ON MONOTONE VS. NONMONOTONE INDUCTION

BY LEO A. HARRINGTON AND ALEXANDER S. KECHRIS

Communicated by S. Feferman, May 11, 1976

1. Introduction. For definitions and notation in what follows, see [4] and [5]. If A is an infinite set and $\varphi(y_1 \cdots y_n, R, Y_1 \cdots Y_m) = \varphi(\overline{y}, R, \overline{Y})$ is a second order relation on A, we call φ operative if R is n-ary. For such a φ let

$$I_{\varphi}^{\xi} = \bigcup_{\eta < \xi} I_{\varphi}^{\eta} U \bigg\{ (\overline{y}, \ \overline{Y}) \colon \varphi \bigg(\overline{y}, \bigg\{ \overline{y} \colon (\overline{y}, \ \overline{Y}) \in \bigcup_{\eta < \xi} I_{\varphi}^{\eta} \bigg\}, \ \overline{Y} \bigg) \bigg\} \quad \text{and} \quad I_{\varphi} = \bigcup_{\xi} I_{\varphi}^{\xi}.$$

If F is a collection of second order relations (for simplicity collection of operators) on A, then F-IND² is the class of all second order relations of the form $\psi(\overline{x}, \overline{Y}) \Leftrightarrow I_{\varphi}(\overline{a}, \overline{x}, \overline{Y})$, for some operative $\varphi(\overline{u}, \overline{x}, R, \overline{Y})$ in F and constants \overline{a} from A. As in [5] F-IND is the class of all relations on A which are in F-IND². We let $F^{m \circ n}$ be the collection of all operative $\varphi(\overline{v}, R, \overline{Y})$ in F which are monotone on R and we put $\neg F = \{ \neg \varphi : \varphi \in F \}$. A collection of operators F on A is adequate if it contains all the $\Pi_1^0(C)$ second order relations, where C is a coding scheme on A and is closed under $\land, \lor, \exists A$ and trivial combinatorial substitutions. Let $WF(S) \Leftrightarrow S$ be a well-founded relation on $A \Leftrightarrow \neg \exists a_0 a_1 a_2 \cdots \forall i(a_{i+1}, a_i) \in S$.

THEOREM 1. Let F be an adequate collection of operators on an infinite set A. If $WF \in \neg F$ and $\neg F \subseteq F^{mon}$ -IND², then F-IND² = F^{mon} -IND².

2. Elementary induction. Let EL be the collection of all the elementary second order relations on a structure $A = \langle A, R_1 \dots R_l \rangle$ and let EL^+ be the subcollection of EL^{mon} consisting of all operative $\varphi(\overline{x}, R, \overline{Y})$ which are definable by positive in R elementary formulas. One usually writes EL^+ -IND² = IND² and EL^+ -IND = IND. Clearly IND² $\subseteq EL^{mon}$ -IND² $\subseteq EL$ -IND² and it is well known that IND² is a tiny part of EL-IND² for (say) almost acceptable A's. By a basic result of Kleene and Spector for ω and Barwise-Gandy-Moschovakis in general (see [4, §8A]), on every countable almost acceptable structure, IND² = EL^{mon} -IND² (= Π_1^1). On the other hand, letting $WF^n(S) \Leftrightarrow S$ is a 2n-ary relation on A which is well founded (viewed as binary on A^n), we have

COROLLARY 1. Let A be an infinite structure such that each WF^n is elementary. Then EL^{mon} -IND² = EL-IND².

AMS (MOS) subject classifications (1970). Primary 02F27.

A more detailed level-by-level version of Corollary 1 is the following, where we just write Σ_m^0 , Π_m^0 instead of $\Sigma_m^0(C)$, $\Pi_m^0(C)$, where C is a hyperelementary coding scheme on A.

COROLLARY 2. Let A be an almost acceptable structure. If $m \ge 2$ and $WF \in \Pi_m^0$, then for all $n \ge m$, Σ_n^0 -IND² = $(\Sigma_n^0)^{m \circ n}$ -IND².

So, for example, in the structure of analysis R this says that Σ_n^1 monotone operators on R inductively define the same relations as arbitrary Σ_n^1 operators, when $n \ge 2$. Similarly for Σ_n^1 . The following rather curious result can be also established by the methods used to prove Theorem 1. If $A = \langle A, R_1 \cdots R_l \rangle$ is a structure, by an *elementary quantifier* Q on A we understand a quantifier on A which viewed as a second-order relation is elementary.

THEOREM 2. Let A be an acceptable structure in which WF is elementary. There is an elementary quantifier Q on A such that for every inductive relation R on A, there is an inductive relation R^* on A such that A $R(x) \Leftrightarrow QyR^*(x,y)$.

This should be compared with a result of Moschovakis [3] in higher type recursion, where "inductive" is replaced by "semirecursive in a total object of type ≥ 3 " and Q becomes the existential quantifier (on an appropriate space).

REMARKS. (i) We conjecture that in Theorem 1 (and correspondingly in Corollary 1) the hypothesis $WF \in \mathbb{T}$ F can be weakened to $WF \in \mathbb{T}$ ($F^{m \, on}$ -IND²). (ii) In a direction opposite to that of Corollary 1 one has the following theorem of Nyberg (unpublished): Let A be almost acceptable. If IND $\not\subseteq -(EL^{m \, on}$ -IND), then $EL^{m \, on}$ -IND = IND. Thus for most structures occurring in practice, $EL^{m \, on}$ -IND is either IND or EL-IND.

3. Further corollaries and applications to Spector classes. An immediate consequence of Theorem 1 is also the following result of Harrington and Moschovakis [2]. (Given a structure A and a quantifier Q on A we abbreviate by Q-IND the class of second order relations which are positive $L^A(Q)$ -inductive (see [4, p. 49]).

COROLLARY 3. (Harrington-Moschovakis [2]). Let A be an almost acceptable structure and let Q be a quantifier on A. If $F = \neg (Q - IND^2)$, then $F - IND^2 = F^{mon} - IND^2$.

This generalizes a result of Grilliot to the effect that over ω , Σ_1^1 -IND² = $(\Sigma_1^1)^{\text{mon}}$ -IND². The original proof of Corollary 2 in [2] yields the stronger statement that for $F = \neg (Q\text{-IND}^2)$, $F\text{-IND}^2 = F^{\text{pos}}\text{-IND}^2$ and also shows that $F\text{-IND}^2 = Q^+\text{-IND}^2$, where Q^+ is the *next quantifier* of Q (see [1]). Turning now to Spector classes we can obtain the following, where the notions involved are explained in [5].

THEOREM 3. Let Γ be a Spector class on A, and let F be a reasonable,

nonmonotone class of operators on A closed under \exists^A . If $WF \in \neg F$, then Γ is F-compact iff Γ is F_*^{mon} -compact, where $F_*^{mon} = \{\varphi(R): \varphi \in F, \varphi \text{ monotone}\}$. In particular if F is typical, nonmonotone, F^{mon} -IND is a Spector class iff F^{mon} -IND = F-IND.

Further applications of the methods developed here to the theory of "second order" Spector classes as well as details and proofs of the results announced here will appear elsewhere.

REFERENCES

- 1. P. Aczel, Quantifiers, games and inductive definitions, Proc. Third Scandinavian Logic Sympos. (S. Kanger, Editor), North-Holland, Amsterdam; American Elsevier, New York, 1975, pp. 1-14. MR 51 #5256.
- 2. L. A. Harrington and Y. N. Moschovakis, On positive induction vs. nonmonotone induction, Mimeographed notes, 1975.
- 3. Y. N. Moschovakis, Hyperanalytic predicates, Trans. Amer. Math. Soc., 129 (1967), 249-282. MR 38 #4308.
- 4. ——, Elementary induction on abstract structures, North-Holland, Amsterdam, 1974.
- 5. ——, On nonmonotone inductive definability, Fund. Math. 82 (1974), 39-83. MR 50 #6853.

DEPARTMENT OF MATHEMATICS, UNIVERSITY OF CALIFORNIA, BERKELEY, BERKELEY, CALIFORNIA 94720

DEPARTMENT OF MATHEMATICS, CALIFORNIA INSTITUTE OF TECHNOLOGY, PASADENA, CALIFORNIA 91125