

Λ -ELIMINATION IN ILLATIVE COMBINATORY LOGIC

M. W. BUNDER

The system of propositional (and predicate) calculus developed in [2] on the basis of the illative combinatory logic in [1] lacked the Λ -elimination rules:

$$\Lambda XY \vdash X \quad (1)$$

and

$$\Lambda XY \vdash Y \quad (2)$$

which are required in some of the applications of this work.¹ In this paper we show that the axioms in [1] allow for stronger versions of

$$\Lambda XY, \mathbf{H}X, \mathbf{H}Y \vdash X \quad (3)$$

and

$$\Lambda XY, \mathbf{H}X, \mathbf{H}Y \vdash Y \quad (4)$$

(which were used in [2]) and that the addition of further axioms can lead to stronger versions still. To obtain (1) and (2), however, we need a strengthening of the axioms which alters the interpretation of the system in terms of the 3 valued truth tables given in [1].

In [1], (3) and (4) were derived using the definition of Λ :

$$\Lambda = [x, y] \mathbf{H}z \supset_z. (x \supset. y \supset x) \supset z$$

and the theorem

$$\mathbf{H}X, \mathbf{H}Y \vdash X \supset. Y \supset X \quad (5)$$

The first aim of [2], however, was to prove the general result:

1. In [3] the extra axiom:

$$\vdash \Lambda xy \supset_{x,y} x$$

was added.