

## COMBINATORY AND PROPOSITIONAL LOGIC

DAVID MEREDITH

The relationship between combinatory and propositional logic is dealt with at length in [1] and tangentially in [2]. The present paper adds nothing essentially new to previous results. It does, however, offer a straightforward procedure, which for any  $\lambda$ -expression in normal form will either lead to its propositional correspondent or determine that this is null. Section 1 presents the hypothesis upon which the correspondence between  $\lambda$ -expressions and propositional formulae is based; our translation procedure is described in section 2, and illustrated in section 3.

1 Hypothesis. In dealing with  $\lambda$ -expressions we assume Church's rules and conventions as given in [3]. With respect to propositional formulae, ' $\Lambda$ ' denotes the null class of formulae, and ' $\Gamma$ ' is used for C. A. Meredith's operator  $D$ : ' $\Gamma PQ$ ' denotes the most general result that can be obtained when *Modus Ponens* is applied with  $P$ , or some substitution in it, as major premiss, and  $Q$ , or some substitution in it, as minor premiss. ' $\sim$ ' denotes correspondence between a  $\lambda$ -expression and a propositional formula. Our basic hypothesis is the following.

*Hypothesis* Where  $L$ ,  $M$  and  $N$  are  $\lambda$ -expressions,  $P$ ,  $Q$  and  $R$  are propositional formulae, and  $\Sigma$  is an operation under the substitution rule which may be null.

1. Let  $L \sim P$ , then for  $M$  with no free variables in common with  $L$ , and all  $N$ ,  $Q$ ,  $R$ . If  $M \sim Q$ ,  $LM = N$ , and  $N \sim R$ , then either  $\Gamma PQ = \Sigma R$  or  $\Gamma PQ = \Lambda$ .
2. Let  $N \sim \Lambda$ , then for  $L$ ,  $M$  with no free variables in common, and all  $P$ ,  $Q$ . If  $L \sim P$ ,  $M \sim Q$ , and  $LM = N$ , then  $\Gamma PQ = \Lambda$ .

The need for the two cases under the first section

- (a)  $L \sim P$ ,  $M \sim Q$ ,  $LM = N$  and  $\Gamma PQ = \Sigma R$  for  $\Sigma$  non-null
- (b)  $L \sim P$ ,  $M \sim Q$ ,  $LM = N$ ,  $N \sim R$  and  $\Gamma PQ = \Lambda$

is unfortunate but unavoidable. With respect to (a): if for  $L \sim P$  we take

$$\lambda abcd.ac(bd) \sim CCpCqrCCsqCpCsr$$