

## WORD PROBLEMS

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**Introduction.** In studying fundamental groups of manifolds, Dehn [4] in 1911 investigated some special cases of a problem which is now known as *the word problem for groups*. Let  $G$  be a group generated by a finite set of elements  $a, b, c, \dots$ . Each element in  $G$  is a product of these generators and their inverses, for example,  $a^{-1}bacb^{-1}c$ . We call such expressions  $w(a, b, c, \dots)$  *words* in the generators.  $G$  is assumed to be given by a finite set of equations or *defining relations*  $r(a, b, c, \dots) = 1$  which the generators satisfy. The question Dehn proposed was to find a uniform test or mechanical procedure (i.e. an *algorithm*) which enables us to decide whether  $w = 1$  for an arbitrary word  $w(a, b, c, \dots)$  in  $G$ . We say that the word problem is *solvable* for the group  $G$  if there is such an algorithm.

A different version of the same kind of problem was posed by Thue a year or two later [35]. Consider an "abstract language"  $L$  having a finite alphabet  $a, b, c, \dots$ . A word in  $L$  is a finite sequence of symbols from the alphabet, for example,  $baabccba$ . A language  $L$  (usually called a *Thue system*) is defined on the alphabet  $a, b, c, \dots$  by a *dictionary*, a finite set of pairs of words. If a word  $w$  is of the form  $usv$ , where  $u, v$  are words and  $s$  is a word which occurs in the dictionary paired with a word  $t$ , then we say that  $w$  can be *transformed* by the dictionary into the word  $utv$  and we write  $usv \rightarrow utv$ . If there is a finite sequence of transformations  $w \rightarrow \dots \rightarrow w'$  connecting two words  $w, w'$ , then we say that  $w, w'$  are *equivalent* in  $L$ . Thue's problem, which we may call *the word problem for the language  $L$* , asks for an algorithm for deciding whether two words in the language are equivalent. We say that the word problem is *solvable* for  $L$  if there is such an algorithm.

These two problems are examples of the same general situation. We have an algebra  $\mathcal{A}$  which is generated by elements  $a, b, c, \dots$ . The elements of the algebra are represented by expressions (*words*) involving the generators and the operations. The algebra satisfies certain axioms and is characterized by certain basic relations  $r = r'$  where the  $r, r'$  are words. We wish to find some effective test for deciding whether two words  $w(a, b, c, \dots), w'(a, b, c, \dots)$  represent the same element of the algebra, i.e. whether  $w = w'$  follows from the axioms and the relations  $r = r'$ . The question of the existence of such an algorithm is called *the word problem* for the algebra. Obviously the word problem for groups is of this type. The word problem for the language  $L$  becomes such a question if we view  $L$  as a semigroup given by generators

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