## AN APPLICATION OF LINEAR PROGRAMMING TO PERMUTATION GROUPS

## W. H. MILLS

Let  $S_N$  denote the symmetric group acting on a finite set X of N elements,  $N \ge 3$ . Let  $\sigma$  and  $\tau$  be elements of  $S_N$ . In a previous paper [1] the following question was raised: If  $\sigma$  and  $\tau$  commute on most of the points of X, does it necessarily follow that  $\tau$  can be approximated by an element in the centralizer  $C(\sigma)$  of  $\sigma$ ?

We define a distance  $D(\sigma, \tau)$  between two elements  $\sigma$  and  $\tau$  in  $S_N$  to be the number of points g in X such that  $g\sigma \neq g\tau$ . (This differs from the distance  $d(\sigma, \tau)$  defined in [1] by a factor of N.) Then  $D(\sigma\tau, \tau\sigma)$  is the number of points in X on which  $\sigma$  and  $\tau$  do not commute. Let  $D_{\sigma}(\tau)$  denote the distance from  $\tau$  to the centralizer  $C(\sigma)$  of  $\sigma$  in  $S_N$ . Thus

$$D_{\sigma}(\tau) = \min_{\lambda \in \mathcal{O}(\sigma)} D(\tau, \lambda)$$
.

It will be shown that the determination of  $D_{\sigma}(\tau)$  is equivalent to the optimal assignment problem in linear programming.

The question raised in [1] can be phrased thus: If  $D(\sigma\tau, \tau\sigma)$  is small, is  $D_{\sigma}(\tau)$  necessarily small? If  $\sigma$  is not the identity we set

$$D_{\sigma} = \max_{\substack{ au \in \mathcal{G}(\sigma)}} D_{\sigma}( au) / D(\sigma au, au \sigma)$$
 .

Now  $D_{\sigma}$  is large unless  $\sigma$  is the product of many disjoint cycles, most of which have the same length. Some examples of this are worked out in detail in [1]. This leads us to study the case where  $\sigma$  is the product of m disjoint cycles of length n, where N = nm and m is large. In [1] it was shown that if  $m \geq 2$ , then

- (a) if n is even, then  $D_{\sigma} = n/4$ , and
- (b) if n is odd,  $n \ge 3$ , then
  - $(n-1)/4 \leq D_{\sigma} \leq n/4.$

In the present paper it is shown that if n is odd,  $n \ge 3$ , and  $m \ge n-2$ , then

$$D_{\sigma} = (n-1)^2/(4n-6)$$
.

1. Relation to linear programming. Let  $\sigma$  be an arbitrary element of the symmetric group  $S_N$ . We write  $\sigma$  as the product of disjoint cycles:

$$\sigma = C_1 C_2 \cdots C_m$$
 ,

Received October 22, 1962.

where  $C_i$  is a cycle of length  $n_i$ , and every point left fixed by  $\sigma$  is counted as a cycle of length 1. Then

$$n_1 + n_2 + \cdots + n_m = N$$
.

Let  $g_i$  be a fixed element of the cycle  $C_i$ ,  $1 \leq i \leq m$ . Then every element of the underlying set X is of the form  $g_i\sigma^a$ , where  $1 \leq i \leq m$  and  $0 \leq a < n_i$ .

Let  $\lambda$  be an element of  $C(\sigma)$ , the centralizer of  $\sigma$  in  $S_N$ . Then since

$$(g_i\sigma^a)\lambda=(g_i\lambda)\sigma^a$$
 ,

it follows that  $\lambda$  is determined by its effect on the  $g_i$ , and that  $\lambda$  permutes the cycles  $C_i$ . Let  $\overline{\lambda}$  be the permutation of 1, 2,  $\cdots$ , *m* such that  $i\overline{\lambda} = j$  if  $\lambda$  maps  $C_i$  onto  $C_j$ . We will call a permutation  $\alpha$  in  $S_m$  admissible if  $\alpha = \overline{\lambda}$  for some  $\lambda \in C(\sigma)$ . It is easy to see that  $\alpha$  is admissible if and only if  $n_i = n_{i\alpha}$ ,  $1 \leq i \leq m$ . Let A denote the group of all admissible permutations.

Let  $\tau$  be a second element of  $S_N$ . We wish to determine

$$D_{\sigma}( au) = \min_{\lambda \in \mathcal{O}(\sigma)} D( au, \lambda)$$
 ,

where  $D(\tau, \lambda)$  is the number of points g in X such that  $g\tau \neq g\lambda$ . Let  $E(\tau, \lambda)$  denote the number of points h in X such that  $h\tau = h\lambda$ , and set

$$E_{\sigma}(\tau) = \max_{\lambda \in \mathcal{O}(\sigma)} E(\tau, \lambda)$$
.

Then

$$D_{\sigma}( au) = N - \max_{\lambda \in \sigma(\sigma)} E( au, \lambda) = N - E_{\sigma}( au)$$
 .

We shall show that the determination of  $E_{\sigma}(\tau)$  is equivalent to the optimal assignment problem in linear programming.

The elements  $\lambda$  in  $C(\sigma)$  are the permutations of the form

$$(g_i \sigma^a) \lambda = g_{i a} \sigma^{a+r_i}$$
,  $1 \leq i \leq m, \, 0 \leq a < n_i$  ,

where  $\alpha$  is admissible and  $r_1, r_2, \dots, r_m$ , are integers. Moreover

$$E( au, \lambda) = \sum_{i=1}^{m} F_i(r_i, ilpha)$$
 ,

where  $F_i(r, j)$  is the number of solutions of

$$(1) \qquad \qquad (g_i \sigma^x) \tau = g_j \sigma^{x+r}, \ 0 \leq x < n_i \ .$$

Set

$$b_{ij} = egin{cases} 0 & ext{if} \;\; n_i 
eq n_j \ \max_r F_i(r,j) & ext{if} \;\; n_i = n_j \;. \end{cases}$$

Thus  $b_{ij}$  is the maximum number of points of  $C_i$  on which an element  $\lambda$  in  $C(\sigma)$ , that maps  $C_i$  onto  $C_j$ , can agree with  $\tau$ . We have

$$E_{\sigma}(\tau) = \max_{\lambda \in \mathcal{O}(\sigma)} E(\tau, \lambda) = \max_{\alpha \in A} \max_{r_1 \cdots r_m} \sum_{i=1}^m F_i(r_i, i\alpha)$$
,

or

(2) 
$$E_{\sigma}(\tau) = \max_{\alpha \in A} \sum_{i=1}^{m} b_{i,i\alpha} .$$

Now let  $\beta$  be an arbitrary permutation of  $1, 2, \dots, m$ . There is an  $\alpha \in A$  such that  $i\alpha = i\beta$  for all *i* such that  $n_i = n_{i\beta}$ . Therefore, since  $b_{ij} = 0$  if  $n_i \neq n_j$ , it follows that we can take the maximum in (2) over the entire symmetric group  $S_m$  instead of over the subgroup A. Thus

(3) 
$$E_{\sigma}(\tau) = \max_{\beta \in S_m} \sum_{i=1}^m b_{i,i\beta}.$$

The determination of a maximum of the form (3) is the optimal assignment problem in linear programming—ordinarily expressed in terms of m individuals to be assigned to m jobs, where  $b_{ij}$  is a measure of how well the *i*th individual can do the *j*th job. (See [2]; or [3], pp. 131-136.) Von Neumann [2] has shown that this problem is equivalent to a certain zero-sum two-person game.

The equality (3) can be rewritten in the form

$$(4) E_{\sigma}(\tau) = \max_{P} \sum_{i,j} e_{ij} b_{ij},$$

where P is the set of all  $m \times m$  permutation matrices  $(e_{ij})$ . The set P is clearly a subset of the set R of all real  $m \times m$  matrices  $(y_{ij})$  such that

(5) 
$$y_{ij} \geqq 0, 1 \leqq i, j \leqq m$$
 ,

(6) 
$$\sum_{i=1}^m y_{ij} = 1$$
 ,  $1 \leq j \leq m$  ,

and

(7) 
$$\sum_{j=1}^m y_{ij} = 1$$
 ,  $1 \leq i \leq m$  .

The matrices of the set R form a convex bounded subset of real  $m^2$ dimensional Euclidean space, whose vertices are the permutation matrices. (This result is due to Garrett Birkhoff. See [2], pp. 8-10.) It follows that

$$E_{\sigma}( au) = \max_{P}\sum_{i,j}e_{ij}b_{ij} = \max_{R}\sum_{i,j}y_{ij}b_{ij}$$
 .

It is now clear that the determination of  $E_{\sigma}(\tau)$  is actually a problem in linear programming. It is easy to see that the equalities (6) and (7) can be replaced by inequalities (see [2], Lemma 1). Thus if Y is the set of all real  $m \times m$  matrices  $(y_{ij})$  satisfying (5),

(8) 
$$\sum_{i=1}^m y_{ij} \leq 1$$
,  $1 \leq j \leq m$ ,

and

(9) 
$$\sum_{j=1}^m y_{ij} \leq 1$$
,  $1 \leq i \leq m$ ,

then

$$E_{\sigma}( au) = \max_{r} \sum\limits_{i,j} y_{ij} b_{ij}$$
 .

For our purposes this is the most useful formulation of the problem.

2. Blocks. By a block of length  $s, s \ge 1$ , we mean a set of the form  $g\sigma, g\sigma^2, \dots, g\sigma^s$ , such that  $\sigma$  and  $\tau$  commute on  $g\sigma, g\sigma^2, \dots, g\sigma^{s-1}$ , but do not commute on g and  $g\sigma^s$ . The length of a block B will be denoted by |B|. If  $\sigma$  and  $\tau$  commute on every point of the cycle  $C_i$ , then we say that  $\sigma$  and  $\tau$  commute on  $C_i$ . In this case the cycle  $C_i$  contains no blocks. On the other hand if  $C_i$  contains exactly q points on which  $\sigma$  and  $\tau$  do not commute,  $q \ge 1$ , then  $C_i$  consists of exactly q blocks, and each point of  $C_i$  belongs to one and only one block. Now  $D(\sigma\tau, \tau\sigma)$  is the number of points in X on which  $\sigma$  and  $\tau$  do not commute. It follows that  $D(\sigma\tau, \tau\sigma)$  is equal to the total number of blocks in all cycles.

If  $\sigma$  and  $\tau$  commute on the points  $g, g\sigma, g\sigma^2, \dots, g\sigma^a$ , then it follows, by induction on a, that

$$(g\sigma^{
u}) au=(g au)\sigma^{
u}$$
 ,  $0\leq 
u\leq a+1$  .

In particular if  $\sigma$  and  $\tau$  commute on the cycle  $C_i$ , and if  $g_i \tau = g_j \sigma^r$ , then

$$g_i\sigma^x au=g_i\sigma^{r+x}$$

for all x. Therefore, in this case, the number of solutions  $F_i(r, j)$  of (1) is  $n_i$ , so that  $b_{ij} = n_i = n_j$ .

Now let  $C_i$  be a cycle on which  $\sigma$  and  $\tau$  do not commute. Then

 $C_i$  is composed of one or more blocks. Let B be one of the blocks of  $C_i$ , and let B consist of the points

$$g_i\sigma^b, g_i\sigma^{b+1}, \cdots, g_i\sigma^{b+s-1}$$

Then |B| = s. Let  $g_i \sigma^b \tau = g_j \sigma^{b+r}$ . Since  $\sigma$  and  $\tau$  commute on  $g_i \sigma^{b+\mu}$ ,  $0 \leq \mu \leq s-2$ , we have

$$g_i \sigma^{b+
u} au = g_i \sigma^{b+r+
u}$$
 ,  $0 \leq v \leq s-1$ 

In particular  $n_j \geq s$ . Moreover if  $n_i = n_j$ , then the number of solutions  $F_i(r, j)$  of (1) is at least s, and hence  $b_{ij} \geq s$ . It follows that if  $n_i = n_j$ , then  $b_{ij}$  is at least the length of the longest block of  $C_i$  that  $\tau$  maps into  $C_j$ .

Moreover since  $\sigma$  and  $\tau$  do not commute on  $g_i \sigma^{b+s-1}$ , we have

$$g_i \sigma^{\scriptscriptstyle b+s} au 
eq g_i \sigma^{\scriptscriptstyle b+s-1} au \sigma = g_j \sigma^{\scriptscriptstyle b+r+s}$$
 .

In particular if  $C_i$  consists of the single block B, then  $s = n_i$ , and

$$g_j\sigma^{\scriptscriptstyle b+r}=g_i\sigma^{\scriptscriptstyle b} au=g_i\sigma^{\scriptscriptstyle b+s} au
eq g_j\sigma^{\scriptscriptstyle b+r+s}$$
 .

It follows that  $s \neq n_j$ . Therefore we must have  $n_j > s = n_i$ . Thus if  $C_i$  consists of a single block B, then  $\tau$  maps B into a cycle  $C_j$  such that  $n_j > n_i$ . This is a generalization of a result noted in [1]: If the cycles  $C_i$  all have the same length, then no cycle can consist of a single block.

3. The case n odd. We now restrict ourselves to the case where  $\sigma$  is the product of m cycles of the same length n, n > 1, N = mn,  $N \ge 3$ . Thus we have  $n_1 = n_2 = \cdots = n_m = n$ , and every permutation in  $S_m$  is admissible, so that  $A = S_m$ . Set

$$D_{\sigma} = \max_{\tau \notin \mathcal{G}(\sigma)} \left\{ D_{\sigma}(\tau) / D(\sigma \tau, \tau \sigma) 
ight\}$$
 .

It was shown in [1] that if n is even and  $m \ge 2$ , then  $D_{\sigma} = n/4$ . We now show that if n is odd and  $m \ge n-2$ , then  $D_{\sigma} = (n-1)^2/(4n-6)$ . Without loss of generality we can take X to be the set of the first N positive integers, and

$$\sigma = (1, 2, \cdots, n)(n+1, \cdots, 2n) \cdots (N-n+1, \cdots, N) .$$

Thus for g in X we have

$$g\sigma = egin{cases} g+1 & ext{if} \ n 
eq g \ , \ g+1-n & ext{if} \ n \mid g \ . \end{cases}$$

We let  $C_i$  denote the *i*th cycle:

$$C_i = (in - n + 1, in - n + 2, \dots, in)$$
.

We must show that

$$\max_{\substack{\tau \notin \mathcal{O}(\sigma)}} \left\{ D_{\sigma}(\tau) / D(\sigma \tau, \tau \sigma) 
ight\} = (n-1)^2 / (4n-6) \; .$$

We break up the proof into two lemmas.

LEMMA 1. If n is odd and  $m \ge n-2$ , then there exists a  $\tau \in S_N$ ,  $\tau \notin C(\sigma)$ , such that

$$D_{\sigma}(\tau)/D(\sigma\tau, \tau\sigma) = (n-1)^2/(4n-6)$$
.

*Proof.* Suppose first that n = 3. Then

$$\sigma = (123)(456) \cdots (N-2, N-1, N)$$
 .

Here we take  $\tau = (12)$ . Then  $\sigma \tau \sigma^{-1} \tau^{-1} = (132)$ , so that  $\sigma$  and  $\tau$  commute on all but three points, and  $D(\sigma \tau, \tau \sigma) = 3$ . Moreover

$$b_{ij} = egin{cases} 0 & ext{if} \ i 
eq j \ , \ 1 & ext{if} \ i = j = 1 \ , \ 3 & ext{if} \ i = j > 1 \ . \end{cases}$$

Hence

$$E_{\sigma}( au) = \max_{_{P}} \sum_{_{i,j}} e_{_{ij}} b_{_{ij}} = \sum_{_{i=1}}^{^{m}} b_{_{ii}} = 3m-2 = N-2$$
 .

Therefore  $D_{\sigma}(\tau) = N - E_{\sigma}(\tau) = 2$ , and

$$D_{\sigma}( au)/D(\sigma au, au\sigma) = 2/3 = (n-1)^2/(4n-6)$$
 .

We can now suppose that  $n \ge 5$ . Set n = 2K + 1. Then  $K \ge 2$ , and  $m \ge 2K - 1$ . Set  $\tau = \tau_1 \tau_2 \cdots \tau_K$ , where

$$egin{array}{ll} au_r = (r,\,n+r,\,2n+r,\,\cdots,\,Kn-n+r,\,K+r\,\,,\ Kn+r,\,Kn+n+r,\,\cdots,\,2Kn-2n+r)\,. \end{array}$$

Thus for g in X we have

$$g au = egin{cases} g+n ext{ if } g=pn+r, 0 \leq p \leq K-2, 1 \leq r \leq K\,, \ K+r ext{ if } g=Kn-n+r, 1 \leq r \leq K\,, \ Kn+r ext{ if } g=K+r, 1 \leq r \leq K\,, \ g+n ext{ if } g=pn+r, K \leq p \leq 2K-3, 1 \leq r \leq K\,, \ r ext{ if } g=2Kn-2n+r, 1 \leq r \leq K\,, \ g ext{ otherwise }. \end{cases}$$

C<sub>K+1</sub>  $C_{2K-1}$  $C_2$ Сĸ n+1Kn-n+1Kn+12Kn - 2n + 11 to to to to to 2Kn - 2n + KKn - n + Kĸ n + KKn + KK + 1 K+1 to 2KFigure 1

The blocks of  $\tau$  are shown schematically in Figure 1.

The permutation  $\tau$  maps the shaded blocks of Figure 1 onto themselves, and it maps the other blocks as indicated by the arrows. The permutations  $\sigma$  and  $\tau$  commute on the cycles  $C_i$  with  $i \ge 2K$ . Hence these cycles contain no blocks and are not shown in the figure. Let c denote the number of cycles on which  $\sigma$  and  $\tau$  commute. Thus c =m - (2K - 1). The number of points on which the identity I agrees with  $\tau$  is

$$E(\tau, I) = cn + 1 + (2K - 2)(K + 1)$$

Clearly I belongs to  $C(\sigma)$ . On the other hand suppose that  $\lambda$  is an arbitrary element of  $C(\sigma)$ . If there exists a cycle  $C_i$  such that  $\tau$  and  $\lambda$  do not agree on any points of  $C_i$ , then

$$E(\tau, \lambda) \leq cn + (2K-2)(K+1)$$
.

If  $\tau$  and  $\lambda$  agree on the point *n*, then

$$E( au, \lambda) \leq cn + 1 + (2K - 2)(K + 1)$$
.

If  $\tau$  and  $\lambda$  do not agree on n, and if  $\tau$  and  $\lambda$  agree on at least one point of every cycle  $C_i$ , then there are at least K - 1 blocks of length K + 1 on which  $\tau$  and  $\lambda$  do not agree. Hence in this case

$$E( au,\lambda) \leq cn + (K-1)(K+1) + K^2 \ = cn + 1 + (2K-2)(K+1) \; .$$

Therefore

$$egin{aligned} E_{\sigma}( au) &= \max_{\lambda \in \mathcal{O}(\sigma)} E( au, \lambda) = E( au, I) = cn + 1 + (2K - 2)(K + 1) \ &= (m - 2K + 1)n + 2K^2 - 1 = N - 2K^2. \end{aligned}$$

Hence

$$D_{\sigma}( au) = N - E_{\sigma}( au) = 2K^2 = rac{1}{2}(n-1)^2$$
 .

We see from Figure 1 that the total number of blocks is

$$2(2K-2) + 3 = 2n - 3$$
.

Since this is equal to  $D(\sigma\tau, \tau\sigma)$ , we have

$$D_{\sigma}(\tau)/D(\sigma\tau, \tau\sigma) = (n-1)^2/(4n-6)$$
.

This proves the lemma.

Lemma 1 establishes that  $D_{\sigma} \ge (n-1)^2/(4n-6)$  if n is odd and  $m \ge n-2$ . Our other lemma, which establishes the opposite inequality, does not depend on the size of m.

LEMMA 2. If n is odd and 
$$\tau \in S_N$$
,  $\tau \notin C(\sigma)$ , then  
 $D_{\sigma}(\tau)/D(\sigma\tau, \tau\sigma) \leq (n-1)^2/(4n-6)$ .

*Proof.* As before we set n = 2K + 1. Let c denote the number of cycles  $C_i$  on which  $\sigma$  and  $\tau$  commute, and let  $Q_s$  denote the total number of blocks of length s. Since the cycles  $C_i$  all have the same length n, it follows from the last paragraph of §2 that there are no blocks of length n. Hence

$$D(\sigma au, au\sigma)=\sum\limits_{s=1}^{n-1}Q_s$$
 ,

since this sum is equal to the total number of blocks. Set

$$G(\tau) = N - rac{(n-1)^2}{4n-6} \sum_{s=1}^{n-1} Q_s$$
.

The desired result holds if and only if

$$E_{\sigma}(\tau) \geq G(\tau)$$
.

By §1 it is sufficient to show that there exists a real  $m \times m$  matrix  $(y_{ij})$  satisfying (5), (8), (9) and

(10) 
$$\sum_{i,j} y_{ij} b_{ij} \ge G(\tau) .$$

Case 1.

$$cn+\sum\limits_{s=1}^{n-1}s^{2}Q_{s}/n \geq G( au)$$
 .

In this case we set  $y_{ij} = n_{ij}/n$ , where  $n_{ij}$  is the number of points of  $C_i$  which are mapped into  $C_j$  by  $\tau$ . Now (5), (6) and (7) hold for this choice of  $(y_{ij})$ . Hence (8) and (9) also hold.

Suppose  $C_i$  is a cycle on which  $\sigma$  and  $\tau$  commute. Suppose  $\tau_i$  maps  $C_i$  onto the cycle  $C_z$ . Then

$$y_{ij} = egin{cases} 1 & ext{if} \ j = z \ , \ 0 & ext{if} \ j 
eq z \ . \end{cases}$$

Moreover  $b_{iz} = n$  by § 2. Hence

$$\sum_{j=1}^m y_{ij}b_{ij} = n$$
,

and therefore

$$\sum\limits_{1}\sum\limits_{j=1}^{m}{y_{ij}b_{ij}}=cn$$
 ,

where  $\Sigma_1$  runs over those c values of i such that  $\sigma$  and  $\tau$  commute on  $C_i$ .

Next suppose that  $C_i$  is a cycle on which  $\sigma$  and  $\tau$  do not commute. Let  $C_z$  be a cycle such that one or more blocks of  $C_i$  are mapped into  $C_z$  by  $\tau$ . Let us denote these blocks by  $B_1, B_2, \dots, B_u$ . We may suppose that these blocks are numbered in such a way that  $B_1$  is the longest of them. Then  $b_{iz} \geq |B_1|$  by § 2. Moreover

$$n_{iz} = |B_1| + |B_2| + \cdots + |B_u|$$
 ,

and

$${y}_{iz} b_{iz} \geqq n_{iz} \, | \, B_1 \, | / n \geqq \sum\limits_{\mu=1}^u | \, B_\mu \, |^{\scriptscriptstyle 2} / n \; .$$

Hence

$$\sum\limits_{2}\sum\limits_{j=1}^{m}y_{ij}b_{ij} \geq \sum\limits_{s=1}^{n-1}s^{2}Q_{s}/n$$
 ,

where the summation  $\Sigma_2$  is taken over those values of *i* such that  $\sigma$  and  $\tau$  do not commute on  $C_i$ . Combining these results we obtain

$$\sum\limits_{i,j} {y}_{ij} b_{ij} \geqq cn + \sum\limits_{s=1}^{n-1} s^2 Q_s/n \geqq G( au)$$
 ,

which disposes of Case 1.

Case 2.

$$cn + \sum\limits_{s=1}^{n-1} s^2 Q_s/n < G( au)$$
 .

Since the total number of points of X that do not belong to any block is cn, we have

$$N=cn+\sum\limits_{s=1}^{n-1}sQ_s$$
 .

Therefore

(11) 
$$G(\tau) = cn + \sum_{s=1}^{n-1} sQ_s - \frac{(n-1)^2}{4n-6} \sum_{s=1}^{n-1} Q_s ,$$

and we have

(12) 
$$\sum_{s=1}^{n-1} s(n-s)Q_s > \frac{n(n-1)^2}{4n-6} \sum_{s=1}^{n-1} Q_s.$$

The inequality (12) cannot hold for n = 3. Hence  $n \ge 5$ ,  $K \ge 2$ .

Let q(i) denote the number of blocks in the cycle  $C_i$ . We denote the blocks of  $C_i$  by  $B_{1i}, B_{2i}, \dots, B_{q(i),i}$ , where we suppose the blocks are ordered in such a way that

$$|B_{1i}| \ge |B_{2i}| \ge \cdots \ge |B_{q(i),i}|$$
 .

We note that if  $\sigma$  and  $\tau$  do not commute on the cycle  $C_i$ , then  $q(i) \geq 2$ ,

$$\sum\limits_{w=1}^{q(i)} |\,B_{wi}\,| = n = 2K + 1$$
 ,

and  $|B_{\mu i}| \leq K$  for  $\mu \geq 2$ . If  $\sigma$  and  $\tau$  commute on the cycle  $C_i$ , then q(i) = 0.

We call  $C_i$  a special cycle if  $\sigma$  and  $\tau$  do not commute on  $C_i$  and  $|B_{1i}| \leq K$ . Let d denote the number of special cycles. Since every cycle that is composed of blocks and is not a special cycle contains exactly one block of length at least K + 1, we have

$$c+d+\sum\limits_{s=K+1}^{n-1}Q_s=m=N/n=c+\sum\limits_{s=1}^{n-1}sQ_s/n$$
 ,

or

(13) 
$$nd - \sum_{s=1}^{K} sQ_s + \sum_{s=K+1}^{n-1} (n-s)Q_s = 0$$
.

We call the block  $B_{wi}$  a special block if  $C_i$  is a special cycle and either

- (a) q(i) = 3, or
- (b) q(i) = 4 and  $w \leq 2$ .

The image  $B\tau$  of a block B is a block of  $\tau^{-1}$ . We call  $B\tau$  a block image. Let v(i) denote the number of block images in the cycle  $C_i$ , and let  $B'_{1i}, B'_{2i}, \dots, B'_{v(i),i}$  denote these block images. We can suppose that

$$|B_{1i}'| \geq |B_{2i}'| \geq \cdots \geq |B_{v(i),i}'|$$
 .

We call the block image  $B'_{wi}$  a special image if it is a special block of  $\tau^{-1}$ . More precisely  $B'_{wi}$  is a special image if  $|B'_{1i}| \leq K$  and either

(a) v(i) = 3, or (b) v(i) = 4 and  $w \le 2$ . If  $\sigma$  and  $\tau$  commute on the cycle  $C_i$  set

$$y_{ij} = egin{cases} 1 & ext{if } au ext{ maps } C_i ext{ onto } C_j ext{ ,} \ 0 & ext{otherwise .} \end{cases}$$

If  $C_i$  consists of blocks and is not a special cycle, then we set

$$y_{ij} = egin{cases} 1 & ext{if } au ext{ maps } B_{\scriptscriptstyle 1i} ext{ into } C_j ext{ ,} \ 0 & ext{otherwise }. \end{cases}$$

If  $C_i$  is a special cycle we set

$$y_{ij} = \Sigma''(K - |B|)/(K - 1)$$
 ,

where the summation  $\Sigma''$  runs over all special blocks B of  $C_i$  that  $\tau$  maps onto special images contained in  $C_j$ . Notice that replacing  $\tau$  by  $\tau^{-1}$  has the effect of replacing the matrix  $(y_{ij})$  by its transpose. Clearly  $y_{ij} \geq 0$  for all i, j. Moreover if the cycle  $C_i$  is not special, then

$$\sum\limits_{j=1}^m {y_{ij}} = 1$$
 .

Now suppose that  $C_i$  is a special cycle. Then

$$\sum\limits_{j=1}^m {y_{ij}} \leq {\Sigma'(K-\mid B\mid)}/{(K-1)}$$
 ,

where  $\Sigma'$  runs over all special blocks B of  $C_i$ . Since  $C_i$  is special we must have  $q(i) \ge 3$ . If q(i) = 3, then every block of  $C_i$  is special,  $\Sigma' |B| = 2K + 1$ , and

$$\Sigma'(K - |B|)/(K - 1) = (3K - \Sigma' |B|)/(K - 1) = 1$$
.

If q(i) = 4, then

$$|B_{1i}|+|B_{2i}|+|B_{3i}|+|B_{4i}|=2K+1$$
 ,

so that

$$\Sigma' \, | \, B \, | = | \, B_{\scriptscriptstyle 1i} \, | + | \, B_{\scriptscriptstyle 2i} \, | \ge K + 1$$
 ,

and

$$\Sigma'(K - |B|)/(K - 1) = (2K - \Sigma' |B|)/(K - 1) \leq 1$$
.

Finally if  $q(i) \ge 5$ , then  $C_i$  contains no special blocks, so that

$$\Sigma'(K - |B|)/(K - 1) = 0$$
.

Thus we have

$$\sum\limits_{j=1}^m y_{ij} \leqq 1$$
 ,  $1 \leqq i \leqq m$  .

By interchanging au and  $au^{-1}$  we obtain

$$\sum\limits_{i=1}^m y_{ij} \leqq 1, 1 \leqq j \leqq m$$
 .

Thus conditions (5), (8), and (9) are satisfied. We must show that (10) is satisfied also.

Let  $T_s$  denote the total number of special blocks of length s. Similarly let  $U_s$  denote the total number of special images of length s. Since there are exactly  $Q_s - U_s$  block images of length s that are not special images, it follows that there are at least

$$T_s-(Q_s-U_s)=T_s+U_s-Q_s$$

special blocks of length s that are mapped onto special images by  $\tau$ .

If  $\sigma$  and  $\tau$  commute on the cycle  $C_i$ , then

$$\sum\limits_{j=1}^m {{y_{ij}}{b_{ij}}} = n$$

If  $C_i$  consists of blocks and is not a special cycle, then  $|B_{1i}| \ge K+1$ , and

$$\sum\limits_{j=1}^m {y_{ij}} b_{ij} \geqq |B_{1i}|$$
 .

If  $C_i$  is a special cycle, then

$$\sum_{j=1}^m {y_{ij} b_{ij}} = \sum_{j=1}^m {{\varSigma''(K-|B|) b_{ij}}/(K-1)} \ \ge {{\varSigma^{\sharp}|B|(K-|B|)/(K-1)}} \;,$$

where  $\Sigma''$  runs over those special blocks B of  $C_i$  that are mapped onto special images contained in  $C_j$  by  $\tau$ , and  $\Sigma^*$  runs over all special blocks B of  $C_i$  that are mapped onto special images by  $\tau$ . It follows that

(14) 
$$\sum y_{ij}b_{ij} \ge cn + \sum_{s=K+1}^{n-1} sQ_s + \sum_{s=1}^{K} s(T_s + U_s - Q_s)(K-s)/(K-1) .$$

To complete the proof of the lemma it is sufficient to show that (10) holds. Suppose that (10) does not hold. Then

$$G( au)>\sum\limits_{i,j}y_{ij}b_{ij}$$
 .

Using (11) and (14) this gives us

$$cn+\sum\limits_{s=1}^{n-1} sQ_s - rac{(n-1)^2}{4n-6} \sum\limits_{s=1}^{n-1} Q_s \ > cn+\sum\limits_{s=K+1}^{n-1} sQ_s + \sum\limits_{s=1}^K s(T_s+U_s-Q_s)(K-s)/(K-1) \; ,$$

or

(15)  
$$\sum_{s=1}^{K} s\{Q_s - (T_s + U_s - Q_s)(K-s)/(K-1)\} > \frac{(n-1)^2}{4n-6} \sum_{s=1}^{n-1} Q_s.$$

We multiply (15) by n-3 and add (12). Since n-3=2(K-1) this gives as

(16)  

$$\sum_{s=1}^{K} s\{(2n-s-3)Q_s - 2(T_s + U_s - Q_s)(K-s)\} + \sum_{s=K+1}^{n-1} s(n-s)Q_s$$

$$> \frac{1}{2}(n-1)^2 \sum_{s=1}^{n-1} Q_s = 2K^2 \sum_{s=1}^{n-1} Q_s .$$

Now we multiply (13) by K-1 and add (16). This yields

(17) 
$$(K-1)nd - V_1 - V_2 + W_1 + W_2 > 0$$
,

where

$$egin{aligned} &V_1 = 2\sum\limits_{s=1}^K sT_s(K-s) \;, \ &V_2 = 2\sum\limits_{s=1}^K sU_s(K-s) \;, \ &W_1 = \sum\limits_{s=1}^K \{s(2n-s-K-2)+2s(K-s)-2K^2\}Q_s \ &= \sum\limits_{s=1}^K \{s(3K-s)+2s(K-s)-2K^2\}Q_s \ &= \sum\limits_{s=1}^K (K-s)(3s-2K)Q_s \;, \end{aligned}$$

and

$$egin{aligned} W_2 &= \sum\limits_{s=K+1}^{n-1} \left\{ (K-1)(n-s) + s(n-s) - 2K^2 
ight\} Q_s \ &= \sum\limits_{s=K+1}^{n-1} (s-1)(K-s+1)Q_s \ . \end{aligned}$$

The effect on (17) of replacing au by  $au^{-1}$  is to interchange  $V_1$  and  $V_2$ .

Now  $D(\sigma\tau, \tau\sigma) = D(\sigma\tau^{-1}, \tau^{-1}\sigma)$  and  $D_{\sigma}(\tau) = D_{\sigma}(\tau^{-1})$ . Thus it is sufficient to prove the desired result with  $\tau$  replaced by  $\tau^{-1}$ . It follows that we can assume, without loss of generality, that  $V_1 \leq V_2$ . Then we obtain

or

(18)  
$$(K-1)nd > \sum_{s=1}^{K} \{(K-s)(2K-3s)Q_s + 4s(K-s)T_s\} + \sum_{s=K+1}^{n-1} (s-1)(s-K-1)Q_s.$$

Let  $Q_s^{(i)}$  denote the number of blocks of length s in the cycle  $C_i$ , and let  $T_s^{(i)}$  denote the number of special blocks of length s in  $C_i$ . Then (18) can be written in the form

(19) 
$$(K-1)nd > \sum_{i=1}^{m} Z_i$$
,

where

$$egin{aligned} Z_i &= \sum\limits_{s=1}^K \left\{ (K-s)(2K-3s)Q_s^{(i)} + 4s(K-s)T_s^{(i)} 
ight\} \ &+ \sum\limits_{s=K+1}^{n-1} (s-1)(s-K-1)Q_s^{(i)} \;. \end{aligned}$$

If  $\sigma$  and  $\tau$  commute on the cycle  $C_i$  we have  $Q_s^{(i)} = T_s^{(i)} = 0$  for all s, so that  $Z_i = 0$ .

If the cycle  $C_i$  contains exactly two blocks,  $B_{1i}$  and  $B_{2i}$ , then we set  $s' = |B_{2i}|$ , and we have  $s' \leq K$ ,  $|B_{1i}| = 2K + 1 - s' \geq K + 1$ ,  $T_s^{(i)} = 0$  for all s, and

$$egin{aligned} Z_i &= (K-s')(2K-3s') + (2K-s')(K-s') \ &= 4(K-s')^2 \geqq 0 \;. \end{aligned}$$

Now suppose that  $C_i$  is a cycle that is not special, but that contains three or more blocks. Thus  $q(i) \ge 3$ , and  $|B_{1i}| > K$ . Set f(x) = (K-x)(2K-3x). The second derivative of the function f is positive, so that f is a convex function. Now  $|B_{2i}| + |B_{3i}| \le n - |B_{1i}| \le K$ . Therefore  $f(|B_{2i}|/2 + |B_{3i}|/2) > 0$ . Now for  $w \ge 4$ , we have  $|B_{wi}| \le K/3$ and  $f(|B_{wi}|) > 0$ . Whence

$$egin{aligned} &Z_i \geqq \sum\limits_{w=2}^{q(i)} f(\mid B_{wi} \mid) \geqq f(\mid B_{2i} \mid) \ &+ f(\mid B_{3i} \mid) \geqq 2f(\mid B_{2i} \mid /2 + \mid B_{3i} \mid /2) > 0 \;. \end{aligned}$$

We have shown that  $Z_i \geq 0$  for every *i* such that  $C_i$  is not a special cycle. Hence these terms can be dropped from the right side of (19). Now there are exactly *d* special cycles. Therefore, by (19), there is a special cycle  $C_i$  such that

$$Z_{\scriptscriptstyle t} < (K-1)n = 2K^{\scriptscriptstyle 2} - K - 1$$
 .

Since  $C_t$  is special we have  $Q_s^{(t)} = 0$  for s > K, and so

(20) 
$$2K^2 - K - 1 > Z_t = \sum_{s=1}^{K} \{ (K - s)(2K - 3s)Q_s^{(t)} + 4s(K - s)T_s^{(t)} \}.$$

Now set q = q(t); and  $s_w = |B_{wt}|, 1 \leq w \leq q$ . Then (20) can be written in the form

(21) 
$$2K^2 - K - 1 > \sum_{w=1}^{q} (K - s_w) H(w)$$

where

$$H(w) = egin{cases} 2K+s_w & ext{if } B_{wt} ext{ is a special block ,} \ 2K-3s_w & ext{if } B_{wt} ext{ is not a special block .} \end{cases}$$

Since  $C_t$  is a special cycle we have  $q = q(t) \ge 3$ .

(A) Suppose  $q \ge 5$ . Then  $C_t$  has no special blocks, and (21) becomes

$$2K^{\scriptscriptstyle 2} - K - 1 > \sum\limits_{w=1}^q f(s_w)$$
 ,

where f(x) = (K - x)(2K - 3x) as before. Since f is a convex function we have

$$\sum\limits_{w=1}^q f(s_w) \geqq qf(\varSigma s_w/q) = qf(n/q)$$
 .

Now f(x) is a decreasing function of x for  $x \leq 5K/6$ , and

$$n/q \leq n/5 = (2K+1)/5 < 5K/6$$
 .

Hence  $f(n/q) \ge f(n/5)$ . Moreover

$$25f(n/5) = (5K - n)(10K - 3n) = (3K - 1)(4K - 3)$$
,

which is positive. Therefore

$$5(2K^2 - K - 1) > 5qf(n/q) \ge 25f(n/5) = (3K - 1)(4K - 3)$$
,

 $\mathbf{or}$ 

$$0>2K^{\scriptscriptstyle 2}-8K+8=2(K-2)^{\scriptscriptstyle 2}$$
 ,

which is impossible. This disposes of the case  $q \ge 5$ . Hence q = 3

or q = 4.

(B) Next suppose that q = 3. Here all blocks of  $C_t$  are special blocks so that (21) gives us

(22)  
$$2K^{2} - K - 1 > \sum_{w=1}^{3} (K - s_{w})(2K + s_{w})$$
$$= 2K \sum_{w=1}^{3} (K - s_{w}) + \sum_{w=1}^{3} s_{w}(K - s_{w}) .$$

Now

$$\sum_{w=1}^{3} (K - s_w) = 3K - \sum_{w=1}^{3} s_w = 3K - n = K - 1$$

We have  $K \ge s_1 \ge s_2 \ge s_3 \ge 1$ ,  $s_1 + s_2 + s_3 = 2K + 1$ , and  $K \ge 2$ . Hence  $s_3 < K$ . Therefore  $1 \le s_3 \le K - 1$ , and we have

$$\sum\limits_{w=1}^3 s_w(K-s_w) \geqq s_3(K-s_3) \geqq K-1$$
 .

Substitution in (22) now gives us

$$2K^{\scriptscriptstyle 2}-K-1>2K(K-1)+K-1$$
 ,

a contradiction. Thus we have eliminated the case q = 3. There remains only q = 4.

(C) Suppose finally that q = 4. Here  $B_{1t}$  and  $B_{2t}$  are special blocks,  $B_{3t}$  and  $B_{4t}$  are not. Thus (21) gives us

(23) 
$$2K^2-K-1>L_1+L_2+M_3+M_4$$
 ,

where  $L_w = (K - s_w)(2K + s_w)$  and

$$M_w = f(s_w) = (K - s_w)(2K - 3s_w)$$
.

If n = 5, then K = 2,  $s_1 = 2$ ,  $s_2 = s_3 = s_4 = 1$ ,  $L_1 = 0$ ,  $L_2 = 5$ ,  $M_3 = M_4 = 1$ , which contradicts (23). Hence  $n \ge 7$  and  $K \ge 3$ .

Now set  $J = s_3 + s_4 = 2K + 1 - s_1 - s_2$ . Then since

$$s_1 \geqq s_2 \geqq s_3 \geqq s_4$$
 ,

we have  $J \leq K$ . Since f(x) is convex we have

$$M_{3}+M_{4}=f(s_{3})+f(s_{4})\geq 2f(J/2)=(2K-J)(4K-3J)/2$$
 .

combining this with (23) we get

$$2K^2>L_1+L_2+M_3+M_4 \ge L_1+L_2+4K^2-5KJ+3J^2/2$$
 ,

or

$$0>2L_{\scriptscriptstyle 1}+2L_{\scriptscriptstyle 2}+4K^{\scriptscriptstyle 2}-10KJ+3J^{\scriptscriptstyle 2}$$
 .

Since  $K \ge 3$ , we have  $2K + 1 \le 7K/3$ , and

$$J \leq 7K/3 - s_1 - s_2$$

Since  $s_1 + s_2 > K$ , we have  $7K/3 - s_1 - s_2 \leq 4K/3$ . Now  $3x^2 - 10Kx$  is a decreasing function of x for  $x \leq 5K/3$ . Hence

$$egin{aligned} 3J^2 - 10KJ &\geq 3(7K/3 - s_1 - s_2)^2 - 10K(7K/3 - s_1 - s_2) \ &= -7K^2 - 4K(s_1 + s_2) + 3(s_1 + s_2)^2 \;. \end{aligned}$$

Combining inequalities we get finally

$$egin{aligned} 0 &> 2L_1 + 2L_2 + 4K^2 + 3J^2 - 10KJ \ &\geqq 2(K-s_1)(2K+s_1) + 2(K-s_2)(2K+s_2) \ &- 3K^2 - 4K(s_1+s_2) + 3(s_1+s_2)^2 \ &= 5K^2 - 6K(s_1+s_2) + s_1^2 + 6s_1s_2 + s_2^2 \ &= 4(K-s_1)(K-s_2) + (s_1+s_2-K)^2 \ . \end{aligned}$$

This is impossible since  $K \ge s_1 \ge s_2$ . This contradiction completes the proof of the lemma.

Lemma 2 shows that  $D_{\sigma} \leq (n-1)^2/(4n-6)$  if n is odd, regardless of the size of m. Combining this with Lemma 1 we obtain our main result:

THEOREM. If  $\sigma$  is the product of m cycles of length n, where n is odd,  $n \geq 3$ , N = nm, and  $m \geq n - 2$ , then

(24) 
$$D_{\sigma} = (n-1)^2/(4n-6)$$
.

In the notation of [1], (24) becomes

$$d_{\sigma} = rac{(n-1)^2}{2n(2n-3)}$$

## References

1. Daniel Gorenstein, Reuben Sandler and W. H. Mills, On Almost-commuting permutations, Pacific J. Math. 12 (1962), 913-923.

2. John von Neumann, A Certain Zero-Sum Two-Person Game Equivalent to the Optimal Assignment Problem, Contributions to the Theory of Games, Vol. 2 (Edited by H. W. Kuhn and A. W. Tucker), pp. 5-12.

3. Samuel Karlin, Mathematical Methods and Theory in Games, Programming and Economics, Vol. 1, (1959).

YALE UNIVERSITY, INSTITUTE FOR DEFENSE ANALYSES