## A MAPPING PROPERTY OF REGRESSIVE ISOLS

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## 1. Introduction

Let $\varepsilon^{*}, \varepsilon, \Lambda, \Lambda_{R}$ and $\Lambda^{*}$ denote the collections of all integers, non-negative integers, isols, regressive isols and isolic integers respectively. Let $f\left(x_{1}, \cdots, x_{n}\right)$ be a recursive function, and let $f_{\Lambda}$ denote the canonical extension of $f$ to a mapping from $\Lambda^{n}$ into $\Lambda^{*}$ [11], [12]. A. Nerode proved in [12] that $f_{\Lambda}$ maps $\Lambda^{n}$ into $\Lambda$ if and only if $f$ is almost recursive combinatorial. In [3], J. Barback proved that if $f(x)$ is a recursive function of one variable, $f_{\Lambda}$ maps $\Lambda_{R}$ into $\Lambda_{R}$ if and only if $f$ is eventually increasing. However the author showed in [10] that the class of recursive functions of two variables mapping $\Lambda_{R} \times \Lambda_{R}$ into $\Lambda_{R}$ is rather limited-trivial cases aside, this class consists of functions eventually of the form

$$
\min (f(x), g(y))+\sum_{i=0}^{x} \sum_{j=0}^{y} d(i, j)
$$

where $\min (x, y)$ is the minimum function, $f(x)$ and $g(y)$ are eventually increasing recursive, and $d(i, j)=0$ for all but finitely many pairs $(i, j)$. The restrictive nature of this last result is not surprising in view of the fact that $\Lambda_{R}$ is not closed under addition or multiplication. J. Barback suggested that it is more natural to ask which recursive functions $f\left(x_{1}, \cdots, x_{n}\right)$ of more than one variable have the property

$$
\begin{equation*}
A_{1}+A_{2}+\cdots+A_{n} \in \Lambda_{R} \Rightarrow f_{\Lambda}\left(A_{1}, A_{2}, \cdots, A_{n}\right) \in \Lambda_{R} \tag{*}
\end{equation*}
$$

The main theorem of this paper characterizes the class of recursive functions with the property (*) as follows: Let $\leq$ be the partial ordering of $\varepsilon^{n}$ obtained by setting

$$
\left(x_{1}, \cdots, x_{n}\right) \leq\left(y_{1}, \cdots, y_{n}\right) \quad \text { iff } \quad x_{i} \leq y_{i}, \quad i=1, \cdots, n
$$

A function $f\left(x_{1}, \cdots, x_{n}\right)$ is called increasing if

$$
\left(x_{1}, \cdots, x_{n}\right) \leq\left(y_{1}, \cdots, y_{n}\right) \Rightarrow f\left(x_{1}, \cdots, x_{n}\right) \leq f\left(y_{1}, \cdots, y_{n}\right)
$$

and eventually increasing if there exists a number $k \in \varepsilon$ such that

$$
f\left(x_{1}+k, \cdots, x_{n}+k\right)
$$

is increasing. A function $g$ of fewer than $n$ variables is called a proper specification of $f\left(x_{1}, \cdots, x_{n}\right)$ if $g$ can be obtained from $f$ by substitution of constants for some of the variables $x_{1}, \cdots, x_{n}$ of $f . \quad f$ is called almost increasing if $f$ and every proper specification of $f$ is eventually increasing.

Theorem. Let $f\left(x_{1}, \cdots, x_{n}\right)$ be recursive. f has the property (*) if and only if $f$ is almost increasing.

We note that for $n=1$ the preceding theorem reduces to the theorem in [3] for functions of one variable. For the sake of simplicity, we shall give the proof only for $n=2$; the proof for $n>2$ requires no essential modification of the techniques to be presented here.

The mapping theorem for $\Lambda$ in [12] can be used to prove the truth in $\Lambda$ of certain $\forall \exists$ sentences with recursive combinatorial Skolem functions. In the final section of this paper we shall apply our main theorem in a similar manner to determine the truth or falsehood in $\Lambda_{R}$ of various sentences with recursive Skolem functions.

## 2. Preliminaries

We assume that the reader is familiar with the basic concepts and results of [1], [3], [6], [7], [10] and [14]. The following definitions and theorems concerning functions of two variables are the natural analogues of the concepts and results in Section 2 of [14] and Section 4 of [1]. The theorems can be proved using the methods of those papers, and will be stated without proof.

By a number-theoretic function of $n$ variables we shall mean a function mapping $\varepsilon^{n}$ into $\varepsilon^{*}$. Every number-theoretic function $f$ can be written as the difference of two combinatorial functions $f^{+}$and $f^{-}$, called the positive and negative parts of $f$. A number-theoretic function $f$ is called recursive if the functions $f^{+}$and $f^{-}$are recursive. For a recursive number-theoretic function $f\left(x_{1}, \cdots, x_{n}\right)$, we can employ the usual canonical extension procedure to define $f_{\Lambda}$, i.e., for any $n$-tuple of isols ( $x_{1}, \cdots, x_{n}$ ),

$$
f_{\Lambda}\left(x_{1}, \cdots, x_{n}\right)=f_{\Lambda}^{+}\left(x_{1}, \cdots, x_{n}\right)-f_{\Lambda}^{-}\left(x_{1}, \cdots, x_{n}\right)
$$

Let $f(x, y)$ be recursive and number-theoretic. For $T, U \in \Lambda_{R}$ we define

$$
\sum_{(r, v)}^{*} f(x, y)=\sum_{(r, v)} f^{+}(x, y)-\sum_{(r, v)} f^{-}(x, y)
$$

By the partial sum function of $f(x, y)$ we mean the function

$$
S_{f}(x, y)=\sum_{i<x} \sum_{j<y} f(i, j)
$$

$\left(S_{f}(x, y)=0\right.$ if $x=0$ or $\left.y=0\right)$.
Proposition 1. Let $f(x, y)$ and $g(x, y)$ be recursive number-theoretic functions and $T, U \in \Lambda_{R}$. Then

$$
\sum_{(T, V)}^{*} f(x, y) \pm \sum_{(T, V)}^{*} g(x, y)=\sum_{(T, v)}^{*}(f(x, y) \pm g(x, y))
$$

Theorem 1. Let $f(x, y)$ be recursive and number-theoretic. Then for all T, $U \in \Lambda_{R}, \sum_{(T, U)}^{*} f(x, y)=\left(S_{f}\right)_{\Lambda}(T, U)$.

For any recursive function $f(x, y)$, we define

$$
\begin{array}{rlrl}
\hat{f}(x, y) & =0, \quad \quad \text { if } x=0 \text { or } y=0 \\
& =f(x-1, y-1), \quad \text { otherwise. } \\
\Delta_{x} f(x, y) & =f(x+1, y)-f(x, y), \\
\Delta_{y} f(x, y) & =f(x, y+1)-f(x, y), \\
D f(x, y) & =\Delta_{x} \Delta_{y} \hat{f}(x, y), & \\
D f^{+}(x, y) & =D f(x, y), \quad D f(x, y) \geq 0 \\
& =0, \quad & \quad \text { otherwise } \\
D f^{-}(x, y) & =-D f(x, y), \quad D f(x, y) \leq 0 \\
& =0, \quad & & \text { otherwise. }
\end{array}
$$

Theorem 2. Let $f(x, y)$ be recursive. Then for $T, U \in \Lambda_{R}$,

$$
f_{\Lambda}(T, U)=\sum_{(T+1, U+1)} D f^{+}-\sum_{(T+1, U+1)} D f^{-}=\sum_{(T+1, U+1)}^{*} D f
$$

In particular, for $n, k \in \varepsilon$,

$$
f(n, k)=\sum_{i=0}^{n} \sum_{j=0}^{k} D f^{+}(i, j)-\sum_{i=0}^{n} \sum_{j=0}^{k} D f^{-}(i, j)
$$

Throughout the remainder of this paper we shall use the notation and terminology introduced in Section 2 of [14]. We shall also use the notations $j(\alpha, \beta)$ for

$$
\{j(x, y) \mid x \in \alpha \text { and } y \in \beta\}
$$

and $\alpha \mid \beta$ for " $\alpha$ is separable from $\beta$ by disjoint r.e. sets".

## 3. Proof of the main theorem for $n=2$

Lemma 1. Let $f(x, y)$ be an increasing recursive function. Let $m, p, k, t$ be numbers such that $0 \leq m \leq p$ and $0 \leq k \leq t$. Define

$$
W=\{(i, j) \mid(i, j) \leq(p, l) \quad \text { and } \quad(i, j) \nsubseteq(m, k)\} .
$$

Then $\sum_{(i, j) \in W} D f(i, j) \geq 0$.
Proof.

$$
\begin{aligned}
\sum_{(i, j) \epsilon W} D f(i, j) & =\sum_{i=0}^{p} \sum_{j=0}^{l} D f(i, j)-\sum_{i=0}^{m} \sum_{j=0}^{k} D f(i, j) \\
& =f(p, l)-f(m, k) \geq 0
\end{aligned}
$$

Lemma 2. Let $A, B \in \Lambda_{R}-\varepsilon$. Let $f(x, y)$ be increasing and recursive. If $A+B \epsilon \Lambda_{R}$, then $f_{\Lambda}(A, B) \in \Lambda$.

Proof. By Theorem 2, we need only prove that

$$
\begin{equation*}
\sum_{(A+1, B+1)} D f^{+} \geq \sum_{(A+1, B+1)} D f^{-} \tag{2.1}
\end{equation*}
$$

Let $\alpha$ and $\beta$ be sets belonging to $A+1$ and $B+1$ respectively, and let $a_{n}$ and $b_{n}$ be regressive functions ranging over $\alpha$ and $\beta$ respectively. Since $A+B \in \Lambda_{R}$, $(A+1) \cdot(B+1) \in \Lambda_{R}$. Let $u_{n}$ be any regressive function ranging over $j(\alpha, \beta)$. Define

$$
\begin{aligned}
& \delta^{+}=U_{(i, j) \epsilon \varepsilon^{2}} j_{3}\left(a_{i}, b_{j}, \nu\left(D f^{+}(i, j)\right)\right) \\
& \delta^{-}=U_{(i, j) \in \varepsilon^{2}} j_{3}\left(a_{i}, b_{j}, \nu\left(D f^{-}(i, j)\right)\right)
\end{aligned}
$$

Clearly (2.1) will follow if we show the existence of a one-to-one partial recursive function $p(x)$ such that
(2.2) $\delta^{-} \subset$ domain $(p), p\left(\delta^{-}\right) \subset \delta^{+}$and $p\left(\delta^{-}\right) \mid \delta^{+}-p\left(\delta^{-}\right)$.

We shall prove the existence of a one-to-one correspondence $z \leftrightarrow g(z)$ which associates with each member $z$ of $\delta^{-}$a member $g(z)$ of $\delta^{+}$in such a manner that
(2.3) given $z$ we can effectively find $g(z)$ and vice-versa, and
(2.4) $g\left(\delta^{-}\right) \mid \delta^{+}-g\left(\delta^{-}\right)$.

This will complete the proof by Proposition 1 of [6].
The description of the correspondence $z \leftrightarrow g(z)$ requires the following series of definitions.

Let $q_{1}, q_{2}$ and $r$ be regressing functions of $a_{n}, b_{n}$ and $u_{n}$ respectively. Let $\zeta$ be any finite subset of $j(\alpha, \beta)$. Define

$$
\zeta_{1}=\{k(x) \mid x \in \zeta\}, \quad \zeta_{2}=\{t(x) \mid x \in \zeta\}
$$

Let $\zeta_{1}=\left\{a_{i(1)}, a_{i(2)}, \cdots, a_{i(s)}\right\}$ and $\zeta_{2}=\left\{b_{j(1)}, \cdots, b_{j(t)}\right\}$ where $i(1)<i(2)<$ $\cdots<i(s)$ and $j(1)<j(2)<\cdots<j(t)$.

Define

$$
\begin{aligned}
& P(\zeta)=\left\{j\left(a_{i}, b_{j}\right) \mid i<i(s) \text { and } j<j(t)\right\} \\
& R(\zeta)=\left\{r^{k}\left(j\left(a_{i}, b_{j}\right)\right) \mid a_{i} \in \zeta_{1} \text { and } b_{j} \in \zeta_{2} \text { and } k \in \varepsilon\right\} \\
& B(\zeta)=P(R(\zeta)), \quad B_{x}(\zeta)=\bigcup_{k=1}^{\infty} B^{k}(\zeta)
\end{aligned}
$$

We note that $\zeta \subset B_{x}(\zeta)$ and that $B^{i}(\zeta) \subset B^{i+1}(\zeta)$ for $i \geq 1$. It is clear that given $\zeta$ we can effectively find $B^{i}(\zeta)$ for any $i \geq 1$. Thus $B_{x}(\zeta)$ is an r.e. subset of the isolated product set $j(\alpha, \beta)$. It follows immediately that given a finite set $\zeta$ we can find $B_{x}(\zeta)$ by generating the sequence of sets $B(\zeta)$, $B^{2}(\zeta), \cdots$ until a repetition appears. Furthermore, the definition of $P$ yields the existence of a pair $\left(a_{k}, b_{l}\right) \in B_{x}(\zeta)$ such that

$$
\begin{equation*}
B_{x}(\zeta)=\left\{j\left(a_{i}, b_{j}\right) \mid(i, j) \leq(k, l)\right\} \tag{2.5}
\end{equation*}
$$

We now define sequences of sets $\left\{S_{n}\right\}$ and $\left\{W_{n}\right\}$ as follows:

$$
\begin{aligned}
S_{0} & =B_{x}\left(\left\{j\left(a_{0}, b_{0}\right)\right\}\right), \\
S_{n+1} & =B_{x}\left(S_{n} \sqcup\left\{j\left(a_{k}, b_{m}\right)\right\}\right)
\end{aligned}
$$

where $j\left(a_{k}, b_{m}\right)=u_{t}$ and $t=(\mu y)\left[u_{y} \oiint S_{n}\right]$,

$$
\begin{aligned}
& W_{0}=S_{0}, \quad W_{n+1}=S_{n+1}-S_{n} \\
& \tilde{W}_{n}=\left\{(i, j) \mid j\left(a_{i}, b_{j}\right) \in W_{n}\right\}
\end{aligned}
$$

Clearly given any member of $W_{n}$ we can use the given regressing functions to list the sets $S_{0}, \cdots, S_{n}$ and thus obtain $W_{0}, W_{1}, \cdots, W_{n}, \tilde{W}_{0}, \cdots, \widetilde{W}_{n}$. We note that by (2.5) and the definition of $W_{n+1}$, for each $n \in \varepsilon$ there exist numbers $m, p, k, l$ such that $(m, k) \leq(p, l)$ and

$$
\tilde{W}_{n+1}=\{(i, j) \mid(i, j) \leq(p, l) \quad \text { and } \quad(i, j) \nsubseteq(m, k)\}
$$

Hence by Lemma 2, for $n>0$,

$$
\begin{equation*}
\sum_{(i, j) \epsilon \tilde{w}_{n}} D f^{+}(i, j)-\sum_{(i, j) \in \tilde{w}_{n}} D f^{-}(i, j) \geq 0 \tag{2.6}
\end{equation*}
$$

For $n=0$, the left-hand expression in (2.6) is merely a member of the range of the recursive $f(x, y)$. Hence (2.6) holds for all $n \in \varepsilon$.

We now define the correspondence $z \leftrightarrow g(z)$. Let $z \epsilon \delta^{-}$. Let $n$ be the unique number such that $j\left(k_{1}(z), k_{2}(z)\right) \in W_{n}$. Define

$$
\begin{aligned}
& \gamma_{1}^{n}=U_{(i, j) \epsilon \tilde{\omega}_{n}} j_{3}\left(a_{i}, b_{j}, \nu\left(D f^{+}(i, j)\right)\right) \\
& \gamma_{2}^{n}=U_{(i, j) \in \tilde{W}_{n}} j_{3}\left(a_{i}, b_{j}, \nu\left(D f^{-}(i, j)\right)\right)
\end{aligned}
$$

Let $c_{1}<c_{2}<\cdots<c_{s}$ and $d_{1}<d_{2}<\cdots<d_{t}$ be finite sequences which array in increasing order all elements of the sets $\gamma_{1}^{n}$ and $\gamma_{2}^{n}$ respectively. By (2.6), $s \geq t$. Pair $c_{i}$ with $d_{i}, i=1, \cdots, t$. Since $z \in \gamma_{2}^{n}, z=d_{k}$ for some $k, 1 \leq k \leq t$. Thus $z$ is paired with $c_{k}$. Set $g(z)=c_{k}$. (Note that $g(z) \epsilon \delta^{+}$, as required.)

That $g$ has properties (2.3) and (2.4) is a consequence of the following observations. Let $y \in \delta^{+} \cup \delta^{-}$be given. Let $j\left(k_{1}(y), k_{2}(y)\right) \in W_{n}$. Then we can find all members of $W_{n}, \widetilde{W}_{n}, \gamma_{1}^{n}$ and $\gamma_{2}^{n}$, and list all pairs $(z, g(z)), z \in \gamma_{2}^{n}$. Examination of this list suffices to determine whether or not $y \in g\left(\delta^{-}\right)$. Furthermore if $y \in \delta^{-}\left(y \in g\left(\delta^{-}\right)\right)$, examination of the list suffices to determine the value of $g(y)\left(g^{-1}(y)\right)$.

Lemma 3. Let $A, B \in \Lambda_{R}-\varepsilon$. If $A \cdot B \in \Lambda_{R}$ and $h(x, y)$ is a recursive function such that $h(x, y) \geq 0$ for $x, y \in \varepsilon$, then $\sum_{A, B} h(x, y) \in \Lambda_{R}$.

Proof. Let $\alpha \in A$ and $\beta \in B$. Let $a_{n}$ and $b_{n}$ be regressive functions with ranges $\alpha$ and $\beta$ respectively. Let $p(x)$ be a regressing function for $a_{n}$ and $q(x)$ a regressing function for $b_{n}$. Let $p^{*}(x)$ and $q^{*}(x)$ be partial recursive functions such that $p^{*}\left(a_{n}\right)=n$ and $q^{*}\left(b_{n}\right)=n$ for $n \in \varepsilon$. Let $u_{n}$ be a regressive function ranging over $j(\alpha, \beta)$. Define $\tilde{a}_{n}=k\left(u_{n}\right)$ and $\tilde{b}_{n}=t\left(u_{n}\right)$. Then $j(\alpha, \beta)=\left\{j\left(\tilde{a}_{n}, \tilde{b}_{n}\right) \mid n \in \varepsilon\right\}$.

Define

$$
\gamma=\bigcup_{n=0}^{\infty} j_{3}\left(\tilde{a}_{n}, \tilde{b}_{n}, \nu\left[h\left(p^{*}\left(\tilde{a}_{n}\right), q^{*}\left(\tilde{b}_{n}\right)\right)\right]\right)
$$

Then $\gamma \in \sum_{A, B} h(x, y)$. We array the elements of $\gamma$ in the following fashion:

$$
\begin{array}{ll}
j_{3}\left(\tilde{a}_{0}, \tilde{b}_{0} 0\right), & \cdots, j_{3}\left(\tilde{a}_{0}, \tilde{b}_{0}, h\left(p^{*}\left(\tilde{a}_{0}\right), q^{*}\left(\tilde{b}_{0}\right)\right)-1\right) \\
\vdots & \\
j_{3}\left(\tilde{a}_{n}, \tilde{b}_{n}, 0\right) & \cdots, j_{3}\left(\tilde{a}_{n}, \tilde{b}_{n}, h\left(p^{*}\left(\tilde{a}_{n}\right), q^{*}\left(\tilde{b}_{n}\right)\right)-1\right) \\
j_{3}\left(\tilde{a}_{n+1}, \tilde{b}_{n+1}, 0\right), \cdots, j_{3}\left(\tilde{a}_{n+1}, \tilde{b}_{n+1}, h\left(p^{*}\left(\tilde{a}_{n+1}\right), q^{*}\left(\tilde{b}_{n+1}\right)\right)-1\right)
\end{array}
$$

It can easily be shown that we can "regress" through this array by proceeding from right to left in each row, and from the $(n+1)^{\text {st }}$ row to the $n^{\text {th }}$.

Theorem 3. Let $A, B \in \Lambda_{R}-\varepsilon$, and let $f(x, y)$ be an increasing recursive function. Then

$$
A+B \epsilon \Lambda_{R} \Rightarrow f_{\Lambda}(A, B) \in \Lambda_{R}
$$

Proof. Observe that the function

$$
h(x, y)=f(x, y)+\sum_{i=0}^{x} \sum_{j=0}^{y} D f^{-}(i, j)
$$

is a recursive function with the property that $D h(x, y) \geq 0$ for all $(x, y) \in \varepsilon^{2}$. Extending the defining equation of $h$ to $\Lambda_{R}$, we obtain

$$
h_{\Lambda}(A, B)=f_{\Lambda}(A, B)+\sum_{(A+1, B+1)} D f^{-}(x, y)
$$

By Lemma $2, f_{\Lambda}(A, B) \in \Lambda$, and by Lemma $3 \sum_{(A+1, B+1)} D f^{-}(x, y) \in \Lambda_{R}$ and $h_{\Lambda}(A, B) \in \Lambda_{R}$.
Hence $f_{\Lambda}(A, B) \leq h_{\Lambda}(A, B)$, and $f_{\Lambda}(A, B) \in \Lambda_{R}$.
Corollary. Let $f(x, y)$ be recursive and almost increasing. Then

$$
\begin{equation*}
A+B \in \Lambda_{R} \Rightarrow f_{\Lambda}(A, B) \in \Lambda_{R} \tag{*}
\end{equation*}
$$

Proof. Let $A$ and $B$ be regressive isols such that $A+B \in \Lambda_{R}$. Let $k$ be a number such that $f(x+k, y+k)$ is increasing. We distinguish two cases.

Case 1. Either $A$ or $B$ is finite. Suppose $A$ is finite. Define $g(y)=f(A, y)$ for $y \epsilon \varepsilon$. Then $g_{\Lambda}(B)=f_{\Lambda}(A, B)$. Since $g$ is a proper specification of $f, g$ is eventually increasing. Hence by Theorem 4 of $[3], f_{\Lambda}(A, B) \in \Lambda_{R}$. The proof is completely similar if $B$ is finite.

Case 2. $\quad A$ and $B$ are both infinite isols. Let $g(x, y)=f(x+k, y+k)$. Then $f_{\Lambda}(A, B)=g_{\Lambda}(A-k, B-k)$. Since $g(x, y)$ is recursive and increasing, Theorem 3 yields $A+B \in \Lambda_{R}$.

Notation. We shall write $(x, y)<(z, w)$ if $(x, y) \leq(z, w)$ and $(x, y) \neq(z, w)$.

Lemma 4. Let $f(x, y)$ be a recursive function which is not eventually increasing. Then there exist infinite regressive isols $A$ and $B$ such that $A+B \in \Lambda_{R}$ and $f_{\Lambda}(A, B) \in \Lambda^{*}-\Lambda$.

Proof. Since $f(x, y)$ is not eventually increasing, we can effectively generate a strictly increasing sequence

$$
\left\{\left(n_{i}, m_{i}, s_{i}, t_{i}\right) \mid i=0,1, \cdots\right\}
$$

of four-tuples such that for $i \in \varepsilon$,

$$
\begin{gather*}
\left(n_{i}, m_{i}\right)<\left(s_{i}, t_{i}\right)<\left(n_{i+1}, m_{i+1}\right), \quad f\left(n_{i}, m_{i}\right)>f\left(s_{i}, t_{i}\right)  \tag{2.7}\\
\lim _{i \rightarrow \infty} n_{i}=\lim _{i \rightarrow \infty} m_{i}=\infty
\end{gather*}
$$

We note that each of the functions $n_{i}, m_{i}, s_{i}$ and $t_{i}$ is recursive and increasing. We define recursive functions $a_{i}$ and $b_{i}$ as follows:

$$
\begin{aligned}
& a_{0}=n_{0}, \quad b_{0}=m_{0} \\
& a_{1}=s_{0}-n_{0}, \quad b_{1}=t_{0}-m_{0} \\
& \vdots \\
& a_{2 k}=n_{k}-s_{k-1}, \quad \vdots \\
& a_{2 k+1}=s_{k}-n_{k}, \quad b_{2 k+1}=m_{k}-t_{k-1} \\
& t_{k}-m_{k}
\end{aligned}
$$

We note that for $k \in \varepsilon$,

$$
\begin{equation*}
n_{k}=\sum_{i=0}^{2 k} a_{i}, \quad s_{k}=\sum_{i=0}^{2 k+1} a_{i}, \quad m_{k}=\sum_{i=0}^{2 k} b_{i}, \quad t_{k}=\sum_{i=0}^{2 k+1} b_{i} \tag{2.8}
\end{equation*}
$$

For $T \in \Lambda_{R}$, we define $A_{T}=\sum_{T} a_{i}$ and $B_{T}=\sum_{T} b_{i}$. Since

$$
A_{T}+B_{T}=\sum_{T}\left(a_{i}+b_{i}\right), A_{T}+B_{T} \in \Lambda_{R} \quad \text { for } T \in \Lambda_{R} .
$$

For $i \in \varepsilon$, define

$$
\begin{aligned}
W_{i} & =\left\{(x, y) \mid(x, y) \leq\left(s_{i}, t_{i}\right) \text { and }(x, y) \nsubseteq\left(n_{i}, m_{i}\right)\right\}, \\
S_{0} & =\left\{(x, y) \mid(x, y) \leq\left(n_{0}, m_{0}\right)\right\}, \\
S_{i+1} & =\left\{(x, y) \mid(x, y) \leq\left(n_{i+1}, m_{i+1}\right) \text { and }(x, y) \neq\left(s_{i}, t_{i}\right) .\right.
\end{aligned}
$$

We note that the sequence of sets $S_{0}, W_{0}, S_{1}, W_{1}, \cdots$ is an infinite sequence of mutually disjoint sets whose union is $\varepsilon^{2}$, and that for $i \in \varepsilon$,

$$
\begin{align*}
S_{0} \sqcup W_{0} \sqcup \cdots \sqcup W_{i-1} \sqcup S_{i} & =\left\{(x, y) \mid(x, y) \leq\left(n_{i}, m_{i}\right)\right\} \\
S_{0} \sqcup W_{0} \sqcup \cdots \sqcup S_{i} \sqcup W_{i} & =\left\{(x, y) \mid(x, y) \leq\left(s_{i}, t_{i}\right)\right\} \tag{2.9}
\end{align*}
$$

We define a recursive function $g(x)$ mapping $\varepsilon$ to $\varepsilon^{*}$ by

$$
g(2 k)=\sum_{(i, j) \epsilon s_{k}} D f(i, j), \quad g(2 k+1)=\sum_{(i, j) \epsilon W_{k}} D f(i, j)
$$

By (2.8) and (2.9),

$$
\begin{align*}
\sum_{i=0}^{2 k} g(i) & =f\left(n_{k}, m_{k}\right)=f\left(\sum_{i=0}^{2 k} a_{i}, \sum_{i=0}^{2 k} b_{i}\right) \\
\sum_{i=0}^{2 k+1} g(i) & =f\left(s_{k}, t_{k}\right)=f\left(\sum_{i=0}^{2 k+1} a_{i}, \sum_{i=0}^{2 k+1} b_{i}\right) \tag{2.10}
\end{align*}
$$

Thus for $k \in \varepsilon$,

$$
\begin{equation*}
\sum_{i=0}^{k} g(i)=f\left(\sum_{i=0}^{k} a_{i}, \sum_{i=0}^{k} b_{i}\right) \tag{2.11}
\end{equation*}
$$

By (2.10) and (2.7), for $k \in \varepsilon$,

$$
g(2 k+1)=f\left(s_{k}, t_{k}\right)-f\left(n_{k}, m_{k}\right)<0
$$

Thus $g$ assumes negative values infinitely often. Extending (2.11) to $\Lambda_{R}$, we obtain for $T \in \Lambda_{R}$,

$$
\sum_{T+1}^{*} g(i)=f_{\Lambda}\left(\sum_{T+1} a_{i}, \sum_{T+1} b_{i}\right)=f_{\Lambda}\left(A_{T+1}, B_{T+1}\right)
$$

It follows from statement (30) in the proof of Theorem 4 of [3] that there exists a regressive isol $T+1$ such that $\sum_{T+1}^{*} g(i) \epsilon \Lambda^{*}-\Lambda$. For such an isol $T+1$,

$$
f_{\Lambda}\left(A_{T+1}, B_{T+1}\right) \in \Lambda^{*}-\Lambda
$$

while $A_{T+1}+B_{T+1} \in \Lambda_{R}$.
Corollary. If $f(x, y)$ is a recursive function which is not eventually increasing, there exist recursive functions $a_{n}$ and $b_{n}$ and an isol $T \in \Lambda_{R}$ such that

$$
f_{\Lambda}\left(\sum_{T} a_{n}, \sum_{T} b_{n}\right) \in \Lambda^{*}-\Lambda
$$

Theorem 4. Let $f(x, y)$ be a recursive function. f has the property

$$
\begin{equation*}
A+B \in \Lambda_{R} \Rightarrow f_{\Lambda}(A, B) \in \Lambda_{R} \tag{*}
\end{equation*}
$$

if and only if $f$ is almost increasing.
Proof. We have already shown that the second condition implies the first. Suppose that $f$ is not almost increasing. If $f$ is not eventually increasing, $f$ does not have the property (*) by Lemma 4. If there is a number $k$ such that the proper specification $f(x, k)$ is not eventually increasing, by Theorem 4 of [3] there exists a regressive isol $T$ such that $f_{\Lambda}(T, k) \notin \Lambda_{R}$. Hence $f$ does not have the property (*). The case in which some specification $f(k, y)$ is not eventually increasing is handled similarly.

## 4. Applications

Myhill and Nerode have shown that if $\Phi\left(x_{1}, \cdots, x_{n}\right)$ is a quantifier-free Horn formula built up from equations between almost combinatorial recursive functions of $x_{1}, \cdots, x_{n}$ and $\Phi\left(x_{1}, \cdots, x_{n}\right)$ is true for all natural numbers, then $\Phi\left(x_{1}, \cdots, x_{n}\right)$ is true for all isols $x_{1}, \cdots, x_{n}$ [11], [12]. Using this result and the fact that any recursive function can be expressed as the difference of two recursive combinatorial functions, one can easily prove the following proposition.

Proposition 2. Let $\mathfrak{\Re}\left(x_{1}, \cdots, x_{n}\right)$ be a quantifier free Horn formula built up from equations between almost increasing recursive functions of $x_{1}, \cdots, x_{n}$. If $\mathfrak{A}\left(x_{1}, \cdots, x_{n}\right)$ is true for all natural numbers, then $\mathfrak{H}\left(x_{1}, \cdots, x_{n}\right)$ is true for all isols $x_{1}, \cdots, x_{n}$ such that $x_{1}+x_{2}+\cdots+x_{n} \in \Lambda_{R}$.

Corollary. Let $\mathfrak{M}\left(x_{1}, \cdots, x_{n}, y\right)$ be a quantifier free Horn formula built up from equations between almost increasing recursive functions of $x_{1}, \cdots, x_{n}, y$. If $(\exists y) \mathfrak{H}\left(x_{1}, \cdots, x_{n}, y\right)$ is true for all natural numbers $x_{1}, \cdots, x_{n}$ and has an
almost increasing recursive Skolem function $f\left(x_{1}, \cdots, x_{n}\right)$, then $(\exists y) \mathfrak{H}\left(x_{1}, \cdots\right.$, $\left.x_{n}, y\right)$ is true for all isols $x_{1}, \cdots, x_{n}$ such that $x_{1}+x_{2}+\cdots+x_{n} \in \Lambda_{R}$.

Proof. Immediate from Proposition 2 and the fact that the class of almost increasing functions is closed under composition. We leave verification of the latter fact to the reader.

Let $p r(n)$ denote the primitive recursive function which enumerates the prime numbers in increasing order. J. C. E. Dekker proved in [8] that $X^{p r(n)} \equiv X(\bmod p r(n))$ for $X \epsilon \Lambda, n \epsilon \varepsilon$. J. Barback has shown in [2] that the class of all prime regressive isols has cardinality $c$ and properly contains the class $\left\{p r_{\Lambda}(T) \mid T \in \Lambda_{R}\right\}$. Hence it is natural to ask if the above congruence holds if $\operatorname{pr}(n)$ is replaced by an infinite regressive prime isol. The following corollary describes a class of regressive isols for which the congruence is satisfied.

Corollary. Let $A$ and $B$ be regressive isols such that $A+B \in \Lambda_{R}$. Let $P=p r_{\Lambda}(B) . \quad$ Then $A^{P} \equiv A(\bmod P)$.

Proof. Consider the formula

$$
\mathfrak{N}=(\exists w)\left[(x+1)^{p r(n)}=w \cdot p r(n)+(x+1)\right] .
$$

This sentence is satisfied by all $x$ and $n$ in $\varepsilon$ and has the recursive increasing Skolem function

$$
W(x, n)=\frac{(x+1)^{p r(n)}-(x+1)}{p r(n)}
$$

By the previous corollary $\mathfrak{Q}$ is satisfied by all $A, B \in \Lambda_{R}$ such that $A+B \in \Lambda_{R}$. This completes the proof in case $A \neq 0$; if $A=0$, the congruence is clearly satisfied.
J. Barback has shown in [2] that there exist $A, B \in \Lambda_{R}$ such that

$$
\min _{\Lambda}(A, B) \not \equiv A+B,
$$

but that the restriction $A+B \in \Lambda_{R}$ is sufficient to guarantee that $\min _{\Lambda}(A, B) \leq A+B$. The following proposition shows that this restriction is not sufficient to guarantee that $\min _{\Lambda}(A, B) \leq \max _{\Lambda}(A, B)$.

Proposition 3. There exist isols $A, B$ such that

$$
A+B \in \Lambda_{R} \quad \text { and } \min _{\Lambda}(A, B) \neq \max _{\Lambda}(A, B) .
$$

Proof. Let $g(x, y)=\max (x, y)-\min (x, y)$. Then the identity

$$
\begin{equation*}
\min _{\Lambda}(A, B)+g_{\Lambda}(A, B)=\max _{\Lambda}(A, B) \tag{**}
\end{equation*}
$$

holds in $\Lambda^{*}$ for $A, B \in \Lambda_{\boldsymbol{R}}$. However $g(x, y)$ is not eventually increasing. Let $A, B \in \Lambda_{R}$ be such that

$$
A+B \in \Lambda_{R} \quad \text { and } \quad g_{\Lambda}(A, B) \in \Lambda^{*}-\Lambda
$$

Since $A+B \in \Lambda_{R}$, both $\min _{\Lambda}(A, B)$ and $\max _{\Lambda}(A, B)$ are regressive. By the identity (**), $\min _{\Lambda}(A, B) \not \max _{\Lambda}(A, B)$.

By the corollary to Lemma 4, Proposition 3 may be strengthened as follows.
Proposition 4. There exist recursive functions $a_{n}$ and $b_{n}$ and a regressive isol T such that

$$
\min _{\Lambda}\left(\sum_{T} a_{n}, \sum_{T} b_{n}\right) \neq \max _{\Lambda}\left(\sum_{T} a_{n}, \sum_{T} b_{n}\right) .
$$

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