

THE UNDECIDABILITY OF THE RECURSIVELY ENUMERABLE DEGREES

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Let $\leq$ be Turing reducibility between subsets of ω and let $\mathcal{R}$ be the collection of all recursively enumerable subsets of ω . For A in $\mathcal{R}$, A is the Turing degree of A ; and $\mathcal{R} = \{A; A \text{ is in } \mathcal{R}\}$. (See [2] for any unexplained notions or notation used above, or below). We also identify $\mathcal{R}$ with the structure $\langle \mathcal{R}, \leq \rangle$.

We wish to announce the result

THEOREM. *The first order theory of $\mathcal{R}$ is undecidable.*

To prove this we show that the theory of partial orderings is reducible to the theory of $\mathcal{R}$, as follows.

Recall, $\mathcal{R}$ is an upper-semi-lattice; thus for a, b in $\mathcal{R}$, their join $a \vee b$ is in $\mathcal{R}$.

LEMMA. *For $\mathcal{P}$ a partial ordering recursive in (say) $0'$, there are a, b_p ($p \in \mathcal{P}$), c, d, e in $\mathcal{R}$ such that*

- (i) $b_p \leq a$,
- (ii) $a \not\leq (b_p \vee c)$,
- (iii) for each $Z \leq a$, either
 - (α) $a \leq (Z \vee c)$, or $\exists p \in \mathcal{P}$ such that
 - (β) _{p} $Z \leq (b_p \vee d)$,
- (iv) for $p \neq q$, $b_p \not\leq (b_q \vee d)$,
- ($\sim$) for p, q in $\mathcal{P}$, $p \leq q$ iff $b_p \leq (b_q \vee d \vee e)$.

Now, for a, b, c, d in $\mathcal{R}$ let $\phi(a, b, c, d) \equiv (b \leq a)$ and $(a \not\leq (b \vee c))$ and $(\neg \exists Z (Z \leq a \text{ and } a \not\leq (Z \vee c) \text{ and } (Z \vee d) > (b \vee d)))$. For a, c, d, e in $\mathcal{R}$, let $Q(a, c, d, e) = \{b \vee d \vee e; \mathcal{R} \models \phi(a, b, c, d)\}$. For a, c, d, e as in the lemma, $Q(a, c, d, e) = \{b_p \vee d \vee e; p \in \mathcal{P}\}$, and $\langle Q(a, c, d, e), \leq \rangle$ is isomorphic to $\mathcal{P}$.

Thus for ψ a sentence of the language of partial orderings: ψ is true of some partial ordering iff (by the usual proof of the completeness theorem) $\mathcal{P} \models \psi$ for some $\mathcal{P}$ recursive in $0'$ iff (by the lemma) $\exists a, c, d, e$ in $\mathcal{R}$ ($\langle Q(a, c, d, e), \leq \rangle \models \psi$).

The lemma is proven, of course, by a priority argument. The type of priority argument used can best be described as an infinite injury argument with a finite injury priority argument on top of it. This kind of construction was first

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