A REMARK ON THE INTERSECTION OF TWO LOGICS

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The intuitionistic logic LJ and Curry's LD (cf. [1], [2]) are logics stronger than Johansson's minimal logic LM (cf. [3]) by the axiom schemes $\land \rightarrow x$ and $y \lor (y \rightarrow \land)$, respectively. However, LM can not be taken literally as the intersection of these two logics LJ and LD, which is stronger than LM by the axiom scheme $(\land \rightarrow x) \lor y \lor (y \rightarrow \land)$. In pointing out this situation, Prof. K. Ono suggested me to investigate the general feature of the intersection of any pair of logics. In this paper, I will show that the same situation occurs in general. I wish to express my thanks to Prof. K. Ono for his kind guidance.

Let A be a logic having logical constants, *implication* (\rightarrow) and *disjunction* (\vee) (and *universal quantification* () for predicate logics), together with all such inference rules with respect them that are admitted in the intuitionistic logic (cf. [5], p. 81). For any logic L, let us denote by Π_L the class of all provable propositions in L.

THEOREM. Let B, C, and D be the logics formed from A by adjoining the axiom schemes

- (1) $(u_1) \cdots (u_p) f(x_1, \ldots, x_s), \quad (p = 0, 1, 2, \ldots),$
- (2) $(v_1) \cdots (v_q) g(y_1, \ldots, y_t), \qquad (q = 0, 1, 2, \ldots),$

(3) $(u_1)\cdots(u_p)f(x_1,\ldots,x_s)\vee(v_1)\cdots(v_q)g(y_1,\ldots,y_t),$

 $(p, q = 0, 1, 2, \ldots),$

respectively; where u_i 's and v_j 's are object variables (p = q = 0 for proposition logics), $(u_1) \cdots (u_p) f(x_1, \ldots, x_s)$ and $(v_1) \cdots (v_q) g(y_1, \ldots, y_t)$ are expressible in A, x_i 's and y_j 's are metalogical variables for propositions, predicates, or relations, and $s \le t$. Then,

 $I. \qquad \qquad \Pi_D = \Pi_B \cap \Pi_C.$

II. B and C formed from D by adjoining the axiom schemes

$$(4)_{\mu} (w_1) \cdots (w_r) (g(y_1, \ldots, y_t) \to f(y_{\mu(1)}, \ldots, y_{\mu(s)})), \quad (r = 0, 1, 2, \ldots),$$

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