

SOME PROOFS OF RELATIVE COMPLETENESS IN MODAL LOGIC

M. J. CRESSWELL

In this paper we adopt for a number of modal systems a method of proving completeness used in [1] (to prove the completeness of S4 relative to T).¹ We prove the completeness of S7, S8 relative to S3, of S3, S6 relative to S2 and of S2 relative to E2.

The system E2 was proposed by E. J. Lemmon² and can be axiomatized as follows;

- A1 *If α is a PC-tautology then $\vdash \alpha$*
 A2 $Lp \supset p$
 A3 $L(p \supset q) \supset (Lp \supset Lq)$

(with uniform substitution or replaced by the equivalent schemata)

- A4 $\vdash \alpha \supset \beta \rightarrow \vdash L\alpha \supset L\beta$
 A5 (**MP**) $\vdash \alpha, \vdash \alpha \supset \beta \rightarrow \vdash \beta$

E2 is in fact the system T (as axiomatized in [4], pp. 533-535) without the rule of necessitation. We obtain S2 by replacing A4 by,

- A6 $\vdash L(\alpha \supset \beta) \rightarrow \vdash L(L\alpha \supset L\beta)$

and adding;

- A7 *If α is a PC-tautology or an axiom then $\vdash L\beta$* ³

For S3 we replace A6 and A3 by

- A8 $L(p \supset q) \supset L(Lp \supset Lq)$

For S6 we add to S2, A9 *MMp*

For S7 we add to S3, A10 *MMp*

For S8 we add to S3, A11 *LMMp*

where A7 does not apply to A9 - A11.

An account of validity for these systems has recently been given by Saul Kripke [7] and it is essentially this account we shall use. We define an E2 model as an ordered quadruple $\langle V, W, R, x_1 \rangle$ where W is a set of objects (worlds), $x_1 \in W$ and R a quasi-reflexive relation over W . By this is meant that for any $x_i \in W$ if any $x_j R x_i$ then $x_i R x_i$. Members x_i of W such that $x_i R x_i$ are called normal.⁴ Quasi-reflexiveness ensures that no world is