

TRANSFER THEOREMS FOR TOPOLOGICAL STRUCTURES

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Transfer theorems are obtained for the following mathematical situations.

\mathcal{A} is a dense substructure of the compact structure \bar{X} . $\{X_i\}$ is the set of all finitely generated substructures of X . F is a structure of functions from Y to the structure X .

The sentences transferred in the above situations are best described as “almost” positive, variables appearing in a negative subformula are quantified in a prescribed manner.

The main tools of this investigation are the manipulation of classical transfer theorems in the context of commutative diagrams, the ultraproduct construction, and the \mathcal{D} -limit operation of Chang and Keisler’s “Continuous Model Theory.”

Introduction. Every subgroup of an abelian group is abelian, any extension ring of a ring with zero divisors has zero divisors, and each homomorphic image of a commutative ring is commutative are instances of classical transfer theorems (Lemma 1). In this paper transfer theorems are obtained for the following mathematical situations:

X is a dense substructure of the compact topological structure \bar{X} ,

$\{X_i\}$ is the set of all finitely generated substructures of X , and F is a structure of functions from Y to the structure X .

The sentences transferred in the above situations are best described as “almost” positive, variables appearing in a negative subformula of the sentence are quantified in a prescribed manner.

The main tools of this investigation are the manipulation of the classical transfer theorems in the context of commutative diagrams, the ultraproduct construction, and the \mathcal{D} -limit operator popularized by Chang and Keisler [3].

Consider a first order language \mathcal{L} and its associated structures. A formula σ is identified with its prenex normal form, i.e., σ is assumed to be of the form

$$Q_1 v_1 Q_2 v_2, \dots, Q_n v_n M,$$

where each Q_i is a quantifier (\forall or \exists), each v_i is a variable, and M is a formula constructed from (positive) atomic formulae and their negations, *negative atomic formulae*, using the connectives “and”

and "or" (\wedge and \vee). Q_1v_1, \dots, Q_nv_n is the *prefix*, and M the *matrix*, of σ . A variable or constant term is *negative* in σ if it occurs in some negative atomic formula of σ . Formulae are classified on the basis of their prefix and matrix.

Assume a formula $\sigma = Q_1v_1Q_2v_2, \dots, Q_nv_nM$. σ is *positive (negative)* if its matrix M is constructed solely from positive (negative) atomic formulae. σ is *universal (existential)* if each quantifier of its prefix is a universal (existential) quantifier. σ is *universal-positive* if it is both universal and positive.

Vector notation will be used when convenient. $\sigma(v_1, \dots, v_n)$ is written $\sigma(\vec{v})$, $\sigma(c_1, \dots, c_n)$ as $\sigma(\vec{c})$, and $\exists v_1, \dots, \exists v_n\sigma(v_1, \dots, v_n)$ as $\exists \vec{v}\sigma(\vec{v})$.

Let \mathcal{A} be a structure. $\mathcal{A} \vdash \sigma(\vec{a})$ denotes that the n -tuple \vec{a} of elements of \mathcal{A} satisfies the formula $\sigma(\vec{v})$. Let φ be a function between the structures \mathcal{A} and \mathcal{B} . φ *preserves* the n -ary operation $F(\vec{v})$ if $F(\varphi(\vec{a})) = \varphi(F(\vec{a}))$ for all n -tuples \vec{a} of elements of \mathcal{A} . φ *preserves* the formula $\sigma(\vec{v})$ if $\mathcal{A} \vdash \sigma(\vec{a})$ implies $\mathcal{B} \vdash \sigma(\varphi(\vec{a}))$. \mathcal{B} is a *homomorphic image* of \mathcal{A} if there exists a function φ from \mathcal{A} onto \mathcal{B} which preserves all operations and positive atomic formulae. $\sigma(\vec{v})$ is *preserved under homomorphic image* if $\sigma(\vec{v})$ is preserved by all onto homomorphisms. \mathcal{A} is a *substructure* of \mathcal{B} (and \mathcal{B} is an *extension* of \mathcal{A}) if there exists a 1-1 function φ from \mathcal{A} to \mathcal{B} which preserves the operations, positive atomic formulae, and negative atomic formulae of \mathcal{A} . $\sigma(\vec{v})$ is *preserved under substructure* if $\mathcal{B} \vdash \sigma(\vec{b})$ and $\varphi(\vec{a}) = \vec{b}$ imply $\mathcal{A} \vdash \sigma(\vec{a})$ for every substructure morphism $\varphi: \mathcal{A} \rightarrow \mathcal{B}$. $\sigma(\vec{a})$ is *preserved under extension* if $\sigma(\vec{v})$ is preserved by every substructure morphism.

LEMMA 1 (Classical Theorems). *Let σ be a formula.*

- (i) (Los, Tarski) σ is preserved under substructures iff σ is universal.
- (ii) (Robinson) σ is preserved under extensions iff σ is existential.
- (iii) (Lyndon) σ is preserved under homomorphic images iff σ is positive.

1. Main results. \mathcal{A} is a *topological structure* if its underlying set A is endowed with a Hausdorff topology wherein each operation of \mathcal{A} is continuous and wherein each n -ary relation of \mathcal{A} is closed as a subset of A^n ; \mathcal{A} is a *compact topological structure* if its underlying topology is compact. \mathcal{A} is a *dense substructure* of the topological structure $\bar{\mathcal{A}}$ if \mathcal{A} is a substructure of $\bar{\mathcal{A}}$ whose underlying set is topologically dense in $\bar{\mathcal{A}}$. The collection of

compact structures $\bar{\mathcal{M}}$ compactifies the collection of structures \mathcal{M} if each structure of \mathcal{M} is a dense substructure of a member of $\bar{\mathcal{M}}$.

The formula σ is \exists -generalized positive if the prefix of σ begins with an initial string of existential quantifiers (i.e., $\sigma = \exists \bar{v} Q_1 v_1, \dots, Q_n v_n M$) such that every negative variable of σ is free or existentially quantified in the initial string $\exists \bar{v}$. The negation of a \exists -generalized positive formula is called a \forall -generalized negative formula; it begins with an initial string of universal quantifiers such that each of its nonnegative variables is free or universally quantified in the initial string.

THEOREM 1. *Let \mathcal{A} be a dense substructure of the compact structure $\bar{\mathcal{A}}$. Every \exists -generalized positive sentence satisfied by \mathcal{A} is satisfied by $\bar{\mathcal{A}}$.*

This generalizes a result of Robinson [7].

THEOREM 2. *Let $\bar{\mathcal{M}}$ compactify the collection of structures \mathcal{M} .*
 (i) *Every \exists -generalized positive sentence satisfied in \mathcal{M} is satisfied in $\bar{\mathcal{M}}$.*
 (ii) *Every \forall -generalized negative sentence valid in $\bar{\mathcal{M}}$ is valid in \mathcal{M} .*

COROLLARIES TO THEOREM 2. *Let σ be a \forall -generalized negative sentence (in an appropriate language).*

2.1. *σ is valid for all extensions of the group \mathcal{G} if σ is valid for all compact groups which are extensions of \mathcal{G} .*

2.2. *σ is valid for all total orderings if σ is valid for all complete total orderings.*

2.3. *σ is valid for all totally ordered lattices extending the ordering \mathcal{F} if σ is valid for all complete totally ordered lattices extending \mathcal{F} .*

2.4. *σ is valid for all distributive lattices if σ is valid for all complete distributive lattices.*

2.5. *σ is valid for all boolean algebras if σ is valid for all complete boolean algebras.*

The corollaries to Theorem 2 follow from:

Every group is a dense subgroup of a compact group (Kelly [6], p. 247).

Every total order can be (densely) embedded in a complete total

order, and every complete total order forms a compact topological space wherein the lattice operations, \wedge and \vee , are continuous (Birkhoff [2] pgs. 240-253).

The lattice (boolean algebra) of all subsets of a fixed set X endowed with the interval topology forms a compact topological space wherein the lattice (boolean algebra) operations are continuous. Every distributive lattice is isomorphic to a ring of sets, and every boolean algebra is isomorphic to a field of sets.

Corollary 2.3 considered the total order as a lattice so as to enlarge the class of sentences preserved. The analog of Corollary 2.3 for arbitrary lattices is false. It is difficult to go from complete lattice to compact-Hausdorff space. In fact,

Counterexample 1.1. (Birkhoff [2], pgs. 126-128.) The completion by cuts of the modular lattice of all subspaces of Hilbert space which have finite dimension or codimension is nonmodular (i.e., does not satisfy the modular equation).

Counterexample 1.2. If in Theorem 1 the hypothesis, the topology of \mathcal{A} is compact, is replaced by the topology of \mathcal{A} is a complete metric space, even positive sentences are not preserved.

Let $\sigma = \forall y \exists x (y^2 - 2)x = 1$. Rationals $\vdash \sigma$, but not Reals $\vdash \sigma$.

Counterexample 1.3. It does not seem that the class of sentences preserved under the hypothesis of Theorem 1 can be enlarged. A reasonable conjecture would be the class of all sentences whose negative variables are existentially quantified.

Consider the sentence in the language of order which states no first element exists,

$$\sigma = \forall x \exists y \exists z (y \leq x \text{ and } z \leq y \text{ and } z \neq y).$$

The open interval $(0, 1)$ is dense in $[0, 1]$; $(0, 1) \vdash \sigma$, but not $[0, 1] \vdash \sigma$.

The collection of substructures $\{\mathcal{A}_i\}$ generates the structure \mathcal{A} if each finite subset of \mathcal{A} is a subset of some \mathcal{A}_i . The formula σ is $\forall^+ \exists$ -generalized positive if $\sigma = Q_1 v_1, \dots, Q_n v_n \exists \bar{v} M$ (variables of $\exists \bar{v}$ are called *terminal*) such that each existentially quantified variable is terminal, or is nonnegative and does not appear in any atomic formula in which a terminal variable appears.

A formula σ is *universal-existential* if $\sigma = \forall \bar{w} \exists \bar{v} M$.

THEOREM 4. Let $\{\mathcal{A}_i\}$ generate the structure \mathcal{A} .

(i) (Los and Susco) Any universal-existential sentence valid in $\{\mathcal{A}_i\}$ is satisfied by \mathcal{A} .

(ii) If, in addition, \mathcal{A} is a compact structure, then every $\forall^+\exists$ -generalized positive sentence valid in $\{\mathcal{A}_i\}$ is satisfied by \mathcal{A} .

COROLLARIES TO THEOREM 4. Let σ be a $\forall^+\exists$ -generalized positive sentence.

4.1. If σ is valid for all finitely generated groups, then σ is valid for all compact groups.

4.2. If σ is valid for all finite total orderings (viewed as a lattice), then σ is valid for all complete totally ordered lattices.

As an application of Corollary 4.2, every totally ordered complete lattice must have a first element. The situation is radically different if we inquire about second elements. We counterexample attempts to weaken conditions on sentences preserved under the hypotheses of Theorem 4 (ii). Examples are taken from the language of total orderings with first element 0 and last element 1.

Counterexample 4.1. Consider the sentence which has a non-terminal existentially quantified negative variable,

$$\sigma = \exists x \forall y (x \neq 0, \text{ and } x \leq y \text{ or } y = 0) .$$

σ asserts the existence of a second element. σ is not satisfied by the complete total ordering $[0, 1]$, but every finite subset of $[0, 1]$ satisfies σ .

Counterexample 4.2. Consider the sentence which has a non-terminal (nonnegative) existentially quantified variable which appears in an atomic formula with a terminal variable,

$$\sigma = \exists x \forall y \exists z (z = x \text{ and } z \neq 0, \text{ and } y = 0 \text{ or } z \leq y) .$$

Again, σ asserts the existence of a second element. $[0, 1]$ doesn't satisfy σ , but each finite subset does satisfy σ .

Let \mathcal{A} be a structure, \mathcal{A}^I denotes the product structure whose underlying set is the set of all functions from I to \mathcal{A} such that operations on functions of \mathcal{A}^I are performed coordinatewise and, such that an atomic relation holds in \mathcal{A}^I iff it holds in each coordinate. F is a *functional structure of range \mathcal{A}* if F is a substructure of \mathcal{A}^I which contains all constant functions. The functional structure $F \subset \mathcal{A}^I$ has *finite solution property* if, for all unequal $f, f' \in F$, $\{i \in I \mid f(i) = f'(i)\}$ is finite.

The formula σ is \forall^+ -generalized positive if no negative variable of σ is existentially quantified.

THEOREM 5. Let F be a functional structure of range \mathcal{A} .

(i) Every \forall^+ -generalized positive sentence satisfied by F is satisfied by \mathcal{A} .

(ii) If, in addition, F has the finite solution property then every $\forall^+\exists$ -generalized positive sentence satisfied by F is satisfied by \mathcal{A} .

The negation of a \forall^+ -generalized positive sentence is a \exists^+ -generalized negative sentence; it has the defining property that each variable appearing in a positive atomic formula is existentially quantified.

COROLLARY 5.1. Every \exists^+ -generalized negative sentence is preserved under direct powers.

COROLLARY 5.2. (i) The ring of real polynomials in one variable is functional over the ring of real numbers, with finite solution property.

(ii) Theorem 5 (ii) may be applied to the ring of functions analytic in the circle, $|z| < 1$.

(iii) Let \mathcal{R} be a ring and M and M' \mathcal{R} -modules. The module of homomorphisms from M to M' is a functional structure of range M' .

If, in addition, M is irreducible (has no nontrivial submodules), the above module of homomorphisms has the finite solution property.

2. Proof of theorems. The function φ from \mathcal{A} to \mathcal{B} is an elementary embedding if, $\mathcal{A} \vdash \sigma(\bar{a})$ implies $\mathcal{B} \vdash \sigma(\varphi(\bar{a}))$ for all formulae $\sigma(\bar{v})$.

Throughout this section I denotes a fixed index set. \mathcal{D} is a filter on I if \mathcal{D} is a subset of the power set of I , such that

- if $X \in \mathcal{D}$ and $X \subseteq Y \subseteq I$, then $Y \in \mathcal{D}$, and
- if $X, Y \in \mathcal{D}$ then $X \cap Y \in \mathcal{D}$.

An ultrafilter on I is a maximal filter on I , i.e., a filter \mathcal{D} on I such that, for all $X \subseteq I$, $X \in \mathcal{D}$ or $I - X \in \mathcal{D}$. $\prod_{i \in I} X_i$ ($\prod_{i \in I} \mathcal{A}_i$) denotes the cartesian product of the sets $\{X_i\}_{i \in I}$ (of the structures $\{\mathcal{A}_i\}_{i \in I}$). When $\forall i \in I X_i = X$, $\prod_{i \in I} X_i$ is written X^I , and similarly for structures. X^I can be identified with the set of all functions from I to X ; an embedding $\text{diag}: X \rightarrow X^I$ is obtained by identifying $x \in X$ with the function with constant value x . For $f \in \prod_{i \in I} X_i$, $f(i)$ denotes the i th coordinate of f .

Assume a set of structures $\{\mathcal{A}_i\}_{i \in I}$ and an ultrafilter \mathcal{D} on I . The ultraproduct of $\{\mathcal{A}_i\}_{i \in I}$ by \mathcal{D} , written $\prod \mathcal{A}_i / \mathcal{D}$, is the homomorphic image of the product structure $\prod_{i \in I} \mathcal{A}_i$ induced by the congruence relation

$$f \equiv f' \text{ iff } \{i \mid f(i) = f'(i)\} \in \mathcal{D}, (f \text{ and } f' \in \prod \mathcal{A}_i).$$

In the case $\forall i \in I \mathcal{A}_i = \mathcal{A}$, the ultraproduct is called the ultrapower of \mathcal{A} by \mathcal{D} and is written $\mathcal{A}^I / \mathcal{D}$. Members of $\prod \mathcal{A}_i / \mathcal{D}$ are written $f / \mathcal{D}, f \in \prod \mathcal{A}_i$.

LEMMA 2 (Los). *The diagonal embedding $\text{diag}: \mathcal{A} \rightarrow \mathcal{A}^I / \mathcal{D}$ is an elementary embedding.*

DEFINITIONS. (i) \mathcal{A} is the *outside* of a \mathcal{B} -sandwich if there exist functions

$$\mathcal{A} \xrightarrow{\varphi} \mathcal{B} \xrightarrow{\psi} \mathcal{A}^I / \mathcal{D}$$

such that φ is a substructure morphism, ψ is a substructure morphism, and

$$\varphi \circ \psi = \text{diag}: \mathcal{A} \longrightarrow \mathcal{A}^I / \mathcal{D}.$$

(ii) \mathcal{A} is a *retract* of \mathcal{B} if there exist functions

$$\mathcal{A} \begin{array}{c} \xrightarrow{\varphi} \\ \psi \\ \longleftarrow \end{array} \mathcal{B}$$

such that φ is a substructure morphism, ψ is an onto homomorphism, and $\varphi \circ \psi = \text{identity of } \mathcal{A}$.

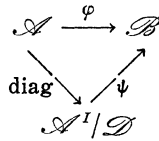
(iii) \mathcal{A} is a *strong retract* of \mathcal{B} if there exists functions

$$\mathcal{A} \begin{array}{c} \xrightarrow{\varphi} \\ \chi \\ \longleftarrow \end{array} \mathcal{B} \xrightarrow{\psi} \mathcal{A}^I / \mathcal{D}$$

such that \mathcal{A} is the outside of a \mathcal{B} -sandwich by $\mathcal{A} \xrightarrow{\varphi} \mathcal{B} \xrightarrow{\psi} \mathcal{A}^I / \mathcal{D}$ and \mathcal{A} is a retract of \mathcal{B} by

$$\mathcal{A} \begin{array}{c} \xrightarrow{\varphi} \\ \chi \\ \longleftarrow \end{array} \mathcal{B}$$

(iv) \mathcal{B} is an ultrahomomorphic image of \mathcal{A} if there exists a commutative diagram



such that φ is a substructure morphism and ψ an onto homomorphism.

THEOREM 6. *Let $\sigma(\bar{v})$ be a formula, \mathcal{A} a substructure of \mathcal{B} by $\mathcal{A} \xrightarrow{\varphi} \mathcal{B}$, and \bar{a} an n -tuple of elements of \mathcal{A} . \bar{a} shall also denote the n -tuple of elements of \mathcal{B} , $\varphi(\bar{a})$. $\mathcal{B} \vdash \sigma(\bar{a})$ implies $\mathcal{A} \vdash \sigma(\bar{a})$ if*

- (i) (Keisler) \mathcal{A} is the outside of a \mathcal{B} -sandwich and σ is a universal-existential formula.
- (ii) (Keisler) \mathcal{A} is a retract of \mathcal{B} and σ is a \forall^+ -generalized positive formula.
- (iii) \mathcal{A} is a strong retract of \mathcal{B} and σ is a $\forall^+\exists$ -generalized positive sentence.
- (iv) If \mathcal{B} is an ultrahomomorphic image of \mathcal{A} and σ is a \exists -generalized positive sentence, then $\mathcal{A} \vdash \sigma(\bar{a})$ implies $\mathcal{B} \vdash \sigma(\bar{a})$.

Proof. Results follow by induction on the number of quantifiers in $\sigma(\bar{v})$. We assume our language has constants for each element of \mathcal{A} , $\{a\}_{a \in \mathcal{A}}$, and constants for each element of \mathcal{B} , $\{b\}_{b \in \mathcal{B}}$. The interpretation of constants will be specified in each situation.

(i) Let $\mathcal{A} \xrightarrow{\varphi} \mathcal{B} \xrightarrow{\psi} \mathcal{A}^I/\mathcal{D}$ be a \mathcal{B} -sandwich.

The constant symbol a is interpreted as $\varphi(a)$ in \mathcal{B} and as $\alpha \circ \beta(a) = \text{diag}(a)$ in $\mathcal{A}^I/\mathcal{D}$ (it is here that we use $\alpha \circ \beta = \text{diag}$, to obtain an unambiguous interpretation of a in $\mathcal{A}^I/\mathcal{D}$). b is interpreted as $\psi(b)$ in $\mathcal{A}^I/\mathcal{D}$ and is not assigned an interpretation in \mathcal{A} .

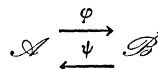
Assume \mathcal{B} satisfies a universal-existential formula, $\mathcal{B} \vdash \forall \bar{w} \exists \bar{v} M(\bar{w}, \bar{v}, \bar{a})$. We show that \mathcal{A} satisfies same formula. It suffices to show that, for each n -tuple \bar{a}' of elements of \mathcal{A} ,

$$\mathcal{A} \vdash \exists \bar{v} M(\bar{a}', \bar{v}, \bar{a}).$$

Clearly, $\mathcal{B} \vdash \exists \bar{v} M(\bar{a}', \bar{v}, \bar{a})$. Thus, for some \bar{b} , $\mathcal{B} \vdash M(\bar{a}', \bar{b}, \bar{a})$.

Since \mathcal{B} is a substructure of $\mathcal{A}^I/\mathcal{D}$, $\mathcal{A}^I/\mathcal{D} \vdash M(\bar{a}', \bar{b}, \bar{a})$. Thus $\mathcal{A}^I/\mathcal{D} \vdash \exists \bar{v} M(\bar{a}', \bar{v}, \bar{a})$. Hence, $\mathcal{A} \vdash \exists \bar{v} M(\bar{a}', \bar{v}, \bar{a})$.

(ii) Let



be a retraction.

a is interpreted as $\varphi(a)$ in \mathcal{B} (note that, since $\varphi \circ \psi = \text{identity}$

of \mathcal{A} , the interpretation of a in \mathcal{A} is unambiguous). b is interpreted as $\psi(b)$ in \mathcal{A} .

Consider a \forall^+ -generalized positive sentence $\sigma(\bar{a}, \bar{b})$ in which no constant of \bar{b} is negative. We show that

$$(1) \quad \mathcal{B} \vdash \sigma(\bar{a}, \bar{b}) \text{ implies } \mathcal{A} \vdash \sigma(\bar{a}, \bar{b}).$$

(1) is trivial if $\sigma(\bar{a}, \bar{b})$ has no quantifiers.

Now consider $\sigma(\bar{a}, \bar{b}, v)$ where no constant of \bar{b} is negative.

Assume that $\forall b \in \mathcal{B}$, (1) holds for $\sigma(\bar{a}, \bar{b}, b)$.

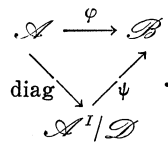
If $\mathcal{B} \vdash \forall v \sigma(\bar{a}, \bar{b}, v)$, then $\forall a \in \mathcal{A}$, $\mathcal{B} \vdash \sigma(\bar{a}, \bar{b}, a)$. By induction hypothesis, $\forall a \in \mathcal{A}$, $\mathcal{A} \vdash \sigma(\bar{a}, \bar{b}, a)$. Thus $\mathcal{A} \vdash \forall v \sigma(\bar{a}, \bar{b}, v)$. This completes the first part of our induction.

Now assume that v is not a negative variable of $\sigma(\bar{a}, \bar{b}, v)$ and assume $\mathcal{B} \vdash \exists v \sigma(\bar{a}, \bar{b}, v)$. For some $b \in \mathcal{B}$, $\mathcal{B} \vdash \sigma(\bar{a}, \bar{b}, b)$. Since v was not a negative variable, b does not appear negatively in $\sigma(\bar{a}, \bar{b}, b)$ and hence $\sigma(\bar{a}, \bar{b}, b)$ falls under the induction hypothesis. Thus $\mathcal{A} \vdash \sigma(\bar{a}, b, b)$ which implies $\mathcal{A} \vdash \exists v \sigma(\bar{a}, \bar{b}, v)$. This completes the induction.

We shall prove (iv) and then prove the difficult (iii).

We assume a language with special symbols for true and false, t and f ; we will also assume these symbols may be subscripted by a formula σ , t_σ and f_σ . p_σ shall vary over $\{t_\sigma, f_\sigma\}$. $\sigma[\varphi/\psi]$ denotes the formula resulting from replacing each occurrence of the formula φ in σ by ψ ; and similarly for finite sets of formulae Φ , $\sigma[\Phi/\psi]_{\sigma \in \Phi}$.

(iv) Assume the ultrahomomorphism



a is interpreted as $\varphi(a) = (\text{diag} \circ \psi)(a)$ in \mathcal{B} . b is not interpreted.

Let $\exists \bar{v} \sigma(\bar{a}, \bar{v})$ be a \exists -generalized sentence, i.e., $\sigma(\bar{a}, \bar{v})$ has no negative constants or variables other than those of \bar{a} and \bar{v} . Note that $\sigma(\bar{a}, \bar{v})$ may contain other quantified nonnegative variables. Assume $\mathcal{A} \vdash \exists \bar{v} \sigma(\bar{a}, \bar{v})$. We show $\mathcal{B} \vdash \exists \bar{v} \sigma(\bar{a}, \bar{v})$. $\mathcal{A} \vdash \exists \bar{v} \sigma(\bar{a}, \bar{v})$ implies $\mathcal{A} \vdash \sigma(\bar{a}, \bar{a}')$ for some \bar{a}' . Let N_σ denote the set of negative atomic formulae appearing in $\sigma(\bar{a}, \bar{a}')$. By the form of $\sigma(\bar{a}, \bar{v})$, each formula α of N_σ has no variables and contains only the constants of \bar{a} and \bar{a}' .

$$(2) \quad \forall \alpha \in N_\sigma, \mathcal{A} \vdash \alpha \text{ iff } \mathcal{B} \vdash \alpha.$$

Write $p_\alpha = t_\alpha$ if $\mathcal{A} \vdash \alpha$; $p_\alpha = f_\alpha$ otherwise ($\alpha \in N_\sigma$), and define $\sigma^*(\bar{a}, \bar{a}') = \sigma(\bar{a}, \bar{a}')[\alpha/p_\alpha]_{\alpha \in N_\sigma}$. Since \mathcal{A} is a substructure of \mathcal{B} and $\mathcal{A} \vdash \alpha \leftrightarrow p_\alpha$,

$$(3) \quad \mathcal{B} \vdash \alpha \longleftrightarrow p_\alpha .$$

Now, $\mathcal{A} \vdash \sigma^*(\bar{a}, \bar{a}')$ whence $\mathcal{A}^I/\mathcal{D} \vdash \sigma^*(\bar{a}, \bar{a}')$. Also $\sigma^*(\bar{a}, \bar{a}')$ is a positive sentence. $\mathcal{B} \vdash \sigma^*(\bar{a}, \bar{a}')$ by Lemma 1 (iii) since \mathcal{B} is a homomorphic image of $\mathcal{A}^I/\mathcal{D}$.

$\mathcal{B} \vdash \sigma^*(\bar{a}, \bar{a}') [\alpha/p_\alpha]_{\alpha \in N_\sigma}$ by (3). Thus, $\mathcal{B} \vdash \sigma(\bar{a}, \bar{a}')$ and $\mathcal{B} \vdash \exists \bar{v} \sigma(\bar{a}, \bar{v})$.

(iii) Let $\mathcal{A} \xrightarrow{\varphi} \mathcal{B} \xrightarrow{\psi} \mathcal{A}^I/\mathcal{D}$, $\mathcal{B} \xrightarrow{\chi} \mathcal{A}$ be a strong retract.

a is interpreted as $\varphi(a)$ in \mathcal{B} and as $\varphi \circ \psi = \text{diag}(a)$ in $\mathcal{A}^I/\mathcal{D}$.

b is interpreted as $\psi(b)$ in $\mathcal{A}^I/\mathcal{D}$ and as $\chi(b)$ in \mathcal{A} . To emphasize interpretation in \mathcal{A} (and to avoid confusion), $\chi(b)$ is written \hat{b} .

Let $\sigma(\bar{a}, \bar{b}) = Q_1 v_1, \dots, Q_n v_n \exists \bar{v} M(\bar{a}, \bar{b}, \bar{v})$ be such that

$$(4) \quad \text{No existentially quantified variable of}$$

$\langle v_1, \dots, v_n \rangle$ and no constant of \bar{b} is negative or appears in any atomic formula with a variable of \bar{v} .

$$(5) \quad \mathcal{B} \vdash \sigma(\bar{a}, \bar{b}) \text{ implies } \mathcal{A} \vdash \sigma(\bar{a}, \hat{b}), \text{ notice } \hat{b} ,$$

is shown by induction on the length of $Q_1 v_1, \dots, Q_n v_n$.

If $n = 0$, $\sigma(\bar{a}, \bar{b})$ is of the form $\exists \bar{v} M(\bar{a}, \bar{b}, \bar{v})$. Assume $\mathcal{B} \vdash \exists \bar{v} M(\bar{a}, \bar{b}, \bar{v})$. Set $A_M =$ the set of positive atomic formulae of $\exists \bar{v} M(\bar{a}, \bar{b}, \bar{v})$ in which variables \bar{v} do not appear. $\forall \alpha \in A_M$ write $p_\alpha = t_\alpha$ if $\mathcal{B} \vdash \alpha(\bar{a}, \bar{b})$; $p_\alpha = f_\alpha$ otherwise, and define

$$\exists \bar{v} M^*(\bar{a}, \bar{v}) = \exists \bar{v} M(\bar{a}, \bar{b}, \bar{v})[\alpha/p_\alpha]_{\alpha \in A_M} .$$

Now, $\mathcal{B} \vdash \exists \bar{v} M^*(\bar{a}, \bar{v})$ and no constant of \bar{b} occurs in $\exists \bar{v} M^*(\bar{a}, \bar{v})$ by (4). $\mathcal{A}^I/\mathcal{D} \vdash \exists \bar{v} M^*(\bar{a}, \bar{v})$ since $\mathcal{A}^I/\mathcal{D}$ is an extension of \mathcal{B} . Thus, $\mathcal{A} \vdash \exists \bar{v} M^*(\bar{a}, \bar{v})$ as no variable of \bar{b} occurs in $M^*(\bar{a}, \bar{v})$. Consider $\exists \bar{v} M^*(\bar{a}, \bar{v})[\alpha/p_\alpha]_{\alpha \in A_M}$. If $\mathcal{B} \vdash \alpha(\bar{a}, \bar{b})$, then $\mathcal{A} \vdash \alpha(\bar{a}, \bar{b})$ since $\alpha(\bar{a}, \bar{b})$ is positive and \mathcal{A} is a homomorphic image of \mathcal{B} ($\alpha(\bar{a}, \bar{b}) \in A_M$). Thus, if $\mathcal{B} \vdash \alpha(\bar{a}, \bar{b})$ (i.e., $p_\alpha = t_\alpha$), then $\mathcal{A} \vdash \exists \bar{v} M^*(\bar{a}, \bar{v})[\alpha/p_\alpha(\bar{a}, \hat{b})]$. Now, if not $\mathcal{B} \vdash \alpha(\bar{a}, \bar{b})$, then $p_\alpha = f_\alpha$ and it does no harm (by the form of M^*) to replace f_α by $\alpha(\bar{a}, \hat{b})$. Thus $\mathcal{A} \vdash \exists \bar{v} M^*(\bar{a}, \bar{v})[\alpha/p_\alpha]_{\alpha \in A_M}$. Thus, $\mathcal{A} \vdash \exists \bar{v} M(\bar{a}, \hat{b}, \bar{v})$.

We complete the induction.

Let $\forall v \sigma(\bar{a}, \bar{b}, v)$ obey (4) and let $\mathcal{B} \vdash \forall v \sigma(\bar{a}, \bar{b}, v)$. $\forall a \in \mathcal{A}$, $\mathcal{B} \vdash \sigma(\bar{a}, \bar{b}, a)$ and $\sigma(\bar{a}, \bar{b}, a)$ obeys (4). Thus by induction

hypothesis, $\forall a \in \mathcal{A}, \mathcal{A} \vdash \sigma(\bar{a}, \bar{b}, a)$. Hence $\mathcal{A} \vdash \forall v \sigma(\bar{a}, \bar{b}, v)$.

Let $\exists v \sigma(\bar{a}, \bar{b}, v)$ obey (4) and let $\mathcal{B} \vdash \exists v \sigma(\bar{a}, \bar{b}, v)$. For some $b \in \mathcal{B}, \mathcal{B} \vdash \sigma(\bar{a}, \bar{b}, b)$ where $\sigma(\bar{a}, \bar{b}, b)$ obeys (4) since $\exists v \sigma(\bar{a}, \bar{b}, v)$ obeys (4). By induction hypothesis, $\mathcal{A} \vdash \sigma(\bar{a}, \bar{b}, \hat{b})$. Thus $\mathcal{A} \vdash \exists v \sigma(\bar{a}, \bar{b}, v)$.

We prove Theorems 1-5 by showing that each situation yields a diagram to which Theorem 6 applies. Let I be a set. $S_o(I)$ denotes the set of all finite subsets of I . \mathcal{D} is a regular ultrafilter on $S_o(I)$ if \mathcal{D} is an ultrafilter on $S_o(I)$ such that

$$\forall i_0 \in I, \{S \in S_o(I) \mid i_0 \in S\} \in \mathcal{D}.$$

It follows that $\{S \in S_o(I) \mid T \subset S\} \in \mathcal{D}$ for each finite $T \subset I$.

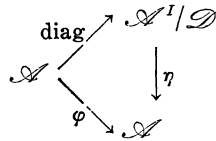
Assume a topological space X and ultrafilter \mathcal{D} on I . For each $x \in X$ and $f \in X^I, x = \mathcal{D}$ -limit f if $\{i \mid f(i) \in V_x\} \in \mathcal{D}$ for each open set V_x containing x .

Essentially the \mathcal{D} -lim operator on sequences $f \in X^I$ acts like the limit operation on countable sequences, with the added property that, if X is a compact space, \mathcal{D} -lim f always exists. The properties of the \mathcal{D} -limit operator are listed below. A more complete treatment is provided in Chang and Keisler [3], pgs. 6-15.

Let X be a topological space with basis δ (i.e., for each open set $V, V = \bigcup_{0 \in \delta} 0$) and let \mathcal{D} be an ultrafilter on I . f will vary over elements of X^I .

- \mathcal{D} -L0 \mathcal{D} -lim f is unique.
- \mathcal{D} -L1 If $\{i \mid f(i) = f'(i)\} \in \mathcal{D}$, then \mathcal{D} -lim $f = \mathcal{D}$ -lim f' (or both don't exist).
- \mathcal{D} -L2 \mathcal{D} -lim $\langle f_1, \dots, f_n \rangle = \langle \mathcal{D}$ -lim f_1, \dots, \mathcal{D} -lim $f_n \rangle$.
- \mathcal{D} -L3 \mathcal{D} -lim f belongs to the closure (in X) of the range of f , (range $f = \{x \in X \mid \exists i \in I, f(i) = x\}$).
- \mathcal{D} -L3' The \mathcal{D} -limit of a constant sequence equals that constant value.
- \mathcal{D} -L4 If F is a continuous function from X to Y , then \mathcal{D} -lim $F(f) = F(\mathcal{D}$ -lim f).
- \mathcal{D} -L4' If R is a closed n -ary relation and $\forall i \in I, R(f_1(i), \dots, f_n(i))$, then $R(\mathcal{D}$ -lim f_1, \dots, \mathcal{D} -lim f_n).
- \mathcal{D} -L5 If X is a compact space, then $\forall f \in X^I$ \mathcal{D} -lim f exists.
- \mathcal{D} -L6 If \mathcal{D} is a regular ultrafilter on $I = S_o(\delta), \delta$ a basis for \bar{X} , then, for each set $S \subseteq \bar{X}$ and limit point x_0 of S , there exists a sequence $f \in S^I$ such that \mathcal{D} -lim $f = x_0$.
- \mathcal{D} -L6' If X is dense in \bar{X} , for each $x_0 \in \bar{X}$, there exists an X -sequence $f \in X^I$ such that \mathcal{D} -lim $f = x_0$ (i.e., the map $X^I \xrightarrow{\mathcal{D}\text{-lim}} \bar{X}$ is onto).

THEOREM 7. Let \mathcal{A} be a dense substructure of the compact structure $\bar{\mathcal{A}}$, $\mathcal{A} \xrightarrow{\varphi} \bar{\mathcal{A}}$. For some set I and ultrafilter \mathcal{D} on I , $\mathcal{A}^I/\mathcal{D}$ is a homomorphic image of $\mathcal{A}^I/\mathcal{D}$ (via the map η) such that



commutes.

Proof. Let $\bar{\mathcal{A}}$ have basis δ . Let \mathcal{D} be any regular ultrafilter on $I = S_\omega(\delta)$.

Define the function $\eta: \mathcal{A}^I/\mathcal{D} \rightarrow \bar{\mathcal{A}}$ by $\eta(f/\mathcal{D}) = \mathcal{D}\text{-lim } f$. Since $\bar{\mathcal{A}}$ is compact, by \mathcal{D} -L5, η is defined $\forall f/\mathcal{D} \in \mathcal{A}^I/\mathcal{D}$. η is well defined by \mathcal{D} -L1, η is a homomorphism of structures by \mathcal{D} -L4 and \mathcal{D} -L4', and η is onto by \mathcal{D} -L6'. Finally, by \mathcal{D} -L3', the diagram of Theorem 7 commutes.

THEOREM 8.

(i) If \mathcal{A} is a dense substructure of the compact topological structure $\bar{\mathcal{A}}$, then $\bar{\mathcal{A}}$ is an ultrahomomorphism of \mathcal{A} .

(ii) If $\{\mathcal{A}_j\}$ generate \mathcal{A} , then \mathcal{A} is the outside of a $\prod_{i \in I} \mathcal{A}_i/\mathcal{D}$ -sandwich for some ultraproduct $\prod_{i \in I} \mathcal{A}_i/\mathcal{D}$.

If, in addition, \mathcal{A} is a compact structure, then \mathcal{A} is a strong retract of $\prod_{i \in I} \mathcal{A}_i/\mathcal{D}$.

(iii) If $F \subseteq \mathcal{A}^I$ is a functional structure of range \mathcal{A} , then \mathcal{A} is a retract of F .

If, in addition, F has the finite solution property, then \mathcal{A} is a strong retract of F .

Proof. (i) follows from Theorem 7.

(ii) Let $\{\mathcal{A}_j\}_{j \in J}$ generate \mathcal{A} . Let $I = S_\omega(\mathcal{A})$ and \mathcal{D} be a regular ultrafilter on I . Assign to each $i \in S_\omega(\mathcal{A})$ some structure \mathcal{A}_i from $\{\mathcal{A}_j\}_{j \in J}$ such that $i \subset \mathcal{A}_i$. Consider $\prod_{i \in I} \mathcal{A}_i/\mathcal{D}$.

\mathcal{A} is a substructure of $\prod_{i \in I} \mathcal{A}_i/\mathcal{D}$ by the map φ , where $\varphi(a) = f/\mathcal{D}$ such that $f(i) = a$ if $a \in i$ and $f(i)$ is arbitrary otherwise. Also, since $\mathcal{A}_i \subset \mathcal{A}$, $\prod_{i \in I} \mathcal{A}_i/\mathcal{D} \subset \mathcal{A}^I/\mathcal{D}$.

Finally, if \mathcal{A} is a compact structure, then \mathcal{D} -lim is a map from $\mathcal{A}^I/\mathcal{D}$ to \mathcal{A} . Since $\prod_{i \in I} \mathcal{A}_i/\mathcal{D}$ contains all constant functions, the restriction of \mathcal{D} -lim to $\prod_{i \in I} \mathcal{A}_i/\mathcal{D}$ is onto.

(iii) Since F contains all constant functions the map $\mathcal{A} \xrightarrow{\varphi} F$ is obvious. Choose any $a \in \mathcal{A}$; the projection map $F \xrightarrow{\psi} \mathcal{A}$, defined

by $\psi(f) = f(a)$, is onto since F contains all constant functions.

If, in addition, $F \subset \mathcal{A}_i$ has the finite solution property, then $\{\{i \mid f(i) \neq f'(i)\} \mid f \neq f' \in F\}$ has the finite intersection property and can be expanded to an ultrafilter \mathcal{D} on I . Clearly, the canonical map $F \rightarrow \mathcal{A}^I/\mathcal{D}$ is 1-1 ($F \subset \mathcal{A}^I$).

Theorems 1-5 readily follow from Theorems 6 and 8.

FINAL REMARKS. Theorems 1 and 2 apply to dense substructures of compact topological structures. We consider the non-compact case. A sentence σ is (\exists -generalized \forall)-positive if $\sigma = \exists \bar{v} \forall \bar{w} M(\bar{v}, \bar{w})$ where M is quantifier-free and every negative variable of M is existentially quantified in $\exists \bar{v}$.

THEOREM 3. *Let \mathcal{A} be a dense substructure of the (Hausdorff) topological structure $\hat{\mathcal{A}}$. Every (\exists -generalized \forall)-positive sentence satisfied by \mathcal{A} is satisfied by $\hat{\mathcal{A}}$.*

The proof of Theorem 3 is similar to the proof of Theorem 1. Instead of a homomorphism from an ultrapower of \mathcal{A} onto $\hat{\mathcal{A}}$, one obtains a homomorphism from a subset of an ultrapower of \mathcal{A} onto $\hat{\mathcal{A}}$ (the domain of the homomorphism is the set of all sequences whose \mathcal{D} -limits do exist). Theorem 3 then follows from a result of Keisler [5].

LEMMA 3 (Keisler). *If \mathcal{A} is a substructure of \mathcal{B} such that \mathcal{B} is a homomorphic image of a substructure of an ultrapower of \mathcal{A} such that $\forall a \in \mathcal{A}$ the image of a under the homomorphism is $a \in \mathcal{B}$, then every (\exists -generalized \forall)-positive sentence satisfied by \mathcal{A} is satisfied by \mathcal{B} .*

It is also shown in Keisler [5] that the sentences preserved under inverse limits are precisely the negations of (\exists -generalized \forall)-positive sentences.

Theorem 6 (iii) and (iv), as well as (i) and (ii), are characterizations, i.e., the classes of sentences described are the set of all sentences preserved under the indicated algebraic operation. Other preservation theorems and concepts may be found in Bell and Slomson [1] and Chang and Keisler [4].

We have just recently been able to extend Theorems 1-4 applied to groups to a weak second-order language where quantification over subgroups and normal subgroups are permitted. One must restrict the notion of subgroups to topologically closed subgroups. These results will appear elsewhere.

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Received December 20, 1974.

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