

## 6. Convergence Concepts in Semi-ordered Linear Spaces. I

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Concerning semi-ordered linear spaces, L. Kantorovitch [1] gave originally two different concepts of convergence, that is, order convergence and star convergence. One of the authors introduced two other concepts, that is, dilatator convergence in [2] and individual convergence in [3], which are essentially equivalent to each other. Combining these concepts, we also obtain star-individual convergence in [4]. In this paper we want to discuss these concepts of convergence and their combinations more systematically. In the sequel we will use the terminologies and notations in the book [4].

Let  $R$  be a continuous semi-ordered linear space. We consider the order convergence basic, that is, for a sequence  $a_\nu \in R$  ( $\nu=0, 1, 2, \dots$ ),  $a_0 = \lim_{\nu \rightarrow \infty} a_\nu$  means

$$a_0 = \bigcap_{\nu=1}^{\infty} \bigcup_{\mu \geq \nu} a_\mu = \bigcup_{\nu=1}^{\infty} \bigcap_{\mu \geq \nu} a_\mu.$$

In the sequel we denote by  $\{a_\nu\}_\nu$  an arbitrary sequence  $a_\nu \in R$  ( $\nu=0, 1, 2, \dots$ ) and  $\{a_\nu\}_{\nu \geq 1}$  means  $a_\nu$  ( $\nu=1, 2, \dots$ ). A mapping  $\alpha$  of all sequences  $\{a_\nu\}_\nu$  to sequences  $\{a_\nu^\alpha\}_\nu$  is called an *operator*, if

$$1) \quad a_0 = \lim_{\nu \rightarrow \infty} a_\nu \quad \text{implies} \quad a_0^\alpha = \lim_{\nu \rightarrow \infty} a_\nu^\alpha,$$

$$2) \quad \{a_\nu^\alpha\}_{\nu \geq 1} \quad \text{depends only upon} \quad \{a_\nu\}_{\nu \geq 1}$$

that is,  $a_\nu = b_\nu$  ( $\nu=1, 2, \dots$ ) implies  $a_\nu^\alpha = b_\nu^\alpha$  ( $\nu=1, 2, \dots$ ). An operator  $\alpha$  is said to be *linear* if

$$(\alpha a_\nu + \beta b_\nu)^\alpha = \alpha a_\nu^\alpha + \beta b_\nu^\alpha \quad (\nu=0, 1, 2, \dots).$$

For two operators  $\alpha, \beta$ , putting

$$\alpha^\alpha \beta = (\alpha^\alpha)^\beta \quad (\nu=0, 1, 2, \dots),$$

we also obtain an operator  $\alpha\beta$ , which will be called the *product* of  $\alpha$  and  $\beta$ . With this definition, we have obviously

$$(\alpha\beta)\gamma = \alpha(\beta\gamma).$$

$\alpha$  is said to *commute*  $\beta$ , if  $\alpha\beta = \beta\alpha$ .

A set  $\mathfrak{A}$  of operators is called a *process*, if for any two sequences  $\{a_\nu\}_\nu, \{b_\nu\}_\nu$  with  $a_0 \neq b_0$  we can find  $\alpha \in \mathfrak{A}$  for which  $a_0^\alpha \neq b_0^\alpha$ . A set  $A$  of processes is called a *modifier*, if for any  $\mathfrak{A}_1, \mathfrak{A}_2 \in A$  we can find  $\mathfrak{A} \in A$  for which  $\mathfrak{A} \subset \mathfrak{A}_1, \mathfrak{A}_2$ . For two modifiers  $A, B$  we write  $A \geq B$ , if for any  $\mathfrak{A} \in A$  we can find  $\mathfrak{B} \in B$  for which  $\mathfrak{A} \supset \mathfrak{B}$ . If  $A \geq B$  and  $B \geq A$  at the same time, we write  $A = B$ .

Let  $A$  and  $B$  be modifiers. For a process  $\mathfrak{A} \in A$  and a system of processes  $\mathfrak{B}_a \in B$  ( $a \in \mathfrak{A}$ ) we see easily that the set

$$\{\alpha\beta: a \in \mathfrak{A}, \beta \in \mathfrak{B}_a\}$$

also is a process, and furthermore that all such processes constitute a modifier, which will be called the *product* of  $A$  and  $B$ , and denoted by  $AB$ . We also see that the system

$$\{ab: a \in \mathfrak{A}, b \in \mathfrak{B}\} \quad (\mathfrak{A} \in A, \mathfrak{B} \in B)$$

is a modifier, which will be called the *direct product* of  $A$  and  $B$  and denoted by  $A \circ B$ .

For modifiers  $A, B, C$  we have obviously by definition

- (1)  $(AB)C = A(BC), (A \circ B) \circ C = A \circ (B \circ C),$
- (2)  $A \circ B \geq AB,$
- (3)  $A \geq B$  implies  $AC \geq BC, CA \geq CB, A \circ C \geq B \circ C, C \circ A \geq C \circ B,$
- (4)  $(AB) \circ C \geq A(B \circ C), A \circ (BC) \geq (A \circ B)C.$

For a modifier  $A$ , a sequence  $\{a_\nu\}_{\nu \geq 1}$  is said to be  $A$ -convergent, if we can find  $a_0 \in R$  and  $\mathfrak{A} \in A$  such that

$$a_0^{\mathfrak{A}} = \lim_{\nu \rightarrow \infty} a_\nu^{\mathfrak{A}} \quad \text{for all } a \in \mathfrak{A}.$$

In this case we see easily that such  $a_0$  is determined uniquely. Thus such  $a_0$  is called the  $A$ -limit of  $\{a_\nu\}_{\nu \geq 1}$  and we write

$$a_0 = A\text{-}\lim_{\nu \rightarrow \infty} a_\nu.$$

With this definition we have obviously

**Theorem 1.** *For two modifications  $A, B$  we have*

$$a_0 = AB\text{-}\lim_{\nu \rightarrow \infty} a_\nu$$

*if and only if we can find  $\mathfrak{A} \in A$  such that*

$$a_0^{\mathfrak{A}} = B\text{-}\lim_{\nu \rightarrow \infty} a_\nu^{\mathfrak{A}} \quad \text{for all } a \in \mathfrak{A}.$$

For two modifiers  $A, B$ , we write  $A \succ B$  if

$$a_0 = A\text{-}\lim_{\nu \rightarrow \infty} a_\nu \quad \text{implies} \quad a_0 = B\text{-}\lim_{\nu \rightarrow \infty} a_\nu;$$

and  $A$  is said to be *equivalent* to  $B$  and denoted by  $A \sim B$ , if  $A \succ B$  and  $B \succ A$  at the same time. With this definition we see easily

- (5)  $A \geq B$  implies  $A \succ B,$
- (6)  $A \succ B$  implies  $CA \succ CB, C \circ A \succ C \circ B,$
- (7)  $A \succ A \circ B \succ AB.$

A modifier  $A$  is said to *commute* an operator  $a$ , if

$$a_0 = A\text{-}\lim_{\nu \rightarrow \infty} a_\nu \quad \text{implies} \quad a_0^{\mathfrak{A}} = A\text{-}\lim_{\nu \rightarrow \infty} a_\nu^{\mathfrak{A}}.$$

With this definition we conclude immediately by Theorem 1

**Theorem 2.** *For two modifiers  $A, B$ , if every operator of  $A$  commutes an operator  $c$  and  $B$  commutes  $c$ , then  $AB$  commutes  $c$ .*

As the simplest operator we have the *identity* 1, that is,  $a_1^{\mathfrak{A}} = a_\nu$  ( $\nu = 0, 1, 2, \dots$ ). The modifier, which consists of only one process [1], is denoted by  $O$ .  $O$ -convergence coincides obviously with the order convergence, that is,  $a_0 = O\text{-}\lim_{\nu \rightarrow \infty} a_\nu$  if and only if  $a_0 = \lim_{\nu \rightarrow \infty} a_\nu$ . Further-

more we have for every modifier  $A$

$$O \succ A, \quad OA = AO = O \circ A = A \circ O = A.$$

For every subsequence  $\{\mu_\nu\}_{\nu \geq 1}$  of  $\{1, 2, \dots\}$ , putting

$$\alpha_0^{\mathfrak{s}} = \alpha_0, \quad \alpha_\nu^{\mathfrak{s}} = \alpha_{\mu_\nu} \quad (\nu = 1, 2, \dots),$$

we obtain an operator  $\mathfrak{s}$ , which will be called a *sub. operator* and denoted by  $\mathfrak{s}\{\mu_\nu\}$ , if we need to indicate  $\{\mu_\nu\}$ . For two sub. operators  $\mathfrak{s}_1, \mathfrak{s}_2$ , the product  $\mathfrak{s}_1\mathfrak{s}_2$  also is a sub. operator. We write  $\mathfrak{s}\{\mu_\nu\} \geq \mathfrak{s}\{\rho_\nu\}$  if  $\{\rho_\nu\}$  is a subsequence of  $\{\mu_\nu\}$ .

We denote by  $S$  the modifier, which consists of all such processes  $\mathfrak{S}$  of sub. operators that

- 1)  $\mathfrak{s} \leq \mathfrak{s}_0 \in \mathfrak{S}$  implies  $\mathfrak{s} \in \mathfrak{S}$ ,
- 2) for any sub. operator  $\mathfrak{s}$  we can find  $\mathfrak{s}_0 \in \mathfrak{S}$  for which  $\mathfrak{s} \geq \mathfrak{s}_0$ .

With this definition we have obviously

$$(8) \quad SS = S \circ S = S.$$

For every projector  $[p]$ , putting  $a_\nu^l = [p]a_\nu$  ( $\nu = 0, 1, 2, \dots$ ), we obtain an operator  $l$ , which will be called a *loc. operator* and denoted by  $l[p]$ , if we need to indicate  $[p]$ . We write  $l[p] \geq l[q]$ , if  $[p] \geq [q]$ . We have obviously  $l[p]l[q] = l[p][q]$  and  $l\mathfrak{s} = \mathfrak{s}l$  for every loc. operator  $l$  and sub. operator  $\mathfrak{s}$ .

We denote by  $L$  the modifier which consists of all such processes  $\mathfrak{L}$  of loc. operators that

- 1)  $l \leq l_0 \in \mathfrak{L}$  implies  $l \in \mathfrak{L}$ ,
- 2) for any loc. operator  $l$  we can find  $l_0 \in \mathfrak{L}$  for which  $l \geq l_0$ .

With this definition we have obviously

$$(9) \quad LL = L \circ L = L.$$

Since  $\mathfrak{s}l = l\mathfrak{s}$  for every loc. operator  $l$  and sub. operator  $\mathfrak{s}$ , we have

$$(10) \quad L \circ S = S \circ L.$$

**Lemma 1.** *Let  $A$  be a modifier, which commutes every loc. operator. In order that*

$$a_0 = LA\text{-lim}_{\nu \rightarrow \infty} a_\nu,$$

*it is necessary and sufficient that we can find a system of projectors  $[p_\lambda]$  ( $\lambda \in A$ ) such that*

$$\bigcup_{\lambda \in A} [p_\lambda] \bigcup_{\nu=1}^{\infty} [a_\nu] = \bigcup_{\nu=1}^{\infty} [a_\nu]$$

$$[p_\lambda]a_0 = A\text{-lim}_{\nu \rightarrow \infty} [p_\lambda]a_\nu \quad \text{for all } \lambda \in A.$$

**Proof.** We need only to prove the sufficiency. For such a system of projectors  $[p_\lambda]$  ( $\lambda \in A$ ), denoting by  $\mathfrak{L}$  the set of all such  $l[p]$  that  $[p] \leq [p_\lambda]$  for some  $\lambda \in A$  or  $[p][p_\lambda] = 0$  for all  $\lambda \in A$ , we see easily that  $\mathfrak{L} \in L$ , and

$$a_0^l = A\text{-lim}_{\nu \rightarrow \infty} a_\nu^l \quad \text{for all } l \in \mathfrak{L},$$

because  $A$  commutes  $l$  by assumption.

For two elements  $p \geq 0 \geq q$  in  $R$ , putting

$$a_\nu^i = (a_\nu \wedge p) \vee q \quad (\nu = 0, 1, 2, \dots),$$

we obtain an operator  $i$ , which will be called an *ind. operator* and

denoted by  $i(p, q)$  if we need to indicate  $p, q$ . We write  $i(p, q) \geq i(r, s)$  if  $p \geq r \geq s \geq q$ . We have obviously

$$i(p, q)i(r, s) = i(p \wedge r, q \vee s)$$

and  $iI = Ii$ ,  $i\mathfrak{s} = \mathfrak{s}i$  for every loc. operator  $I$  and sub. operator  $\mathfrak{s}$ .

We denote by  $I$  the modifier which consists of only one process of all ind. operators. With this definition we have obviously

$$(11) \quad II = I \circ I = I.$$

From the proof of Theorem 1.1 in [3], we conclude easily

$$(12) \quad I \sim L.$$

**Lemma 2.** *In order that  $a_0 = I\text{-lim } a_\nu$ , it is necessary and sufficient that we can find a sequence  $0 \leq p_1 \leq p_2 \leq \dots$  such that*

$$(a_0 \wedge p_\mu) \vee (-p_\mu) = \lim_{\nu \rightarrow \infty} (a_\nu \wedge p_\mu) \vee (-p_\mu) \quad \text{for all } \mu = 1, 2, \dots,$$

$$\lim_{\mu \rightarrow \infty} (x \wedge p_\mu) \vee (-p_\mu) = x \quad \text{for all } x \in [a_1, a_2, \dots]R.$$

**Proof.** We need only to prove the sufficiency. Putting  $i_\mu = i(p_\mu, -p_\mu)$  ( $\mu = 1, 2, \dots$ ), we obtain by assumption for any ind. operator  $i$

$$(\overline{\lim}_{\nu \rightarrow \infty} a_\nu^{i_\mu})^{i_\mu} = \overline{\lim}_{\nu \rightarrow \infty} a_\nu^{i_\mu i_\mu} = (\overline{\lim}_{\nu \rightarrow \infty} a_\nu^{i_\mu})^{i_\mu} = a_0^{i_\mu i_\mu} = (a_0^{i_\mu})^{i_\mu}$$

Thus, making  $\mu \rightarrow \infty$ , we obtain  $\overline{\lim}_{\nu \rightarrow \infty} a_\nu^i = a_0^i$ . We conclude similarly also that  $\underline{\lim}_{\nu \rightarrow \infty} a_\nu^i = a_0^i$ . Therefore  $a_0 = I\text{-lim } a_\nu$  by definition.

As  $iI = Ii$  and  $I$  consists of only one process, we have by definition

$$(13) \quad I \circ L = L \circ I = LI.$$

Recalling (12), we obtain by (9), (11)

$$(14) \quad LI \sim IL \sim I.$$

As  $i\mathfrak{s} = \mathfrak{s}i$ , we have

$$(15) \quad I \circ S = S \circ I = SI.$$

As  $I \circ S \geq IS$  by (2), we have hence  $SI \succ IS$  by (5). Now we shall prove

$$(16) \quad SI \sim IS.$$

We suppose  $a_0 = IS\text{-lim } a_\nu$ . Putting  $p_\mu = \mu \sum_{\nu=1}^{\mu} |a_\nu|$  ( $\mu = 1, 2, \dots$ ), we see easily that the sequence  $0 \leq p_1 \leq p_2 \leq \dots$  satisfies the condition of Lemma 2. For any sub. operator  $\mathfrak{s}$ , we can find by assumption a sequence of operators  $\mathfrak{s}_1 \geq \mathfrak{s}_2 \geq \dots$  such that

$$(a_0 \wedge p_\mu) \vee (-p_\mu) = \lim_{\nu \rightarrow \infty} (a_\nu^{\mathfrak{s}_\mu} \wedge p_\mu) \vee (-p_\mu) \quad (\mu = 1, 2, \dots).$$

Then we can find by the diagonal method a sub. operator  $\mathfrak{s}_0 \leq \mathfrak{s}$  such that

$$(a_0 \wedge p_\mu) \vee (-p_\mu) = \lim_{\nu \rightarrow \infty} (a_\nu^{\mathfrak{s}_0} \wedge p_\mu) \vee (-p_\mu) \quad (\mu = 1, 2, \dots).$$

Thus we have  $a_0 = SI\text{-lim } a_\nu$ , and therefore  $IS \succ SI$  by definition.

A modifier is said to be *regular*, if it commutes every sub., loc. and ind. operators. The modifier  $O$  is obviously regular.

**Lemma 3.** *If a modifier  $A$  is regular, then all  $SA$ ,  $LA$  and  $IA$  are regular, and  $S \circ A \prec A$ ,  $L \circ A \prec A$ ,  $I \circ A \prec A$ .*

**Proof.** By virtue of Theorem 2, both  $LA$  and  $IA$  are regular. To prove that  $SA$  is regular, we need only to show that  $SA$  commutes every sub. operator. We suppose that  $a_0 = SA\text{-}\lim_{\nu \rightarrow \infty} a_\nu$ . Then we can find by Theorem 1 a process  $\mathfrak{S} \in S$  such that

$$a_0^{\mathfrak{S}} = A\text{-}\lim_{\nu \rightarrow \infty} a_\nu^{\mathfrak{S}} \quad \text{for all } \mathfrak{S} \in \mathfrak{S}.$$

For any sub. operator  $\mathfrak{S}_0$ , we obtain hence

$$a_0^{\mathfrak{S}_0 \mathfrak{S}} = A\text{-}\lim_{\nu \rightarrow \infty} a_\nu^{\mathfrak{S}_0 \mathfrak{S}} \quad \text{for } \mathfrak{S}_0 \mathfrak{S} \in \mathfrak{S}.$$

Putting  $\mathfrak{S}_0 = \{\mathfrak{S}: \mathfrak{S}_0 \mathfrak{S} \in \mathfrak{S}\}$ , we see easily that  $\mathfrak{S}_0 \in S$ . Thus we have  $a_0^{\mathfrak{S}_0} = SA\text{-}\lim_{\nu \rightarrow \infty} a_\nu^{\mathfrak{S}_0}$ . Therefore  $SA$  commutes every sub. operator. If  $A$  is regular, then we have obviously  $S \circ A \prec A$ ,  $L \circ A \prec A$ ,  $I \circ A \prec A$  by definition.

**Lemma 4.** *If  $R$  is super-universally continuous and a modifier  $A$  commutes every loc. operator, then we have*

$$(L \circ S)A \sim LSA \succ SLA.$$

**Proof.** We suppose that  $a_0 = LSA\text{-}\lim_{\nu \rightarrow \infty} a_\nu$ . As  $R$  is super-universally continuous by assumption, we can find  $[p_\mu]$  ( $\mu=1, 2, \dots$ ) such that

$$[p_\mu]a_0 = SA\text{-}\lim_{\nu \rightarrow \infty} [p_\mu]a_\nu \quad (\mu=1, 2, \dots), \quad \bigcup_{\mu=1}^{\infty} [p_\mu] \geq \bigcup_{\nu=1}^{\infty} [a_\nu].$$

Then we can find  $\mathfrak{S}_\mu \in S$  by definition such that

$$[p_\mu]a_0^{\mathfrak{S}_\mu} = A\text{-}\lim_{\nu \rightarrow \infty} [p_\mu]a_\nu^{\mathfrak{S}_\mu} \quad \text{for all } \mathfrak{S}_\mu \in \mathfrak{S}_\mu \quad (\mu=1, 2, \dots).$$

Denoting by  $\mathfrak{S}$  the intersection of all  $\mathfrak{S}_\mu$  ( $\mu=1, 2, \dots$ ), we see easily by the diagonal method that  $\mathfrak{S} \in S$ . Denoting by  $\mathfrak{V}$  the set of all  $I[p]$  such that  $[p] \leq [p_\mu]$  for some  $\mu=1, 2, \dots$  or  $[p][p_\mu]=0$  for all  $\mu=1, 2, \dots$ , we see easily that  $\mathfrak{V} \in L$ , because  $A$  commutes every loc. operator by assumption. Thus we have

$$a_0^{\mathfrak{V} \mathfrak{S}} = A\text{-}\lim_{\nu \rightarrow \infty} a_\nu^{\mathfrak{V} \mathfrak{S}} \quad \text{for all } I \in \mathfrak{V}, \mathfrak{S} \in \mathfrak{S},$$

and hence  $a_0 = (L \circ S)A\text{-}\lim_{\nu \rightarrow \infty} a_\nu$ . Therefore we have  $LSA \succ (L \circ S)A$ .

On the other hand we have  $(L \circ S)A \succ LSA$  by (2), (3). Consequently  $(L \circ S)A \sim LSA$ . As  $L \circ S = S \circ L \geq SL$ , we obtain hence  $LSA \succ SLA$ .

A modifier is said to be *standard*, if it is composed only of  $O$ ,  $S$ ,  $L$ ,  $I$  by the product and the direct product.

**Theorem 3.** *If  $R$  is super-universally continuous, then every standard modifier is equivalent to one of  $O$ ,  $S$ ,  $L$ ,  $LS$ ,  $SL$ .*

**Proof.** We need only to show  $SLS \sim LSL \sim ILS \sim ISL \sim SL$ . As  $LS \succ SL$  by Lemma 4, we obtain by (6), (8), (7):  $SLS \succ SSL = SL \succ SLS$ , and by (9), Lemma 3:  $LSL \succ SLL = SL \succ LSL$ . As  $L \sim I$  by (12), we have by (6), (16), (11):  $ISL \sim ISI \sim IIS = IS \sim SI \sim SL$ . As  $IL \leq LI$  by

(13), (2), we have by (3), (16), (12):  $ILS \leq LIS \sim LSI \sim LSL \sim SL$ . On the other hand we have  $ILS \succ ISL$  by Lemma 4 and (6), and  $ISL \sim SL$ , as proved just above.

**Theorem 4.** *If  $R$  is super-universally continuous and complete,\*<sup>o</sup> then every standard modifier is equivalent to one of  $O$  and  $S$ .*

**Proof.** If  $R$  is super-universally continuous and complete, then we see easily  $I \sim L \sim O$ . Thus we obtain by Lemmas 3 and 4

$$S \succ LS \succ SL \sim SO = S.$$

Therefore we conclude our assertion from Theorem 3.

### References

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\*<sup>o</sup>) A semi-ordered linear space is said to be *complete* if every orthogonal sequence of elements is bounded (cf. [5]).