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10. Simple Type Theory of Gentzen Style with the Inference of Extensionality

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(Comm. by Zyoiti SUETUNA, M. J. A., Feb. 12, 1968)

The definitions of types and expressions are the same as in [2]. We also refer to expressions of type 1 as formulas. As in LK or GLC, the form

$$A_1, \cdots, A_m \rightarrow B_1, \cdots, B_n,$$

where $A_1, \dots, A_m, B_1, \dots, B_n(m, n \ge 0)$ are formulas, is called a sequent. The inference rules of our system are as follows:

(I) Structual inference rules

$$\frac{\Gamma \rightarrow \Delta}{A, \Gamma \rightarrow \Delta} \qquad \frac{\Gamma \rightarrow \Delta}{\Gamma \rightarrow \Delta, A} \\
\frac{A, A, \Gamma \rightarrow \Delta}{A, \Gamma \rightarrow \Delta} \qquad \frac{\Gamma \rightarrow \Delta, A, A}{\Gamma \rightarrow \Delta, A} \\
\frac{\Gamma, A, B, \Pi \rightarrow \Delta}{\Gamma, B, A, \Pi \rightarrow \Delta} \qquad \frac{\Gamma \rightarrow \Delta, A, B, \Lambda}{\Gamma \rightarrow \Delta, B, A, A} \\
(Cut) \qquad \frac{\Gamma \rightarrow \Delta, A \quad A, \Pi \rightarrow \Lambda}{\Gamma, \Pi \rightarrow \Delta, \Lambda}$$

(II) Inference rules on logical symbols

$$\frac{\Gamma \rightarrow \varDelta, A}{\neg A, \Gamma \rightarrow \varDelta} \qquad \frac{A, \Gamma \rightarrow \varDelta}{\Gamma \rightarrow \varDelta, \neg A} \\
\frac{A, \Gamma \rightarrow \varDelta}{A \lor B, \Gamma \rightarrow \varDelta} \\
\frac{A, \Gamma \rightarrow \varDelta}{A \lor B, \Gamma \rightarrow \varDelta} \\
\frac{\Gamma \rightarrow \varDelta, A}{F \rightarrow \varDelta, A \lor B} \qquad \frac{\Gamma \rightarrow \varDelta, B}{\Gamma \rightarrow \varDelta, A \lor B} \\
\frac{A(a^{\tau}), \Gamma \rightarrow \varDelta}{\exists x^{\tau} A(x^{\tau}), \Gamma \rightarrow \varDelta}, \qquad \frac{\Gamma \rightarrow \varDelta, A(e^{\tau})}{\Gamma \rightarrow \varDelta, \exists x^{\tau} A(x^{\tau})},$$

where a^{τ} does not occur in the lower sequent. where e^{τ} is an arbitrary expression of type τ .

(III) Inference of comprehension

$$\frac{A(e_1^{\tau_1}, \cdots, e_n^{\tau_n}), \Gamma \longrightarrow \Delta}{(e_1^{\tau_1}, \cdots, e_n^{\tau_n} \in \lambda x_1^{\tau_1} \cdots x_n^{\tau_n} A(x_1^{\tau_1}, \cdots, x_n^{\tau_n})), \Gamma \longrightarrow \Delta}$$
$$\frac{\Gamma \longrightarrow \Delta, A(e_1^{\tau_1}, \cdots, e_n^{\tau_n})}{\Gamma \longrightarrow \Delta, (e_1^{\tau_1}, \cdots, e_n^{\tau_n} \in \lambda x_1^{\tau_1} \cdots x_n^{\tau_n} A(x_1^{\tau_1}, \cdots, x_n^{\tau_n}))}$$

(IV) Inference of extensionality