## A note on symmetric codes over GF(3)

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Let q be a prime power such that  $q \equiv 2 \pmod{3}$  and  $q \equiv 1 \pmod{4}$ , GF(q) a field of q elements and  $\mu$  the quadratic character of  $GF(q)^x$  with  $\mu(0)=0$ .

Let T be a matrix of degree q defined by  $T(a,b) = \mu(b-a)$ , where  $a, b \in GF(q)$ , and

$$S = \begin{pmatrix} 0 & 1 & \cdots & 1 & 1 \\ 1 & & & & \\ \vdots & & T & & \\ 1 & & & & \\ 1 & & & & \end{pmatrix}.$$

Let C(q) be the code generated by (I, S) over GF(3), which is introduced by V. Pless in [3] and I denotes the identity matrix of degree q+1.

The purpose of this note is to show that the minimum weight of C(q) is not smallet than  $\sqrt{q}$ .

§ 1. Let  $C^*(q)$  be the code generated by (I, S) over  $GF(3^2)$ . Let i be a primitive fourth root of unity in  $GF(3^2)$ . Then we may choose

$$\begin{pmatrix} -I-iS, & iI-S \\ -I+iS, & -iI-S \end{pmatrix}$$

as generators of  $C^*(q)$ , since (-S, I) is contained in C(q) (See [3]). We notice that -i(-I-iS)=iI-S and i(-I+iS)=-iI-S. Let U and L be the subcodes of  $C^*(q)$  generated by  $(-I-iS,\ iI-S)$  and  $(-I+iS,\ -iI-S)$  respectively. Then any codevector of  $C^*(q)$  has a form  $(x+y,\ -i(x-y))$ , where  $(x,\ -ix)\in U$  and  $(y,\ iy)\in L$ .

LEMMA 1. Let w denote the weight function. Then we have that

$$w(x+y, -i(x-y)) \ge w(x)$$
 and  $w(y)$ .

PROOF. We may label elements of  $GF(3^2)$  as follows:  $a_1=0$ ,  $a_2=1$ ,  $a_3=-1$ ,  $a_4=i$ ,  $a_5=i+1$ ,  $a_6=i-1$ ,  $a_7=-i$ ,  $a_8=-i+1$  and  $a_9=-i-1$ . Now

let  $n_i$  be the number of  $a_i$  in x and  $n_{ij}$  the number of  $a_j$  in the portion of y which corresponds to  $n_i$  coordinates of x giving  $a_i$   $(i, j = 1, \dots, 9)$ . Then we have that  $w(x) = n_2 + \dots + n_9$ ,  $w(y) = n_{12} + \dots + n_{19} + n_{22} + \dots + n_{29} + \dots + n_{92} + \dots + n_{99}$ ,  $w(x+y) = n_{12} + \dots + n_{19} + n_{21} + n_{22} + n_{24} + \dots + n_{29} + \dots + n_{91} + \dots + n_{94} + \dots + n_{96} + \dots + n_{99}$  and  $w(x-y) = n_{12} + \dots + n_{19} + n_{21} + n_{23} + \dots + n_{29} + \dots + n_{91} + \dots + n_{98}$ . Thus we have that  $w(x+y, -i(x-y)) = 2(n_{12} + \dots + n_{19}) + (2(n_{21} + \dots + n_{29}) - n_{22} - n_{23} + \dots + 2(n_{91} + \dots + n_{99}) - n_{95} - n_{99} \ge w(x)$  and w(y).

- -I-iS and -I+iS are equivalent under the field automorphism, and so they have the same minimum weight.
- § 2. The minimum weight of  $-I_1-iT$  is equal to or smaller than by one that of -I-iS, where  $I_1$  is the identity matrix of degree q. Now let G be a generalized quadratic residue code of J. H. van Lint and F. J. Mc-Williams [2] of GF(q) over  $GF(3^2)$ . Since  $(1, \dots, 1)$  belongs to both  $-I_1-iT$  and G, and since G=(-I-iT)+J, where J is the all 1 matrix of degree q, we see that  $-I_1-iT$  is a generator for G. Thus by a theorem of J. H. van Lint and F. J. McWilliams [2, Theorem 2, i] the minimum weight of -I-iT is at least  $\sqrt{q}$ .

Remark. The case  $q \equiv 3 \pmod{4}$  is treated in [1].

## **Bibliography**

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