1, 2),$$

$$Y := Y_1 \cap Y_2.$$

Assume that $\mathcal{H}_{[Y_j]}^j(\mathcal{M}) = 0$ for $j \neq j_0$. Then we can compute

$$\mathcal{H}_{[Y]}^{j_0}(\mathcal{M}) = \mathcal{H}_{[Y_2]}^{j_0}(\mathcal{H}_{[Y_1]}^{j_0}(\mathcal{M}))$$

explicitly. The following computation was carried out by using Kan ([24]).

Example 5.3. Put $X = K^3$, $f_1 := x^2 - y^3$, $f_2 := y^2 - z^3$, and $Y := \{(x, y, z) \in X \mid f_1(x, y, z) = f_2(x, y, z) = 0\}$. Then we have $\mathcal{H}_{[Y]}^j(\mathcal{O}_X) = 0$ for $j \neq 2$ and

$$\mathcal{H}_{[Y]}^2(\mathcal{O}_X) \simeq \mathcal{D}_X / \mathcal{I},$$

where \mathcal{I} is the left ideal of \mathcal{D}_X generated by f_1, f_2 and

$$9x\partial_x + 6y\partial_y + 4z\partial_z + 30,$$

$$9z^2y^2\partial_x + 6z^2x\partial_y + 4yx\partial_z.$$

Let u_j be the residue class of f_j^{-1} in $\mathcal{H}_{[Y_j]}^1(\mathcal{O}_X) = \mathcal{O}_X[f_j^{-1}] / \mathcal{O}_X$ with $Y_j := \{(x, y, z) \mid f_j(x, y, z) = 0\}$. Then the b -function for f_2 and u_1 at $0 = (0, 0, 0)$ is

$$(s + 1) \left(s + \frac{1}{12}\right) \left(s + \frac{5}{12}\right) \left(s + \frac{7}{12}\right)$$

$$\left(s + \frac{5}{6}\right) \left(s + \frac{11}{12}\right) \left(s + \frac{7}{6}\right),$$

while the b -function for f_1 and u_2 at 0 is

$$(s + 1) \left(s + \frac{7}{18}\right) \left(s + \frac{11}{18}\right) \left(s + \frac{13}{18}\right)$$

$$\left(s + \frac{5}{6}\right) \left(s + \frac{17}{18}\right) \left(s + \frac{19}{18}\right) \left(s + \frac{7}{6}\right) \left(s + \frac{23}{18}\right).$$

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