

Research Article

A Modified p -Median Model for the Emergency Facilities Location Problem and Its Variable Neighbourhood Search-Based Algorithm

Menghao Xi,^{1,2} Feng Ye,¹ Zhong Yao,¹ and Qiuhong Zhao¹

¹ School of Economics and Management, Beihang University, Beijing 100191, China

² Department of Economics and Management, Institute of Disaster Prevention, East Beijing 101601, China

Correspondence should be addressed to Qiuhong Zhao; qhzhao@buaa.edu.cn

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Emergency incidents, including natural disasters, terrorist attacks, public health outbreaks, and industrial and mining accidents, and so forth, result in severe human casualties and property losses. Emergency facilities, which provide relief materials and services, play an important role in rescue management. The decision of where to locate the emergency rescue facilities is very important, as it determines the efficiency and effectiveness of the emergency management process. This paper develops a modified p -median problem model that accounts for rescue time limitations. A variable neighbourhood search- (VNS-) based algorithm is developed for the model considered. The modified VNS algorithm exhibits good performance on p -median benchmark problems. A case from Western China is studied, and a reasonable location decision is then made for emergency rescue facilities using the modified VNS algorithm. The paper also compares the results with and without considering the rescue time limitation.

1. Introduction

Disaster is a part of human history. Disasters, which can be divided into natural catastrophes and man-made disasters, have long threatened human survival and societal development. In recent years, many serious catastrophes have occurred throughout the world. On May 12, 2008, an 8.0-magnitude earthquake struck in Wenchuan, China; in the fall of 2009, a serious drought affected several provinces in southwest China, including Yunnan, Guizhou, and Guangxi; on March 11, 2011, a 9.0-magnitude earthquake occurred, triggering a tsunami that struck Fukushima, Japan, resulting in heavy casualties and a serious nuclear disaster. The acceleration of economic and industrial globalisation increases the likelihood of numerous types of disasters.

To reduce the losses caused by disasters, it is important to deliver rescue resources in a timely manner. The rescue resources, including mining equipment, tents, blankets, living supplies, heaters, and staffs, should be stored in the nearby the disaster areas. Therefore, time/distance is a key factor in emergency management. Meanwhile, to avoid waste, it is

important to site a suitable number of facilities and minimise transportation costs to ensure operational efficiency.

Given the importance of the choice of rescue facility locations in emergency management, this paper considers a modified p -median model that limits the maximum travel time to the demand points. While total transportation distance should be minimised, the number of facilities, denoted p in the p -median problem, is also a variable in the model considered here. A variable neighbourhood search- (VNS-) based algorithm is developed for the original p -median model and a modified model. Numerical studies demonstrate the good performance of the VNS algorithm. A case study from Western China is analysed. This analysis demonstrates that reasonable rescue facility locations can be identified using the proposed VNS algorithm.

The remainder of the paper is organised as follows. Section 2 reviews the literature related to emergency location problems. The p -median problem and the algorithms for p -median problem are then discussed. Both the original p -median model and the modified p -median model are presented in Section 3. A VNS-based algorithm is developed

TABLE 1: The computational results for the ORLIB pmed1–40 instances.

Number	n	p	Opt.	Result	Error (%)	Iterations
pmed1	100	5	5819	5819	0	28
pmed2	100	10	4093	4093	0	311
pmed3	100	10	4250	4250	0	355
pmed4	100	20	3034	3034	0	59691
pmed5	100	33	1355	1355	0	520295
pmed6	200	5	7824	7824	0	46
pmed7	200	10	5631	5631	0	364
pmed8	200	20	4445	4445	0	172840
pmed9	200	40	2734	2734	0	5851762
pmed10	200	67	1255	1255	0	19672781
pmed11	300	5	7696	7696	0	335
pmed12	300	10	6634	6634	0	9216
pmed13	300	30	4374	4374	0	1329065
pmed14	300	60	2968	2968	0	12960824
pmed15	300	100	1729	1729	0	1121054
pmed16	400	5	8162	8162	0	246
pmed17	400	10	6999	6999	0	5761
pmed18	400	40	4809	4809	0	5767360
pmed19	400	80	2845	2845	0	489788
pmed20	400	133	1789	1789	0	750728
pmed21	500	5	9138	9138	0	85
pmed22	500	10	8579	8579	0	7461
pmed23	500	50	4619	4619	0	36678
pmed24	500	100	2961	2961	0	32685
pmed25	500	167	1828	1828	0	3725367
pmed26	600	5	9917	9917	0	159
pmed27	600	10	8307	8307	0	2551
pmed28	600	60	4498	4498	0	36304
pmed29	600	120	3033	3033	0	5343570
pmed30	600	200	1989	1989	0	2146715
pmed31	700	5	10086	10086	0	57
pmed32	700	10	9297	9297	0	20694
pmed33	700	70	4700	4700	0	568705
pmed34	700	140	3013	3013	0	61781
pmed35	800	5	10400	10400	0	296
pmed36	800	10	9934	9934	0	15857
pmed37	800	80	5057	5057	0	1162436
pmed38	900	5	11060	11060	0	2430
pmed39	900	10	9423	9423	0	642
pmed40	900	90	5128	5129	0.019501	1149577
Average	—	—	—	—	0.0005	1575673

for the p -median models in Section 4. Numerical analyses are conducted in Section 5 including p -median benchmark instances. A case from Western China is studied in Section 6. Finally, Section 7 concludes the paper.

2. Literature Review

Location problems have been widely addressed in the literature. In the first study in this vein, Weber [1] studied the problem of determining the location of a warehouse to

minimise the total distance between the warehouse and the customers (which is called a Weber problem). Hakimi [2] found the optimum location of a “switching centre” in a communication network and located the best place to construct a “police station” in a highway system. Berman and Krass [3] considered a generalisation of the maximal cover location problem, which allows for partial customer coverage, with the degree of coverage being a nonincreasing step function of the distance to the nearest facility. The topic of emergency facility location has received substantial attention in recent

TABLE 2: The computational results for the fl1400 and pcb3038 instances.

Number	n	p	VNDS	Results	Error (%)	Iterations
fl1400	1400	10	101249.5	101249.50	$4.94E-05$	93
	1400	20	57857.55	57857.81	0.0004	5049
	1400	30	44086.53	44013.46	-0.1657	9614
	1400	40	35005.82	35002.55	-0.0093	16324
	1400	50	29089.78	29090.32	0.0019	11665
	1400	60	25166.15	25165.63	-0.0021	119233
	1400	70	22125.53	22126.02	0.0022	42814
	1400	80	19877.88	19876.73	-0.0058	66605
	1400	90	17987.94	17988.58	0.0036	30333
	1400	100	16551.2	16553.96	0.0167	220361
	1400	150	12032.65	12035.64	0.0248	131721
	1400	200	9360.01	9362.86	0.0304	188086
	1400	250	7742.7	7747.86	0.0666	572625
	1400	300	6624.52	6629.74	0.0788	406619
	1400	350	5727.02	5742.36	0.2679	870917
	1400	400	5020.5	5061.02	0.8071	247492
	Average	—	—	—	0.00698	183721
PCB3038	3038	10	1213082	1213082	$-8.2E-06$	249254
	3038	20	841349.1	840881.3	-0.0556	225928
	3038	30	680540.1	678613.8	-0.2830	629376
	3038	40	573407.4	573232.3	-0.0305	306695
	3038	50	507655.2	507965.3	0.0611	1960712
	3038	60	462232.9	461699.1	-0.1155	367235
	3038	70	428062.7	427791	-0.0635	260389
	3038	80	397990.3	397950.1	-0.0101	825286
	3038	90	373847	373806	-0.0110	12664323
	3038	100	353255.2	353800.1	0.1542	1449169
	3038	150	281772.1	282045.3	0.0970	3227972
	3038	200	238623	238973.3	0.1468	6881769
	3038	250	209343.3	209906.8	0.2691	854546
	Average	—	—	—	0.0012	2300204
rl5934	5934	10	9794951	9795160.966	0.0021	168012
	5934	20	6729282.5	6723997.648	-0.0785	3718
	5934	30	5405661.5	5390772.822	-0.2754	5973
	5934	40	4574374	4551995.656	-0.4892	19943
	5934	50	4053917.75	4048484.215	-0.1340	83309
	5934	60	3655898.75	3655516.508	-0.0105	30982
	5934	70	3353885	3352712.578	-0.0350	145637
	5934	80	3104877.75	3098857.868	-0.1939	6956913
	5934	90	2903895.25	2896253.012	-0.2632	6331784
	5934	100	2733817.25	2732673.696	-0.0418	237586
	5934	150	2151018.5	2153245.918	0.1036	1293267
	5934	200	1809064.38	1813514.215	0.2460	1196488
	5934	250	1571813.5	1573956.982	0.1364	3125323
	5934	300	1394715.12	1398625.701	0.2804	1037643
	Average	—	—	—	-0.0538	1474041

TABLE 3: The final solutions for the emergency facilities location problem.

The final solutions														p	Total distance (km)	Maximal distance (km)
19	I	E	3l	D	G	K	35	F	26	14	Q	1	N	14	353.6	14.5



FIGURE 1: The sites in the case study.

TABLE 4: The emergency facilities and the demand points they serve.

Emergency facilities	Demand points to serve
19	L, J, 19, 20
I	I, 15, 18, 16
E	E, 8, 7, 6
31	R, 31, 32, 33
D	D, 4, 3, 5, 2
G	G, 12, 11
K	K, 17, 22, 21
35	34, 35
F	F, 10, 9
26	P, 26, 27
14	13, 14, H
Q	Q, 28, 30, 29
1	C, O, A, B, 1
N	23, 25, M, N

decades. According to Jia et al. [4], emergency facility location problems can be divided into three types depending on the objective function of the location models: covering models,

p -median models, and p -centre models. The objective of coverage models is to provide “coverage” to demand points, and the earliest work on the subject is by Toregas et al. [5]. A demand point is only considered covered if a facility is available to service the demand point within a specified distance limit. The p -median problem, introduced by Hakimi [6], is to determine the location of p facilities to minimise the average (total) distance between demand points and facilities. Subsequently, ReVelle and Swain [7] formulated the p -median problem as a linear integer program and used a branch-and-bound algorithm to solve the problem. In contrast to p -median models, which concentrate on optimising the overall performance of the system, the p -centre model attempts to minimise the maximum distance between each demand point and its nearest facility, so the p -centre model is also called the minimax model.

The problem addressed in this paper is a revised version of the p -median problem, which minimises both the total distance and the number of the emergency facilities, under a constraint on the maximum amount of time required to access the demand points. The problem studied here is common in reality because both fixed and the variable costs should be minimised when establishing emergency facilities

TABLE 5: The final solutions for the emergency facilities location problem.

The final solutions														p	Total distance (km)	Maximal distance (km)
18	20	3	26	17	6	1	31	Q	34	9	G	23	14	14	343.6	20.9

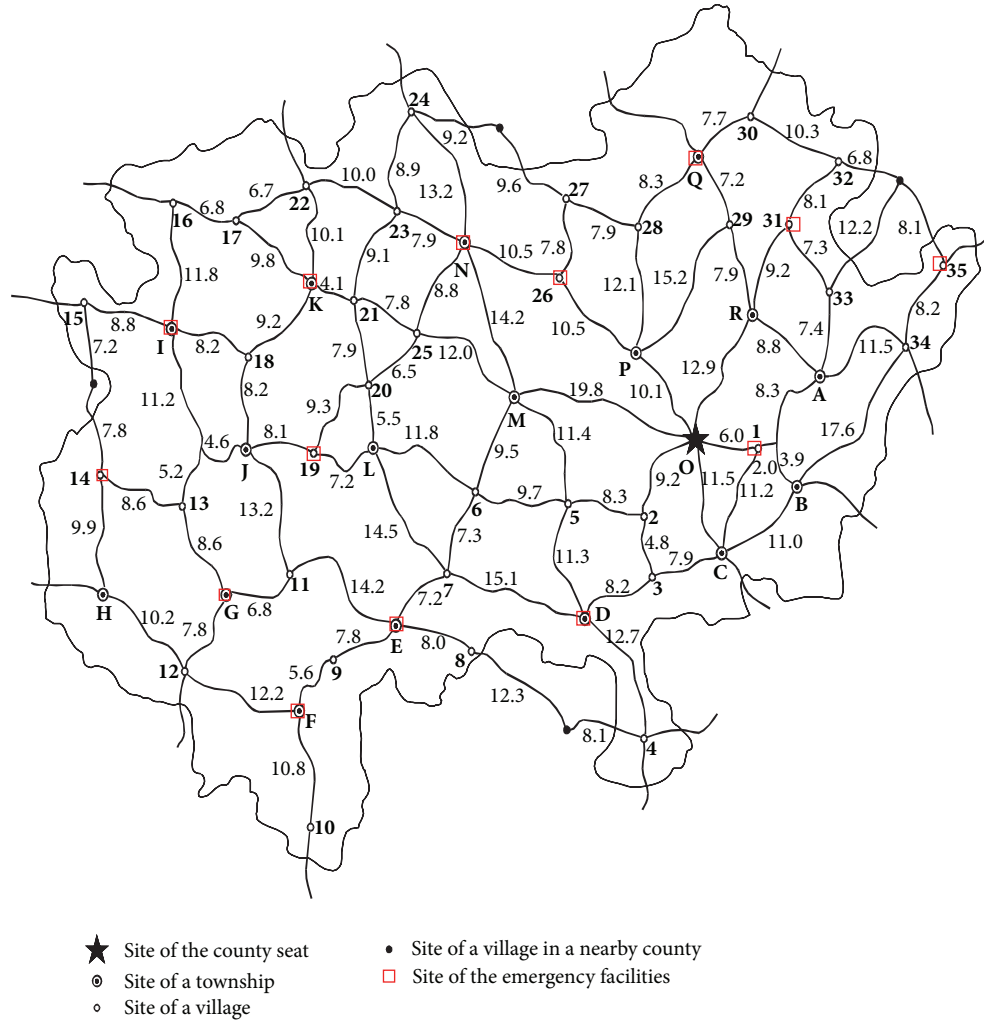


FIGURE 2: Locations of the emergency facilities in the case study with maximum time limitation.

while guaranteeing the efficiency of any possible rescue operation. In 1979, Kariv and Hakimi [8] demonstrated that the p -median problem is NP-hard. Since then, many heuristic algorithms have been proposed to solve the p -median problem and its various versions, including the genetic algorithm, ant algorithms, tabu search, artificial neural network.

The VNS algorithm initially proposed by Mladenović and Hansen [9] is among the most advanced methodologies for solving optimisation problems. Because its principle is simple and easy to understand and implement, various VNS-based algorithms have been successfully applied to many optimisation problems [10]. Hansen and Mladenović [11] were the first to apply VNS to the p -median problem. Hansen et al. [12] developed the variable neighbourhood decomposition search (VNDS) to solve the p -median problem. Mladenović et al. [13] summarised a variety of heuristic methods for the p -median problem and concluded that the VNS outperformed other algorithms.

3. The Models

3.1. The Original p -Median Problem Model. The original p -median problem model assumes that there are total of m potential sites, from which p candidate facilities are selected to serve n customers. Let M be the set of the potential m sites, N the set of the n customers, d_{ij} the distance from facility i to customer j , $i \in M$, and $j \in N$. It is generally assumed that $m = n$ and $M = N$; that is, the locations of the m customers are employed as the potential facility sites. The mixed integer model for the p -median problem is as follows:

$$\text{Min } z = \sum_{i \in M} \sum_{j \in M} d_{ij} x_{ij} \quad (1)$$

$$\text{s.t. } \sum_{i \in M} x_{ij} = 1, \quad \forall j \in M \quad (2)$$

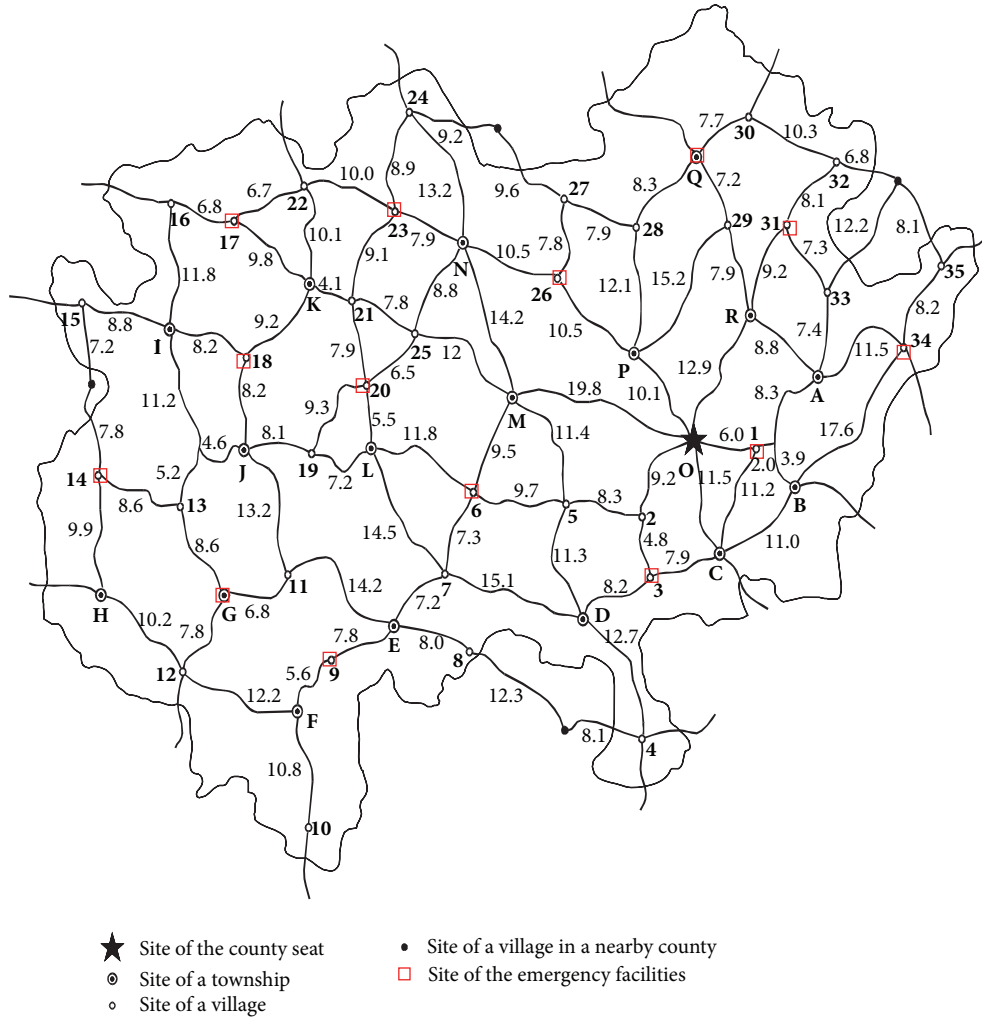


FIGURE 3: Locations of the emergency facilities in the case study with a maximum time limitation.

$$x_{ij} \leq y_i, \quad \forall i, j \in M \quad (3)$$

$$\sum_{i \in M} y_i = p \quad (4)$$

$$x_{ij}, y_j \in \{0, 1\} \quad \forall i, j \in M. \quad (5)$$

In the above model, $y_i = 1$ indicates that location i is selected as the service facility, $y_i = 0$ otherwise. When demand in location j needs to be served by facility i , $x_{ij} = 1$, $x_{ij} = 0$ otherwise. Therefore, the objective function in formulation (1) is to minimise total transportation distance from the service facilities to all demand points.

Constraint (2) ensures that any demand point is served by one and only one facility. Constraint (3) guarantees that whenever location i services a demand point j , it should be selected as the facility. Constraint (4) ensures that the total number of the service facilities equals p .

The mathematical model above was first proposed by ReVelle and Swain [7], and variants of this model have been developed by other researchers. For example, Church [14] studied the condensed Balinski constraints with the reduction of assignment variables using equivalent variable substitution; in 2008, Church [15] proposed a new

mathematical formula for the p -median problem: Both Exact and Approximate Model Representation (BEAMR).

Important modifications need to be made to adapt the p -median model to emergency management. When the average (total) distance decreases, the accessibility and effectiveness of the emergency facilities will increase. However, rescue time is also a key factor in emergency management, which is included in the modified p -median model below.

3.2. The p -Median Problem Model with a Maximum Time Limitation. In emergency management, it is very important that the rescue teams and rescue resources are able to access the demand point in time. Therefore, there should be some limitation to the maximum rescue time in contrast to the traditional p -median problem, which in turn means that p becomes a variable. Without loss of generality, it is assumed that the maximum time limitation is equivalent to the distance constraint. Denote L as the longest distance allowed (equivalent to the maximum time limitation); the corresponding mathematical model is as follows:

$$\text{Min } z = \alpha p + \sum_{i \in M} \sum_{j \in M} d_{ij} x_{ij} \quad (6)$$

Initialisation. Define the set of neighbourhood structures $N_k, k = 1, \dots, k_{\max}$, that will be used in the search; find an initial solution Y_0 ; set $Y \leftarrow Y_0$; and choose a stop condition. Repeat the following until the stop condition is met:
 (1) Set $k \leftarrow 1$; (2) Until $k = k_{\max}$ repeat the following steps:
 (a) Shaking. Generate a point Y' at random from the k th neighbourhood of Y ($Y \in N_k(Y)$);
 (b) Local search. Apply some local search method with Y' as the initial solution; denote with Y'' the local optimum that is obtained;
 (c) Move or not. If Y'' is better than Y , move there ($Y \leftarrow Y''$), and continue the search from $N_1(x)$ ($k \leftarrow 1$); otherwise, set $k \leftarrow k + 1$;

ALGORITHM 1: The basic VNS scheme.

TABLE 6: The emergency facilities and the demand points they serve.

Emergency facilities	Points to serve
18	I, K, J
20	19, 21, 25, L
3	2, D, C, 4
26	27
17	16, 22
6	7, 5, M
1	O, B, A
31	R, 32, 33
Q	30, 28, 29
34	35
9	F, 10, E, 8
G	12, 11
23	24, N
14	15, 13, H

$$\text{s.t. } \sum_{i \in M} x_{ij} = 1, \quad \forall j \in M \quad (7)$$

$$x_{ij} \leq y_i, \quad \forall i, j \in M \quad (8)$$

$$\sum_{i \in M} y_i = p \quad (9)$$

$$\max_{i \in M} \{d_{ij} x_{ij} \mid \forall j \in M\} < L \quad (10)$$

$$x_{ij}, y_j \in \{0, 1\} \quad \forall i, j \in M. \quad (11)$$

In the above model, there is an additional item αp in the objective function that is absent from formulation (1), while α can be considered as the fixed cost one of establishing an emergency facility. As mentioned above, p is a variable in formulation (6), as additional emergency facilities may be required due to the limitation on the maximum travel distance to any demand point. Consequently, α , the fixed cost of establishing one emergency facility, should be considered and initially set a sufficiently large value to minimise the number of emergency facilities. Therefore, in reality, α can be set to a very large value, such as 10,000,000.

The constraints in the modified p -median problem model are identical to those in the original p -median model except

that constraint (10) is added. Constraint (10) requires that the distance from any demand point to its nearest emergency facility be within the predefined distance L .

4. A Variable Neighbourhood Search-Based Algorithm

VNS is a type of metaheuristic devoted to the combinatorial problems, such as those described by the two models in Section 3. Briefly, VNS is a heuristic strategy that can be used to conduct an efficient stochastic search for better solutions through predefined and systematic changes of neighbourhoods. It implies the general principle of moving from smaller to larger sets to explore the neighbourhoods of the incumbent solution. Algorithm 1 presents the basic scheme of the VNS algorithm in Hansen et al. [12].

The VNS algorithm has also been applied to solve p -median problems. Hansen and Mladenović [11] proposed several VNS heuristics for the p -median problem and compared them to the greedy-plus-Interchange algorithm and two tabu search heuristics. García-López et al. [16] considered several parallelisation strategies of the VNS for the p -median problem and coded them in C using OpenMP. Crainic et al. [17] proposed a cooperative multisearch method for the VNS metaheuristic based on the central-memory mechanism. The p -median problem serves as a test case. The results indicate that, compared to sequential VNS, the cooperative strategy yields significant gains in terms of computation time without a loss in solution quality. Kochetov et al. [18] considered the p -median problem using a modified VNS algorithm. A new large neighbourhood for the graph partition problem is proposed. Computational experiments conducted by Kochetov et al. [18] show that the local improvement algorithm with the neighbourhood is fast and finds feasible solutions with little relative error. Hansen et al. [19] solved the clustering problem as a large-scale p -median model, using a new approach based on the VNS metaheuristic.

In this paper, we propose a modified VNS algorithm for the p -median problem model with a maximum time limitation. It can also be used to solve the original p -median problem model. The neighbourhood of the current solution in the proposed algorithm is determined by a stochastic rule. For example, more opportunities will be provided to the nearest neighbour if the current solution is simply being updated; otherwise, more distant neighbours may need to

Function EFLP

```

(1) set  $p = p_0$ , exist = false;
(2) randomize;
(3) initmap;
(4)  $C_1 = \text{initcluster}(p)$ ;
(5) while !exist do
(6)    $S = \text{pmp\_search}(p, C_1)$ 
(7)   if  $S \neq \phi$  then
(8)      $p := p - 1$ ;  $S' = S$ ;
(9)      $C_1 = \text{initcluster}(p)$ ;
(10)  else
(11)    exist = true;
end EFLP

```

ALGORITHM 2: The modified VNS algorithm for the p -median problem with a time limitation.

be explored. The proposed VNS scheme is embedded in the algorithm for the original p -median problem and the p -median problem model with a maximum time limitation. Accordingly, the modified VNS algorithm is illustrated in four parts. The first, called function **EFLP**, is the overall scheme of the modified algorithm (shown in Algorithm 2); the second, called function **pmp_search**, is a subfunction in **EFLP** that is used to find the global optimum for a given p (shown in Algorithm 3); the third and fourth parts, called function **LocalSearch** and function **ChangeCluster**, respectively, are subfunctions in **pmp_search**. The first part conducts the local search, and the second performs the variable neighbourhood search to change sites in the current emergency facilities set (shown, resp., in Algorithms 4 and 5).

In Algorithm 2, the function **EFLP** outputs the overall results of the modified algorithm, including the optimal number of the emergency facilities, the sites of these facilities, the demand clusters serviced by each facility, and the total travel distance. In Algorithm 2, p is initially at an initial value p_0 that is sufficiently large such that the solution S acquired in line 6 is feasible. The functions **randomize** and **initMap** initialise all of the sites (line 2 and line 3). In line 4, C_1 is the set of the facility sites, numbering p , obtained by randomly selecting p sites from the full set of available points. From line 5 to line 11, the algorithm repeatedly completes the steps to find the optimal solution S given p and C_1 , which is obtained through the function **pmp_search** (line 6), and the feasibility of the solution is then checked (line 7). The loop will continue with $p - 1$ if all demand points can be accessed within the maximum time limitation, and the current solution is saved as S' , which includes both the minimal total travel distance and the emergency facilities set, along with the allocation of the demand points to the emergency facilities (line 8), and the initial cluster with updated p is obtained for the next loop (line 9); otherwise, the algorithm stops, and the final solution is obtained.

Algorithm 3, called function **pmp_search**, is the core of the function **EFLP**. In **pmp_search**, the subfunctions **LocalSearch** and **ChangeCluster** are run first, while **LocalSearch** outputs the new emergency facilities set C'_1 and the overall travel distance f_1 by conducting a local search (see Algorithm 4 for detail). Next, **ChangeCluster** searches

the updated neighbourhood around the current solution C'_1 , while $f_1 \cdot \text{count}$ is the times that f_1 is unchanged. Algorithm 5 describes **ChangeCluster** function in detail.

In line 2 in Algorithm 3, REPEAT is the predefined value of repeating the steps from line 3 to line 15. In each loop, the local search around the current solution C_i is conducted first, with the output of C'_i . If the solution is updated and is within the maximum time limitation (i.e., $f_i \cdot \text{objv} < f_{i-1} \cdot \text{objv}$ and $f_i \cdot \text{maxlength} \leq L$), then the neighbourhood of C'_i is searched, and $f_i \cdot \text{count}$ is recorded as 1; otherwise, $f_i \cdot \text{count}$ is accumulated. If the accumulation is as large as RESTART, a predefined number, the search will start with a new initial solution (line 9); otherwise, the neighbourhood of C'_i/C_i is searched if a generated random number is smaller (not smaller) than Pc, a predefined decimal and $\text{Pc} < 0.5$. In line 15, the best results of the first i loops are recorded as min, and line 16 outputs the final optimum after the total REPEAT loops.

Algorithm 4, called function **LocalSearch**, conducts local search around solution C . In this algorithm, each demand point is first allocated to its nearest facility site $c_i \in C$, $i = 1, 2, \dots, p$, and the corresponding cluster set, T_i , $i = 1 \dots p$, which includes the facility site c_i and its service sites, is formed. For each T_i , $i = 1 \dots p$, an updated c_i that minimises the total length in T_i is searched. If C is modified, a new search of cluster set, T_i , $i = 1 \dots p$, is restarted; otherwise the local search is finished with the final output of c_i , T_i , $i = 1 \dots p$.

Algorithm 5, called function **changecluster**, conducts a neighbourhood search around the solution C . In the algorithm, the neighbourhood of the solution C is defined as the number of the emergency facilities to be moved from C , which is determined by equation $\text{CN} = \lceil \text{count} * \text{ChangeNum} / \text{RESTART} \rceil$ in line 1, while “count” is the number of times that C has remained unchanged (the maximal number of unchanged repetitions allowed is RESTART); ChangeNum is a predefined value that is selected randomly as 1 or 2 (extensive computational studies indicate that it is difficult to improve on the solution if ChangeNum is larger than 2). When CN is determined, both the sites to be chosen as emergency facilities and those to be removed from C are selected randomly. The loop repeats $\text{CN} + 1$ times. It can be seen from Algorithm 5 that although the neighbourhood of the solution C is not determined because ChangeNum is selected randomly from 1 and 2, it is more likely that more sites (a larger neighbourhood) will be selected if the value of “count” increases.

As can be observed from Algorithm 2 and the subfunctions (Algorithms 3 to 5), the proposed modified VNS algorithm successfully combines the variable neighbourhood search scheme with some random mechanics. However, all of the necessary steps can be implemented easily.

5. Numerical Studies

Algorithm 2, which is designed for the emergency rescue facilities location problem, originated from the algorithm developed for the p -median problem. Accordingly, mechanics employed by the two algorithms (i.e., for the original p -median problem and for the p -median problem with


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Function pmp_search ( $p, C_1$ )
(1)  $(f_1, C'_1) = \text{localsearch}(C_1)$ ;  $f_1 \cdot \text{count} = 1$ ;  $C_2 = \text{changecluster}(C'_1, f_1 \cdot \text{count})$ ;
(2) for  $i = 2$  to REPEAT do
(3)    $(f_i, C'_i) = \text{localsearch}(C_i)$ ;
(4)   if  $f_i \cdot \text{objv} < f_{i-1} \cdot \text{objv}$  and  $f_i \cdot \text{maxlength} \leq L$  then
(5)      $f_i \cdot \text{count} = 1$ ;  $C_{i+1} = \text{changecluster}(C'_i, f_i \cdot \text{count})$ ;
(6)   else
(7)      $f_i \cdot \text{count} := f_i \cdot \text{count} + 1$ ;
(8)     if  $f_i \cdot \text{count} = \text{RESTART}$  then
(9)        $C_{i+1} = \text{initcluster}$ ;
(10)    else
(11)      if rand < Pc then
(12)         $C_{i+1} = \text{changecluster}(C'_i, f_i \cdot \text{count})$ ;
(13)      else
(14)         $C_{i+1} = \text{changecluster}(C_i, f_i \cdot \text{count})$ ;
(15)   $\min = \min_i \{f_i \cdot \text{objv}\}$ 
(16)  return  $(f_{\min}, C'_{\min})$ .
end pmp_search

```

ALGORITHM 3: The modified VNS algorithm for the p -median problem.**Function** localsearch (C)

```

(1)  $T_i = \Phi$ ;  $i = 1 \dots p$ 
(2) for each  $s \in M$ 
(3)   find the shortest edge  $(s, c)$ ,  $c \in C$ 
(4)    $T_i = T_i \cup s$ 
(5) for each  $T$ 
(6)   find the best  $c$  to make the total length in  $T_i$  shortest
(7) if  $C$  has no change then exit else Go to (2)
end localsearch

```

ALGORITHM 4: The modified VNS algorithm for the p -median problem.**Function** changecluster (C, count)

```

(1) CN =  $\lceil \text{count} * \text{ChangeNum} / \text{RESTART} \rceil$ ;
(2) for  $i = 0$  to CN do
(3)   random choose  $v$  from  $V$ , but  $v \notin C$ 
(4)    $C_{\text{random}} = v$ 
end changecluster

```

ALGORITHM 5: The modified VNS algorithm for the p -median problem.

time limitation) to determine the global optimisation solution with a given p , which is the **pmp_search** function in Algorithm 2, are identical. In this section, the benchmark instances ORLIB pmed1-40, fl1400 and pcb3038 are used to evaluate the performance of the proposed algorithm. The CPU of the computer used for computation is 2.1GHz, and the code is written in C. The parameters in the function **pmp_search** are as follows: REPEAT = 50,000,000, RESTART = 80,000, and Pc = 0.3.

As in Hansen et al. [12], the relative error = $((f - f^*)/f^*) * 100\%$ is used for evaluation, where f is the objective

value obtained by the proposed algorithm and f^* is the optimal objective value under comparison. In Tables 1 and 2, the results generated by the **pmp_search** function (i.e., the proposed p -median algorithm) are compared to the optimal solutions of the ORLIB pmed1-40 and the results for fl1400, pcb3038, and rl5934 reported in Hansen et al. [12], respectively.

In Table 1, the computational results gained by the proposed p -median algorithm for the ORLIB pmed1-40 are compared to the optimal solutions. Table 1 demonstrates that the proposed p -median algorithm can obtain the optimal solutions, except the solution for the last instance. However, the result of the proposed model is only 0.019501% higher than the optimum. Additionally, to demonstrate the computational efficiency of the proposed algorithm, the last column of Table 1 lists the number of iterations required to obtain the final results. Table 1 shows that p is the main factor in the global search. When p is small, the function **pmp_search** is highly efficient. For pmed1-40 instances, 1,000,000 iterations require approximately 10 minutes.

Table 2 compares the computational results for the fl1400, pcb3038, and rl5934 instances with those in Hansen et al. [12]. While some of the results obtained by Hansen et al. [12] are updated by the proposed algorithm, the two algorithms exhibit similar performance with respect to computational results. See the last row for each data set. Moreover, as the proposed algorithm is simple to realise and outputs the final results in a short period of time, overall, it outperforms the algorithm in Hansen et al. [12]. The same is true of the results in Table 1.

6. A Case Study

The real case studied here comes from a county in Western China. In this county, there are a total of 18 townships and 35 villages for a total of 54 sites with the inclusion of the county seat. Figure 1 plots the sites, along with the routes and the distances between the sites. In Figure 1, the county seat

is denoted 0, the townships are denoted $A, B, \dots, R$, and the villages are denoted $1, 2, \dots, 35$.

6.1. Solution with Maximum Time Limitation. The objective of this study is to determine the minimum number of emergency facilities such that the total distance from the emergency facilities to the demand points are minimised, and the distance from any emergency facility to any demand point it serves is less than the restriction of 15 km.

This case is solved using the algorithm proposed in Section 4. In the final solution, a total of 14 sites are selected for the emergency facilities, as shown in the left column in Table 3. Figure 2 plots the locations of these emergency facilities. Table 3 indicates that the total distance in the solution is 353.6 kilometres, while the maximal distance from any emergency facility to any demand location it services is 14.5 km. In Table 4, we list the emergency facilities and the demand points they serve. In addition, the case was solved with the exact method using CPLEX. There is no deviation from the optimum.

6.2. Solution without a Maximum Time Limitation. For comparison, in this subsection, no maximum time limitation is assumed. The computational results are listed in Tables 5 and 6, while the locations of the emergency sites are plotted in Figure 3.

The results shown in Table 5 indicate that compared to the results in Table 3, the total distance in the solution declines, while the maximal distance from any emergency facility to any demand point it services increases. While the value of p is identical in these two cases, the sites of these emergency facilities differ, as do the demand points they serve. See Tables 4 and 6 for details.

7. Conclusions

In this paper, we consider an emergency facilities locations problem. The location of emergency rescue facilities plays an important role in rescue management because it determines the efficiency and effectiveness of the emergency management process. A reasonable location decision can, on the one hand, improve the efficiency and effectiveness of emergency management and, on the other, reduce both the fixed and variable costs of establishing these facilities and conducting emergency rescue operations. This paper proposes a modified p -median problem model capable of considering a rescue time limitation. A VNS-based algorithm is developed for the considered model. The VNS algorithm exhibits good performance on p -median benchmark problems. A case from Western China is studied, and the VNS algorithm determined a reasonable set of emergency rescue facility locations.

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