



# A family of 2-arc transitive pentagraphs with unbounded valency

Andries E. Brouwer      Johannes Huizinga

## Abstract

We construct polygonal graphs on the points of a generalized polygon in general position with respect to a polarity.

**Keywords:** pentagraphs, generalized polygon, polarity

**MSC 2000:** 05E18, 20B25, 51E12

## 1 Polygonal graphs

Let  $(X, L, I)$  be a generalized  $n$ -gon with polarity  $\sigma$ . Let  $Z$  be the set of points in general position with respect to  $\sigma$ , i.e.,  $Z = \{x \in X \mid d(x, x^\sigma) \geq n - 1\}$ , with distances measured in the point-line incidence graph  $\Sigma$  of  $(X, L, I)$ . (Thus, if  $n$  is even then  $d(x, x^\sigma) = n - 1$  and if  $n$  is odd then  $d(x, x^\sigma) = n$  for  $x \in Z$ .) Define a graph  $\Gamma$  with vertex set  $Z$  by letting distinct vertices  $x, y \in Z$  be adjacent (notation  $x \sim y$ ) when  $x I y^\sigma$ .

**Theorem 1.1.** *If  $n$  is odd, then  $\Gamma$  has girth  $g \geq n$  and each edge is contained in a unique  $n$ -gon. If  $n$  is even, then  $\Gamma$  has girth  $g \geq n + 1$  and each 2-path is contained in a unique  $(n + 1)$ -gon.*

*Proof.* Let us first collect information about the vertex set  $Z$ .

**Step 1.** *If  $x_0 I x_1^\sigma I x_2 I \dots I x_{l-1} I x_0^\sigma I x_1 I \dots I x_{l-1}^\sigma I x_0$  is a self-polar  $2l$ -circuit in  $\Sigma$ , and  $l \leq n + 1$ , then  $x_i \in Z$  ( $0 \leq i \leq l - 1$ ).*

(Indeed, if  $d_\Sigma(x_i, x_i^\sigma) = m$ , then we find an  $(m + l)$ -circuit in  $\Sigma$ , so that  $m + l \geq 2n$ .)

**Step 2.** *If  $n$  is even, and  $x \in Z$ , and  $x I x_1^\sigma I \dots I x_{n-2} I x^\sigma$  is the unique path of length  $n - 1$  joining  $x$  to  $x^\sigma$  in  $\Sigma$ , then  $x_i \notin Z$  ( $1 \leq i \leq n - 2$ ).*

(Indeed, applying  $\sigma$  to this path, we find another path that must coincide with this path, so that  $x_i^\sigma = x_{n-1-i}$  ( $1 \leq i \leq n - 2$ ).

Now look at the graph  $\Gamma$ . Note that if  $x \sim y \sim z$  in  $\Gamma$ , then  $x I y^\sigma I z$  in  $\Sigma$ .

**Step 3.**  *$\Gamma$  does not have even circuits of length less than  $2n$  and no odd circuits of length less than  $n$ . In particular, if two vertices have distance less than  $n$  in  $\Gamma$ , then there is a unique shortest path in  $\Gamma$  joining them.*

(Indeed, if  $x_0 \sim x_1 \sim \dots \sim x_{l-1} \sim x_0$  is an  $l$ -circuit in  $\Gamma$ , and  $l$  is even, then  $x_0 I x_1^\sigma I x_2 I \dots I x_{l-1}^\sigma I x_0$  is an  $l$ -circuit in  $\Sigma$ , and it follows that  $l \geq 2n$ . If  $l$  is odd, then  $x_0 I x_1^\sigma I x_2 I \dots I x_{l-1} I x_0^\sigma$  is an  $l$ -path in  $\Sigma$ , and by Step 2 we have  $l \geq n$ .)

**Step 4.** *If  $n$  is odd, then each edge is contained in a unique  $n$ -gon.*

(Indeed, if  $n$  is odd, and  $xy$  is an edge in  $\Gamma$ , then  $d_\Sigma(x, y) = n - 1$  and in  $\Sigma$  there is a unique geodesic  $x = x_0 I x_1^\sigma I x_2 I \dots I x_{n-1} = y$  joining  $x$  and  $y$ . This geodesic is part of the self-polar  $2n$ -circuit

$$x_0 I x_1^\sigma I x_2 I \dots I x_{n-1} I x_0^\sigma I x_1 I x_2^\sigma I \dots I x_{n-1}^\sigma I x_0$$

in  $\Sigma$ . Thus, by Step 1,  $x_0 \sim x_1 \sim \dots \sim x_{n-1} \sim x_0$  is the unique  $n$ -gon on the edge  $xy$  in  $\Gamma$ .)

**Step 5.** *If  $n$  is even, then each 2-path is contained in a unique  $(n + 1)$ -gon.*

(Indeed, if  $x \sim y \sim z$  in  $\Gamma$ , then  $d_\Sigma(x, y) = d_\Sigma(y, z) = n$  (since by Step 2 the unique point on  $y^\sigma$  that has distance  $n - 2$  to  $y$  is not in  $Z$ ). Let  $x = x_0 I x_1^\sigma I x_2 I \dots I x_{n-1}^\sigma = z^\sigma$  be the unique path of length  $n - 1$  in  $\Sigma$  joining  $x$  to  $z^\sigma$ . Then

$$x_0 I x_1^\sigma I x_2 I \dots I x_{n-1}^\sigma I y I x_0^\sigma I x_1 I \dots I x_{n-1} I y^\sigma I x_0$$

is a self-polar  $(2n + 2)$ -circuit in  $\Sigma$ . Thus, by Step 1,  $x_0 \sim x_1 \sim \dots \sim x_{n-1} = z \sim y \sim x$  is the unique  $(n + 1)$ -gon on the path  $x \sim y \sim z$  in  $\Gamma$ .)

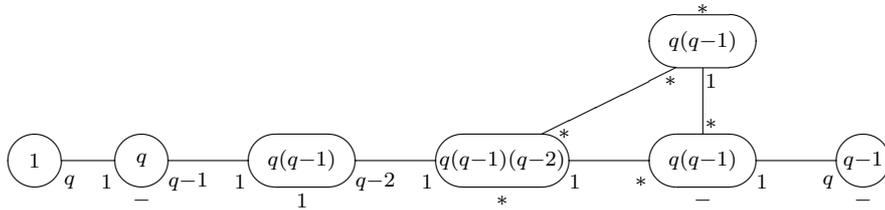
This completes the proof. □

If  $(X, L, I)$  is a  $2m$ -gon, then  $\Gamma$  is a single edge. If  $(X, L, I)$  is a  $(2m + 1)$ -gon, then there are two possible polarities  $\sigma$ ; for one choice of  $\sigma$  the graph  $\Gamma$  consists of a single vertex; for the other choice it is a  $(2m + 1)$ -gon itself.

## 2 Pentagraphs

Now let us specialize to the finite case  $n = 4$ , i.e., let  $(X, L, I)$  be a generalized quadrangle of order  $q$  with a polarity  $\sigma$ . Then  $2q$  is a square, cf. Payne [4]. Examples exist when  $q$  is an odd power of 2, cf. Tits [9]. We define the graph  $\Gamma$  as before. As we shall see,  $\Gamma$  is a *pentagraph*, that is, any 2-path in  $\Gamma$  is contained in a unique pentagon. (For this concept, and other examples, and some theory, see Perkel [5, 6, 7, 8] and Ivanov [3].)

**Theorem 2.1.**  $\Gamma$  is a pentagraph of valency  $q$  on  $q^3 + q$  vertices, and has distance distribution diagram



*Proof.* Recall that a point or line is called *absolute* (for  $\sigma$ ) if it is incident with its image (under  $\sigma$ ). We shall use  $\sim$  for adjacency in  $\Gamma$ , and  $\perp$  for collinearity in  $(X, L)$ .

**Step 1.** Each line contains a unique absolute point, and, dually, each point is on a unique absolute line.

(Indeed, if  $x$  is absolute, then  $x^\sigma$  is the only absolute line on  $x$ , and if  $x$  is not absolute then the unique line on  $x$  meeting  $x^\sigma$  is the only absolute line on  $x$ .)

**Step 2.** The set  $A$  of absolute points under  $\sigma$  is an ovoid in  $(X, L)$ . The graph  $\Gamma$  has  $v = q(q^2 + 1)$  vertices.

(Indeed, each  $l \in L$  meets  $A$  in a unique point. It follows that  $|A| = q^2 + 1$ . But  $|X| = (q + 1)(q^2 + 1)$ .)

**Step 3.**  $\Gamma$  is regular of valency  $q$ , and does not contain triangles. Adjacent vertices are non-collinear.

(Indeed, the neighbours of  $x$  are the  $q$  nonabsolute points of  $x^\sigma$ .)

**Step 4.**  $\Gamma$  does not have quadrangles, and any two vertices at distance 2 determine a unique pentagon. Two vertices have distance 2 if and only if they are collinear and the line joining them is non-absolute.

(Indeed, if  $x \sim y \sim z$ , then  $x$  and  $z$  are joined by the line  $y^\sigma$ . In particular,  $y$  is the only common neighbour of  $x$  and  $z$ . Let  $z \perp p \in x^\sigma$ . Then  $p \notin A$  because the unique absolute point on  $x^\sigma$  is collinear to  $x$ . Also the line  $l = zp$  is not absolute because  $z^\sigma$  passes through  $y$  and  $p \neq y$ . It follows that  $x \sim y \sim z \sim l^\sigma \sim p \sim x$  is the unique pentagon on  $x$  and  $z$ .)

Let us describe the distribution of vertices in  $\Gamma$  around a vertex  $x$ . Let  $m$  be the absolute line on  $x$ , and let  $x' = m^\sigma = x^\sigma \cap A$  be its absolute point. The vertex set of  $\Gamma$  is partitioned into the following seven parts:  $X_0 = \{x\}$ ,  $X_1 = x^\sigma \setminus A$ ,  $X_2 = x^\perp \setminus (A \cup m)$ ,  $X_5 = m \setminus (A \cup \{x\})$ ,  $X_{4a} = \{x'\}^\perp \setminus (A \cup m \cup x^\sigma)$ ,  $X_{4b} = \{y \in X \setminus A \mid y \sim z \in X_1 \text{ and } yz \text{ is absolute}\}$ , and  $X_3$ , consisting of the remaining points. Our aim is to show that  $X_i$  consists of the vertices at distance  $i$  from  $x$  in  $\Gamma$ , where  $X_{4a}$  and  $X_{4b}$  are distinguished by the fact that points in  $X_{4a}$  have neighbours in  $X_5$ . (Note however that for  $q = 2$  we have  $X_3 = \emptyset$ , and the graph  $\Gamma$  is the disjoint union of two pentagons. If  $p$  is in the relation  $4a$  to  $x$ , then  $x$  is in relation  $4b$  to  $p$ , i.e., relations  $4a$  and  $4b$  are paired, while the remaining relations are self-paired.)

**Step 5.** We have  $|X_0| = 1$ ,  $|X_1| = q$ ,  $|X_2| = q(q - 1)$ ,  $|X_3| = q(q - 1)(q - 2)$ ,  $|X_{4a}| = |X_{4b}| = q(q - 1)$ ,  $|X_5| = q - 1$ .

(Indeed, the claims are clear for  $X_i$  with  $i \leq 2$ . The only vertices that do not have distance 2 to some vertex of  $X_1$ , are the vertices that either are collinear to the point  $x' = x^\sigma \cap A$  (i.e., are in  $X_{4a} \cup X_5$ ), or are joined to a vertex on  $x^\sigma$  by an absolute line (i.e., are in  $X_{4b}$ ). The absolute line  $m$  on  $x$  contains  $q$  vertices,  $q - 1$  other than  $x$ , and none of them is collinear to a point in  $X_0 \cup X_1 \cup X_2$ , so these vertices have distance at least 5 to  $x$ . The vertices adjacent to some vertex in  $X_5$  are the  $q(q - 1)$  vertices of  $X_{4a}$ . The vertices of  $X_3$  are collinear to a unique vertex of  $x^\sigma$ , so this determines  $|X_3|$ .)

**Step 6.** Each vertex in  $X_3 \cup X_{4b}$  has a unique neighbour in  $X_{4a}$ .

(Indeed, let  $p \in X_3 \cup X_{4b}$ . Then  $p^\sigma$  does not pass through  $x'$  (since  $p \notin m$ , i.e.,  $p \notin X_0 \cup X_5$ ), so  $x'$  is collinear with a unique point  $z \in p^\sigma$ . The line  $x'z$  is not absolute (since  $z \notin m$  because  $p \notin X_{4a} \cup X_1$ ) and the point  $z$  is not absolute (since the line  $x'z$  contains only one absolute point), so  $z$  is the unique neighbour of  $p$  in  $X_1 \cup X_{4a}$ . Clearly  $z \in X_1$  iff  $p \in X_2$ .)

This proves everything claimed in the diagram. □

Now let us look at the special case where  $q = 2^{2e} + 1$  and  $(X, L)$  is the  $Sp(4, q)$  generalized quadrangle. The centralizer in  $Sp(4, q)$  of the polarity  $\sigma$  is the Suzuki group  $Sz(q)$  of order  $(q^2 + 1)q^2(q - 1)$ . This group is 2-transitive on



**Step 3.** If  $p \in X_2 \cup X_3 \cup X_{4a}$  then  $p$  has  $\frac{1}{2}q - 1$  neighbours in  $X_{3a}$ .

(Indeed, if  $p \in X_2 \cup X_3 \cup X_{4a}$ , then  $x \notin p^\sigma$ , and the plane  $\langle x, p^\sigma \rangle$  is a secant plane. In this plane, the point  $x$  is on one tangent, and on  $\frac{1}{2}q$  secants. One of these secants contains  $p'$ ; the remaining  $\frac{1}{2}q - 1$  contain each one neighbour of  $p$ .)

This determines the entire diagram. □

**Remark 2.3.** The graph  $\Gamma$ , and the fact that it is 2-arc transitive for  $Sz(q)$ , was found independently by Fang Xin Gui, a student of Cheryl Praeger.

**Remark 2.4.**  $\text{Aut } \Gamma$  is not primitive: the spread  $\{a^\sigma \mid a \in A\}$  is a system of blocks of imprimitivity. However,  $\text{Aut } \Gamma$  acts 2-transitively on the set of blocks, so that we do not find a nontrivial graph structure on the quotient.

**Remark 2.5.** Of course we also get finite heptagraphs (of valency  $q = 3^{2e} + 1$ ) starting from a generalized hexagon (of type  $G_2(q)$ ) with a polarity.

### 3 Addendum

The above was written in April 1992. In the meantime, Xin Gui Fang & C. E. Praeger [1, 2] appeared where the above graphs are found in the classification of certain 2-arc transitive graphs (and they refer to this work). As far as we know, the relation to generalized polygons with polarity still does not appear in the literature.

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Andries E. Brouwer  
*e-mail*: aeb@cwi.nl

Johannes Huizinga  
*e-mail*: J.Huizinga@dordrecht.nl