

A NOTE ON Π_1^1 ORDINALS

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In reference [3] Tanaka proves (among other things) that the Π_1^1 ordinals are precisely the ordinals recursive in Kleene's set O .¹ The purpose of this note is to show how this result may be neatly obtained as a corollary of reference [2]. Here is some background on the matter.

An ordinal α is called *recursive* (Σ_1^1 , Π_1^1 , recursive in O , etc.) if there is a recursive (Σ_1^1 , Π_1^1 , recursive in O , etc.) well-ordering of natural numbers with order type α . O is the set of notations of Kleene's system S_3 , and ω_1 (ω_1^O) is the least ordinal that is not recursive (recursive in O). Some well-known facts of ordinal notation theory are the following, where each set is an initial segment of ordinals.

$$(1) \quad \{\alpha : \alpha \text{ is recursive}\} = \{\alpha : \alpha \text{ is } \Sigma_1^1\} \subset \{\alpha : \alpha \text{ is } \Pi_1^1\} \subseteq \{\alpha : \alpha \text{ is recursive in } O\}.$$

Tanaka's result concerns the final inclusion in (1):

PROPOSITION. $\{\alpha : \alpha \text{ is } \Pi_1^1\} = \{\alpha : \alpha \text{ is recursive in } O\}$

*Proof, derived from [2].*² We show that every ordinal less than ω_1^O is Π_1^1 . In the notation of [2], $W[A]$ is the set of all natural numbers e such that the partial recursive function $\{e\}$ is defined on $A \times A$, and $\{(x, y) : \{e\}(x, y) = 0\}$ well-orders A . If A is infinite, then the order types of such well-orderings comprise a segment of ordinals beginning with ω . The least upper bound of the segment is denoted by " $|W[A]|$ ". Remark 4.8 and theorem 7.3 of [2] imply that

$$(2) \quad |W[O]| = \omega_1^O,$$

which is exactly the needed fact:

1. This fact is also proved in §VI.1 of [1].

2. This proof, more direct than the one appearing in [1], was suggested to the author by Professor Richter.

If α is an infinite ordinal less than ω_1^0 , then, by (2), there is an $e \in W[\mathbf{O}]$ such that α is the order type of the Π_1^1 well-ordering $\{(x, y) : x \in \mathbf{O} \ \& \ y \in \mathbf{O} \ \& \ \{e\}(x, y) = 0\}$. Q.E.D.

To generalize the proposition, we set $\mathbf{O}(0) =_{\text{def}} \phi$ and $\mathbf{O}(n+1) =_{\text{def}} \mathbf{O}^{\mathbf{O}(n)}$. Then one may prove by induction, beginning with (2), that

$$\{\alpha : \alpha \text{ is } \Pi_1^1 \mathbf{O}^{\mathbf{O}(n)}\} = \{\alpha : \alpha \text{ is recursive in } \mathbf{O}(n+1)\}.$$

REFERENCES

- [1] Gass, F. S., *The Present State of Ordinal Notation Theory*, Ph.D. Thesis, Dartmouth College (1968).
- [2] Richter, W., "Extensions of the constructive ordinals," *The Journal of Symbolic Logic*, vol. 30 (1965), pp. 193-211.
- [3] Tanaka, H., "On analytic well-orderings," *The Journal of Symbolic Logic*, vol. 35 (1970), pp. 198-204.

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